

Graphical Models

Discrete Inference and Learning

MVA

2022 – 2023

<http://thoth.inrialpes.fr/~alahari/disinflearn>

Lecturers



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Organization

- 7 lectures of 3 hours each
 - Today + 24/1, 31/1, 7/2, 28/2, 7/3, 14/3
- 13:45 – 17:00 (except today) with a short break or two
- Last lecture: 14th March

<http://thoth.inrialpes.fr/~alahari/disinflearn>

Requirements

- Solid understanding of mathematical models
 - Linear algebra
 - Integral transforms
 - Differential equations
- Ideally, a basic course in discrete optimization

Topics covered

- Basic concepts, Bayesian networks, Markov random fields
- Inference algorithms: belief propagation, tree-reweighted message passing, graph cuts, move-making algorithms, Parameter learning
- Deep learning in graphical models, graph neural networks, other recent advances
- Causality

Evaluation

- Projects
- In groups of at most 3 people
- Report and presentation – Date TBD
- Topics: your own or see list on 25/1
- Bonus points for excellent class participation

What you will learn?

- Fundamental methods
- Real-world applications
- Also, pointers to using these methods in your work

Your tasks

- Following the lectures and participating actively
- Reading the literature
- Doing well in the project

Graphical Models

Discrete Inference and Learning

Lecture 1

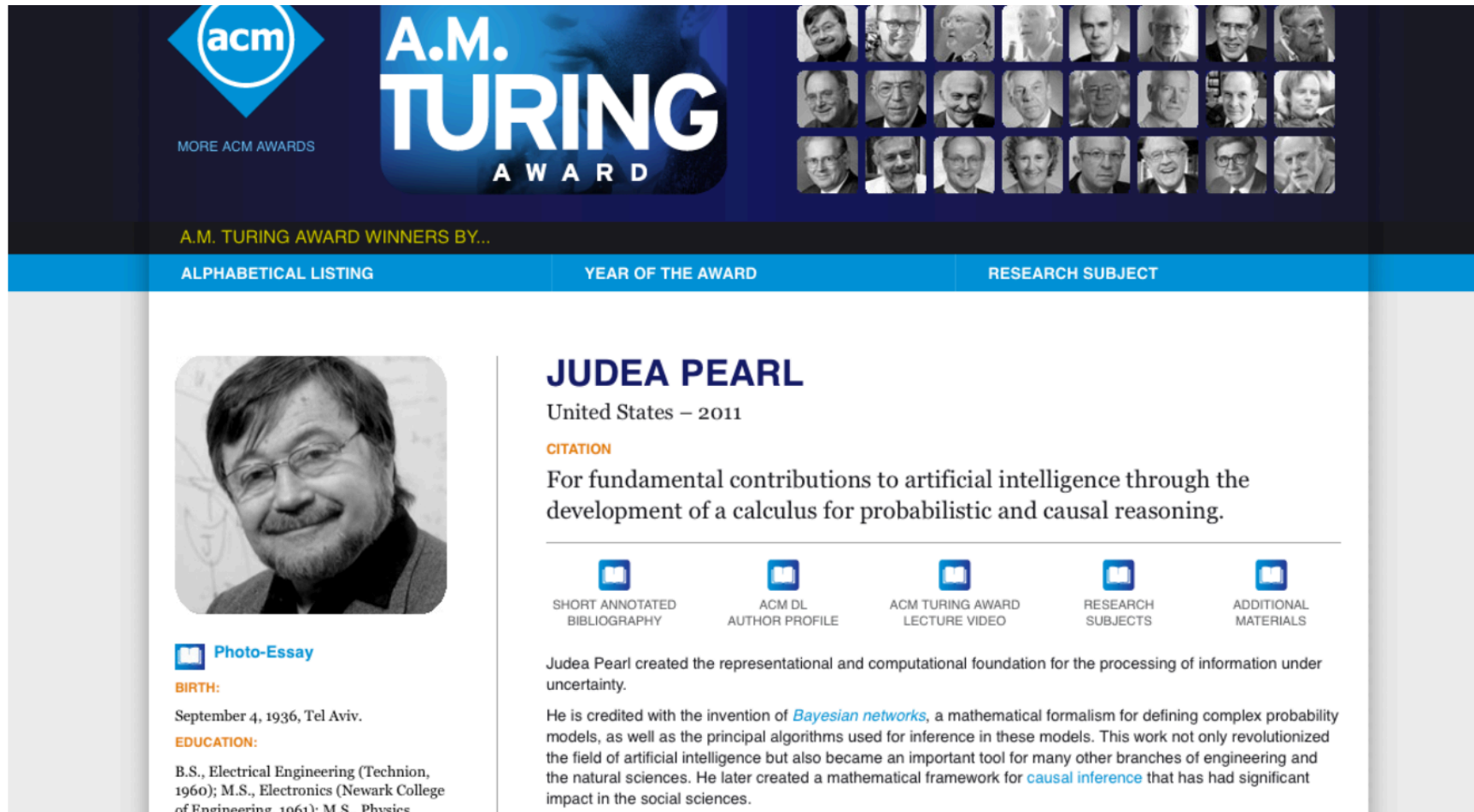
MVA

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Slides based on material from Stephen Gould, Pushmeet Kohli, Nikos Komodakis, M. Pawan Kumar, Carsten Rother, Daphne Koller, Dhruv Batra

Graphical Models ?



The screenshot shows the ACM Turing Award website. At the top, there is a navigation bar with the ACM logo, the text "A.M. TURING AWARD", and a grid of 24 small portraits of past winners. Below this is a sub-header "A.M. TURING AWARD WINNERS BY..." followed by three tabs: "ALPHABETICAL LISTING", "YEAR OF THE AWARD", and "RESEARCH SUBJECT". The main content area displays the profile of Judea Pearl, the 2011 winner. On the left is a large portrait of Pearl. To the right of the portrait, his name "JUDEA PEARL" is displayed in large blue letters, followed by "United States – 2011". Below this is a "CITATION" section with the text: "For fundamental contributions to artificial intelligence through the development of a calculus for probabilistic and causal reasoning." Underneath the citation are five icons with corresponding links: "SHORT ANNOTATED BIBLIOGRAPHY", "ACM DL AUTHOR PROFILE", "ACM TURING AWARD LECTURE VIDEO", "RESEARCH SUBJECTS", and "ADDITIONAL MATERIALS". At the bottom left of the profile, there is a "Photo-Essay" link. Below that, the "BIRTH:" information is given as "September 4, 1936, Tel Aviv." and the "EDUCATION:" information lists his degrees: "B.S., Electrical Engineering (Technion, 1960); M.S., Electronics (Newark College of Engineering, 1961); M.S., Physics".

acm
MORE ACM AWARDS

A.M. TURING AWARD

A.M. TURING AWARD WINNERS BY...

ALPHABETICAL LISTING YEAR OF THE AWARD RESEARCH SUBJECT

JUDEA PEARL

United States – 2011

CITATION

For fundamental contributions to artificial intelligence through the development of a calculus for probabilistic and causal reasoning.

SHORT ANNOTATED BIBLIOGRAPHY **ACM DL AUTHOR PROFILE** **ACM TURING AWARD LECTURE VIDEO** **RESEARCH SUBJECTS** **ADDITIONAL MATERIALS**

Photo-Essay

BIRTH:
September 4, 1936, Tel Aviv.

EDUCATION:
B.S., Electrical Engineering (Technion, 1960); M.S., Electronics (Newark College of Engineering, 1961); M.S., Physics

What this class is about?

- Making **global** predictions from **local** observations

Inference

- Learning such models from large quantities of data

Learning

Motivation

- Consider the example of medical diagnosis



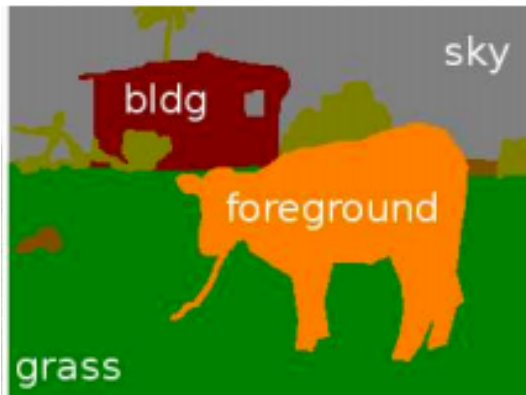
Predisposing factors
Symptoms
Test results



Diseases
Treatment outcomes

Motivation

- A very different example: image segmentation



Millions of pixels
Colours / features



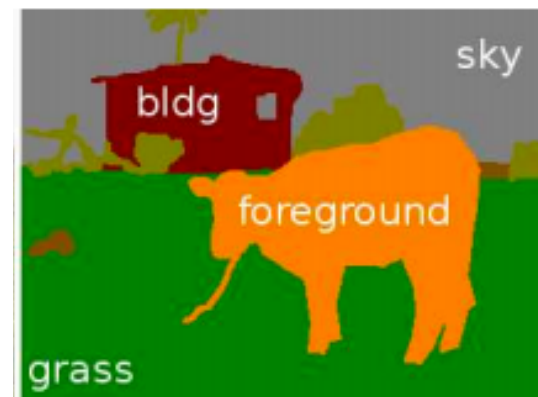
Pixel labels
{building, grass, cow, sky}

e.g., [He et al., 2004; Shotton et al., 2006; Gould et al., 2009]

Slide inspired by PGM course, Daphne Koller

Motivation

- What do these two problems have in common?



Slide inspired by PGM course, Daphne Koller

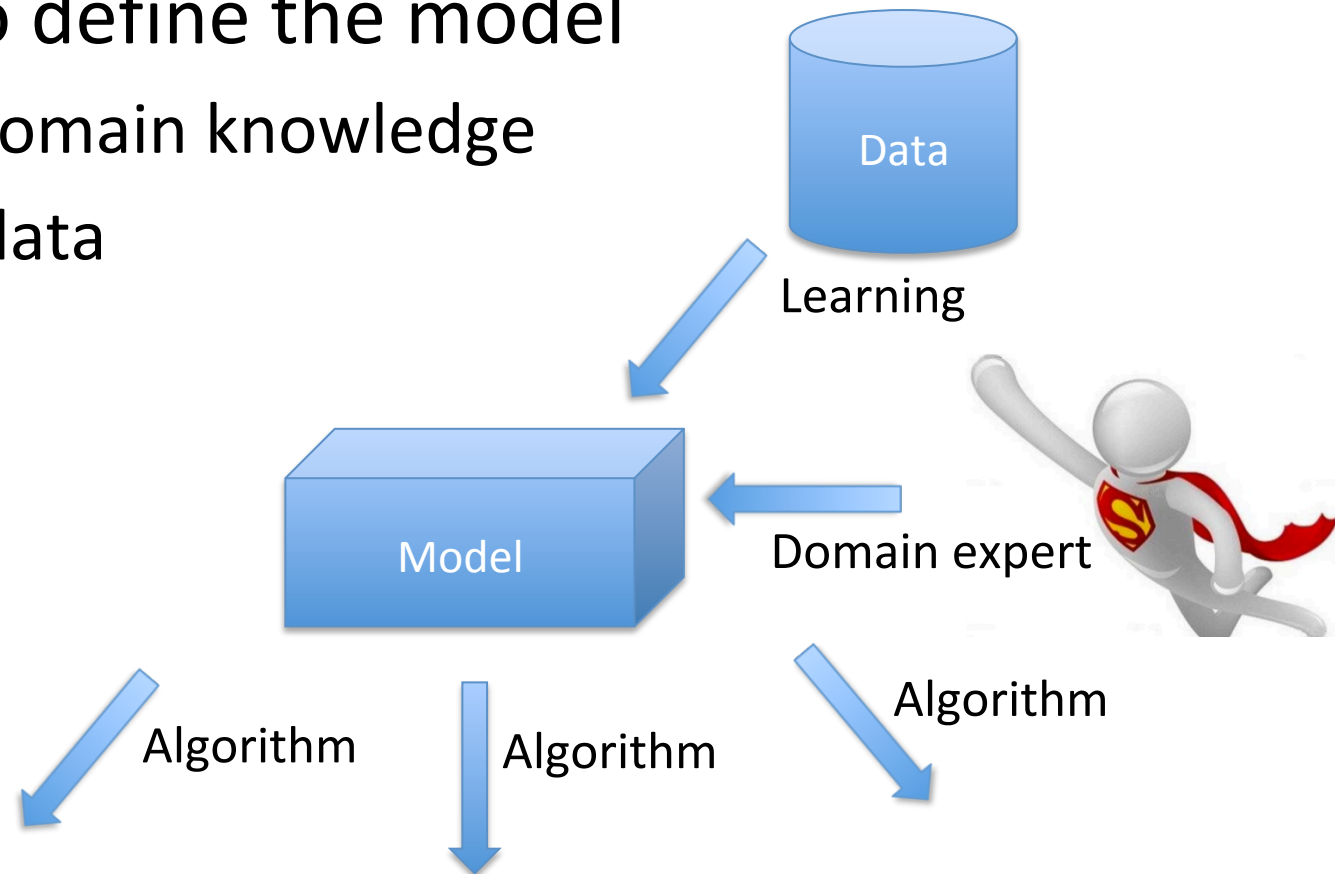
Motivation

- What do these two problems have in common?
 - Many variables
 - Uncertainty about the correct answer

Graphical Models (or Probabilistic Graphical Models)
provide a framework to address these problems

(Probabilistic) Graphical Models

- First, it is a model: a declarative representation
- Can also define the model
 - with domain knowledge
 - from data



(Probabilistic) Graphical Models

- Why probabilistic ?
- To model uncertainty
- Uncertainty due to:
 - Partial knowledge of state of the world
 - Noisy observations
 - Phenomena not observed by the model
 - Inherent stochasticity

(Probabilistic) Graphical Models

- Probability theory provides
 - Standalone representation with clear semantics
 - Reasoning patterns (conditioning, decision making)
 - Learning methods

(Probabilistic) Graphical Models

- Why graphical ?
- Intersection of ideas from probability theory and computer science
 - To represent large number of variables

Predisposing factors

Symptoms

Test results

Millions of pixels

Colours / features

Random variables Y_1, Y_2, \dots, Y_n

Goal: capture uncertainty through joint distribution $P(Y_1, \dots, Y_n)$

(Probabilistic) Graphical Models

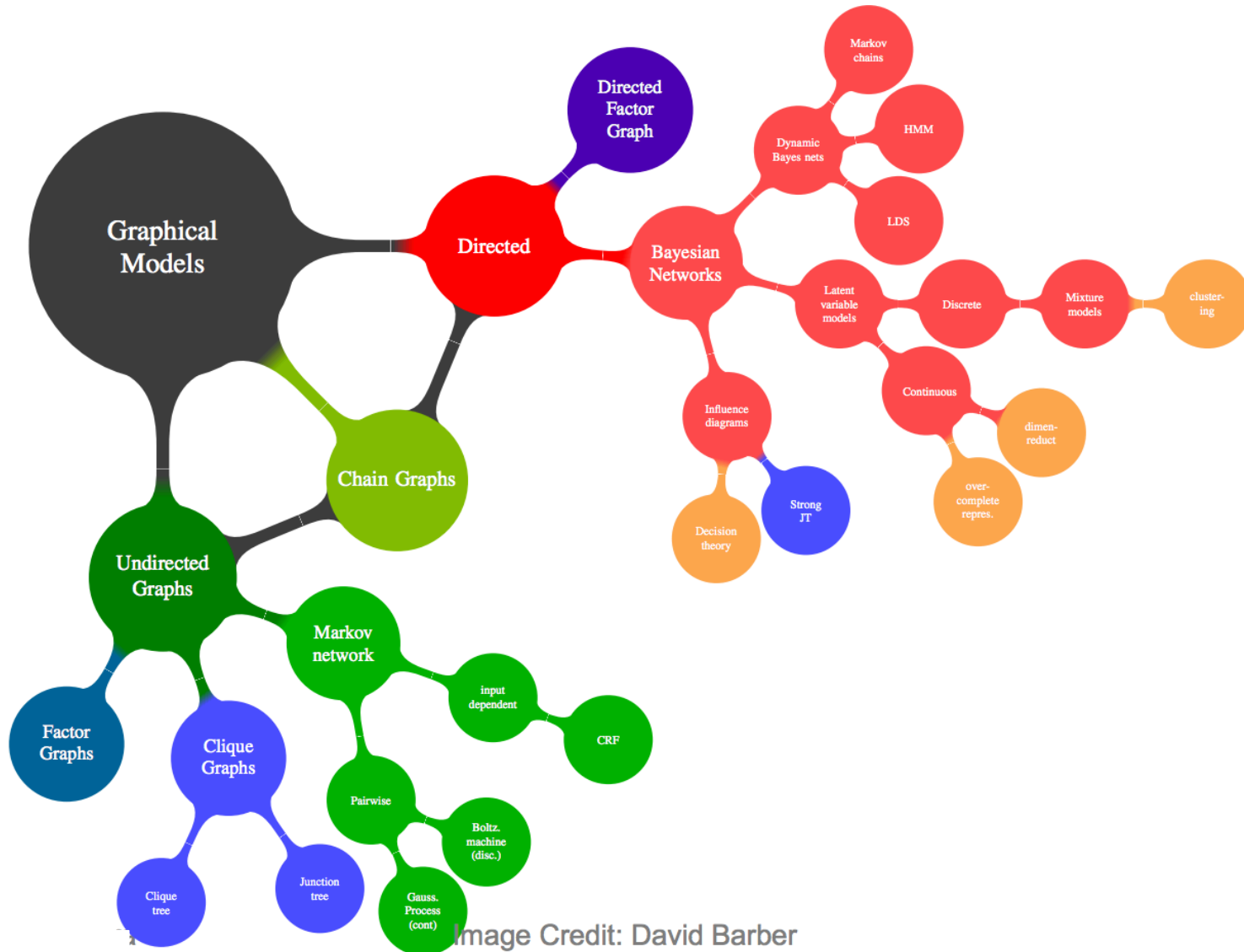
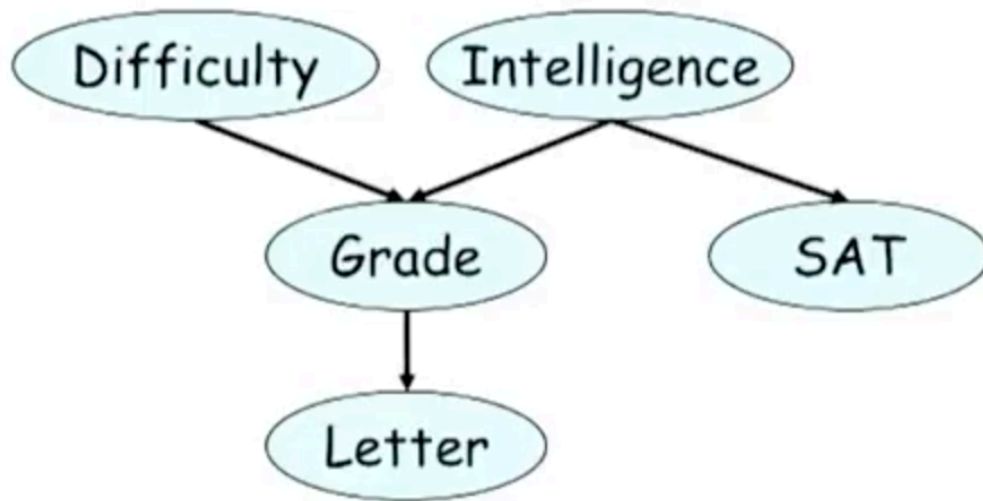


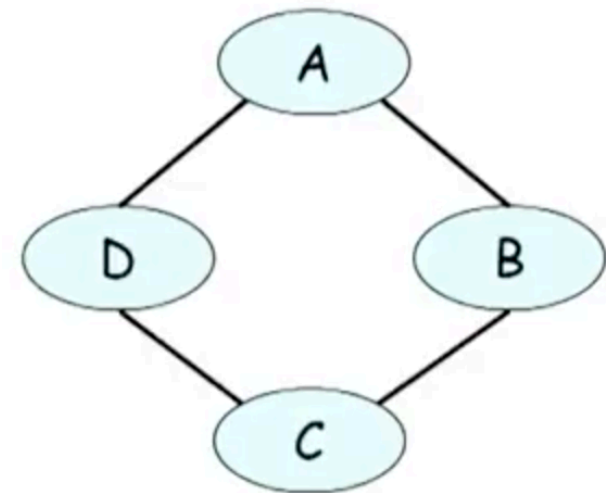
Image Credit: David Barber

(Probabilistic) Graphical Model

- Examples



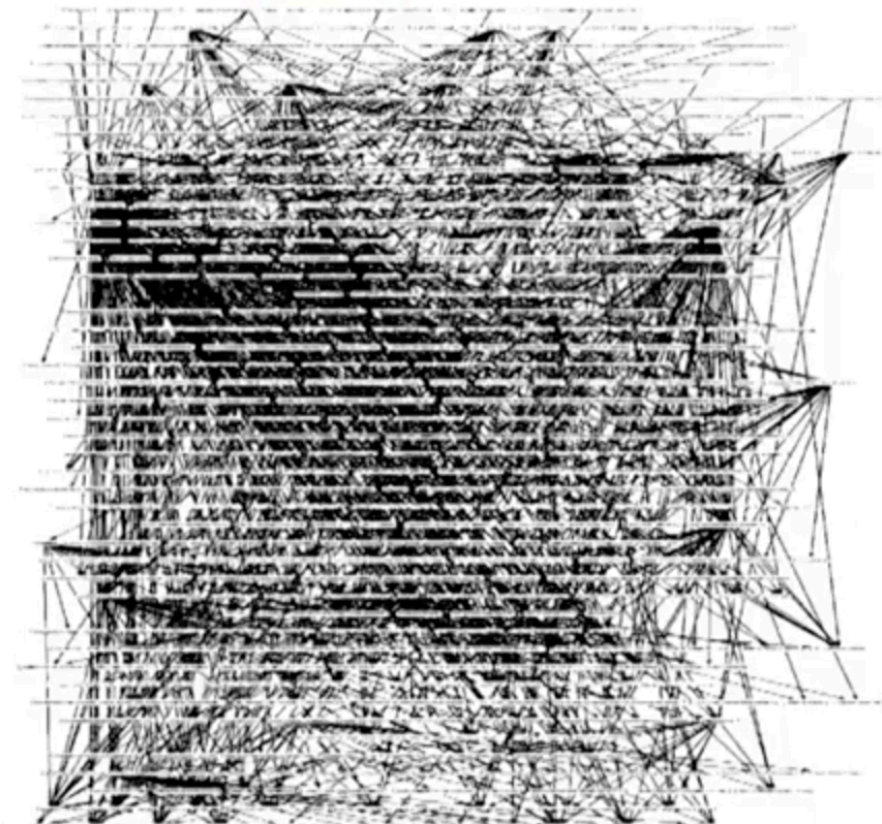
Bayesian network
(directed graph)



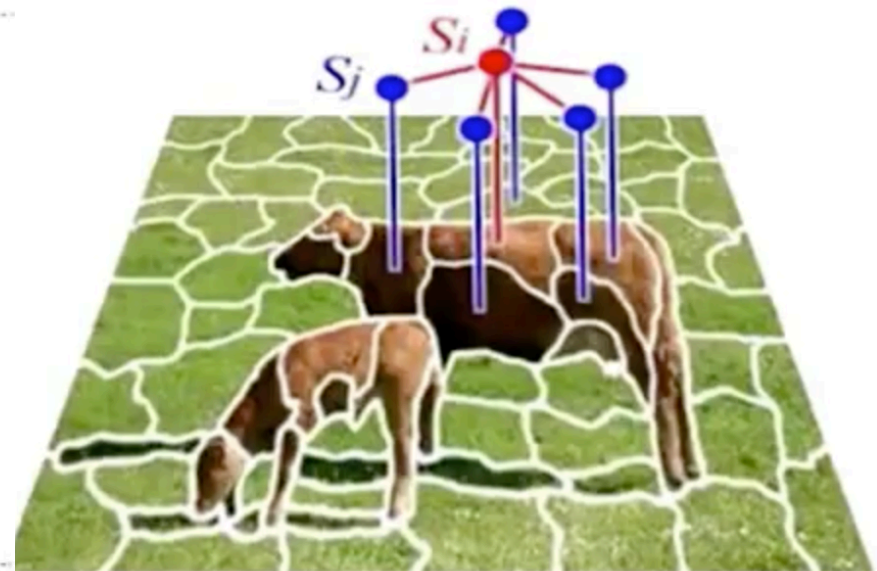
Markov network
(undirected graph)

(Probabilistic) Graphical Model

- Examples



Diagnosis network: Pradhan et al., UAI'94



Segmentation network (Courtesy D. Koller)

(Probabilistic) Graphical Model

- Intuitive & compact data structure
- Efficient reasoning through general-purpose algorithms
- Sparse parameterization
 - Through expert knowledge, or
 - Learning from data

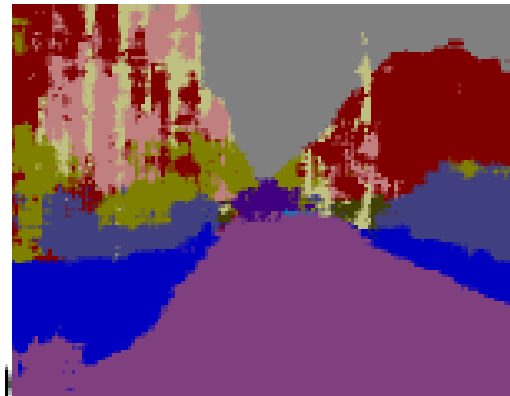
(Probabilistic) Graphical Model

- Many many applications
 - Medical diagnosis
 - Fault diagnosis
 - Natural language processing
 - Traffic analysis
 - Social network models
 - Message decoding
 - Computer vision: segmentation, 3D, pose estimation
 - Speech recognition
 - Robot localization & mapping

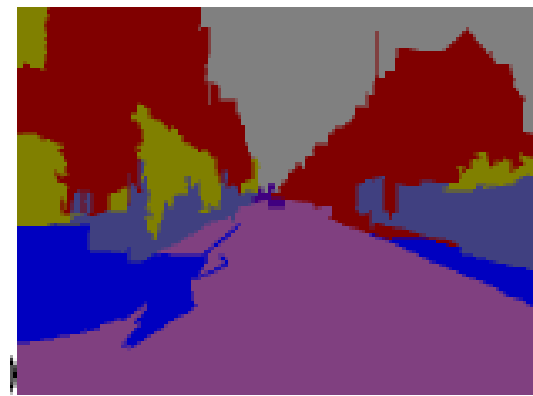
Image segmentation



Image



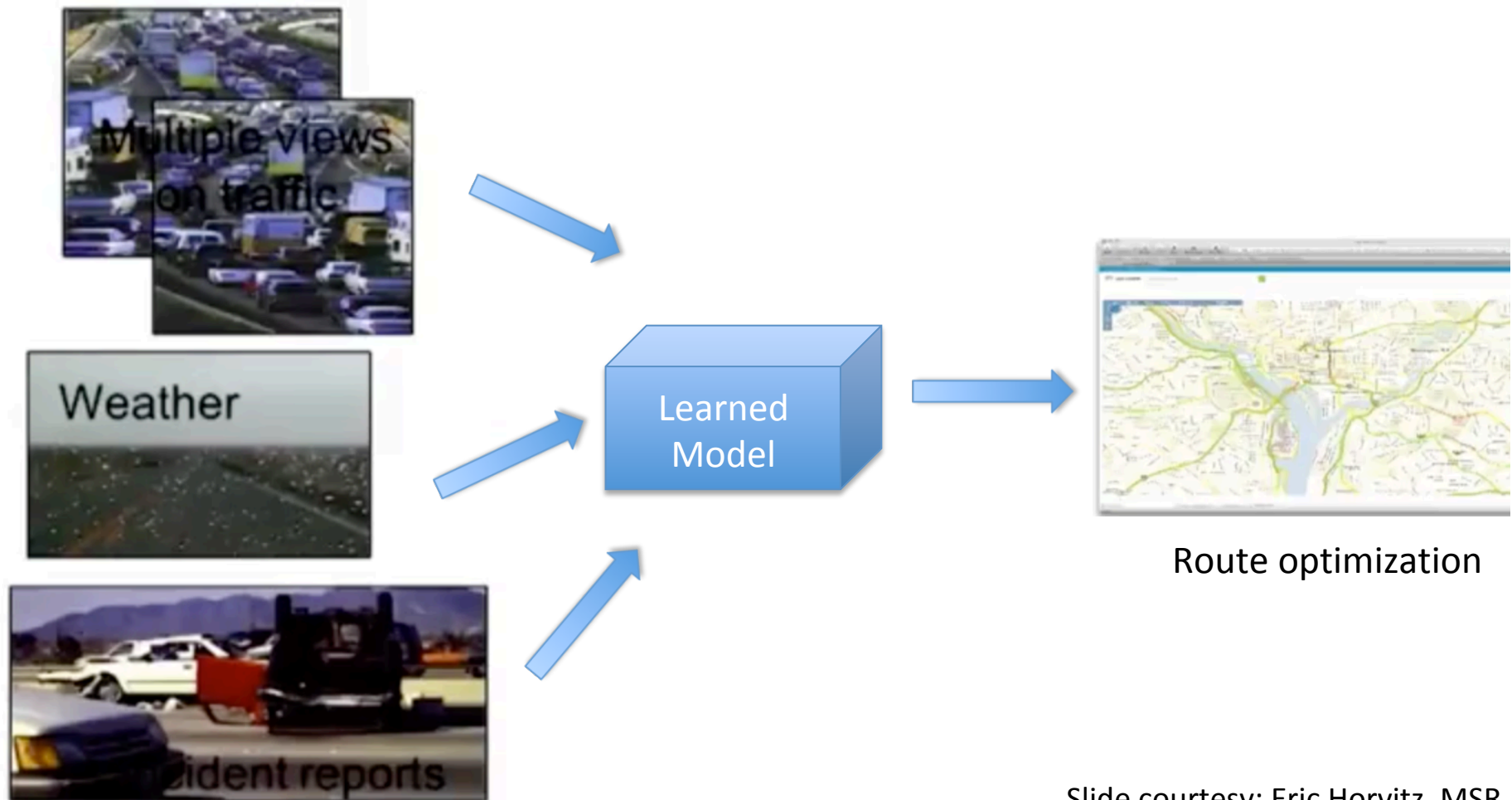
No graphical model



With graphical model

Multi-sensor integration: Traffic

- Learn from historical data to make predictions



Going global: Local ambiguity

- Text recognition

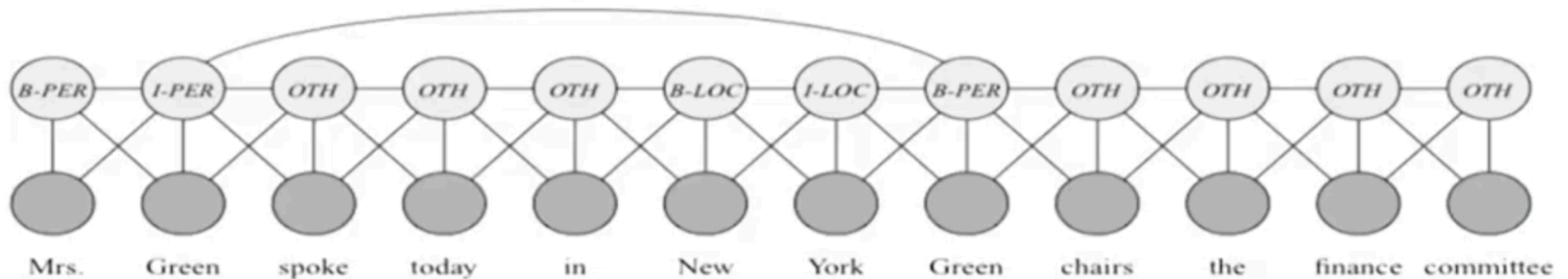
The image shows the words 'THE' and 'CAT' in a simple, blocky font. The letters are colored in two ways: green and red. In 'THE', the 'T' and 'E' are green, while the 'H' is red. In 'CAT', the 'C' and 'T' are green, while the 'A' is red. This creates a visual ambiguity where the words can be read as 'THE' and 'CAT' or 'TAE' and 'CAT'.

Smyth et al., 1994

Going global: Local ambiguity

- Textual information extraction

e.g., Mrs. Green spoke today in New York. Green chairs the financial committee.



Overview

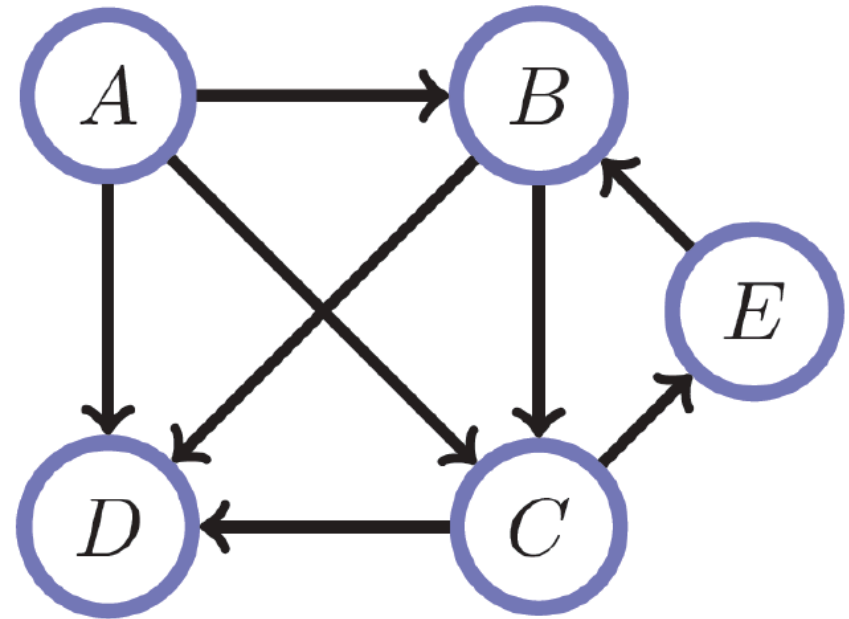
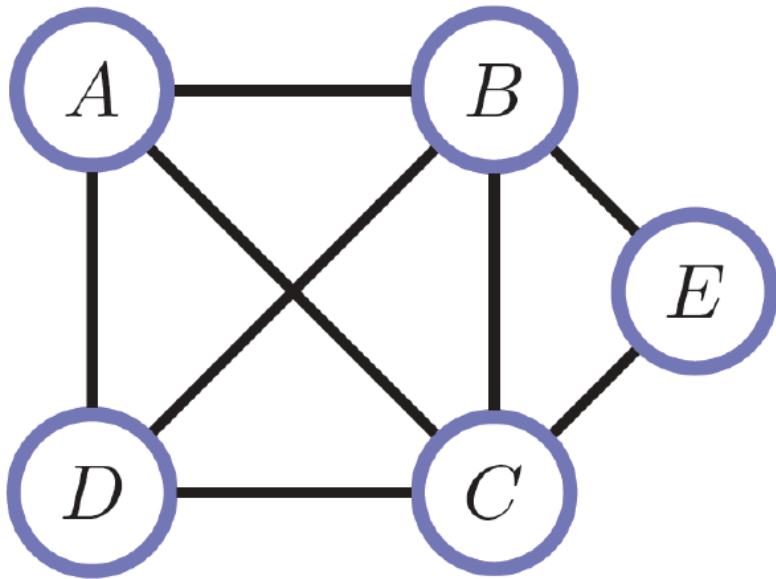
- Representation
 - How do we store $P(Y_1, \dots, Y_n)$
 - Directed and undirected (model implications/assumptions)
- Inference
 - Answer questions with the model
 - Exact and approximate (marginal/most probable estimate)
- Learning
 - What model is right for data
 - Parameters and structure

First, a recap of basics

Graphs

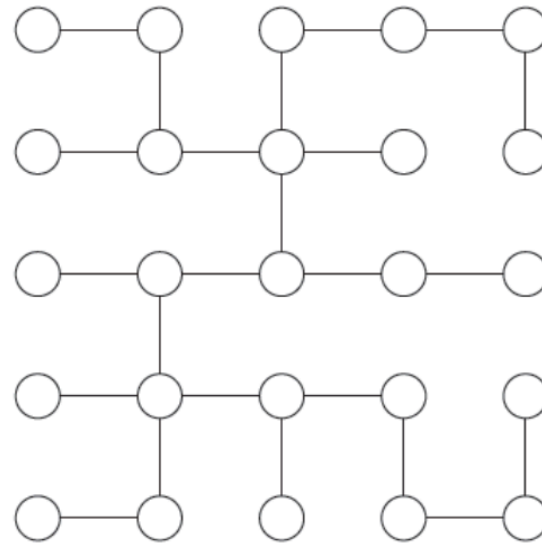
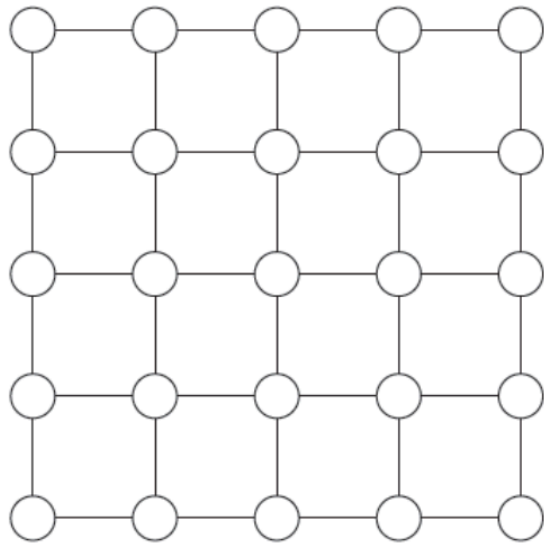
- Concepts
 - Definition of G
 - Vertices/Nodes
 - Edges
 - Directed vs Undirected
 - Neighbours vs Parent/Child
 - Degree vs In/Out degree
 - Walk vs Path vs Cycle

Graphs

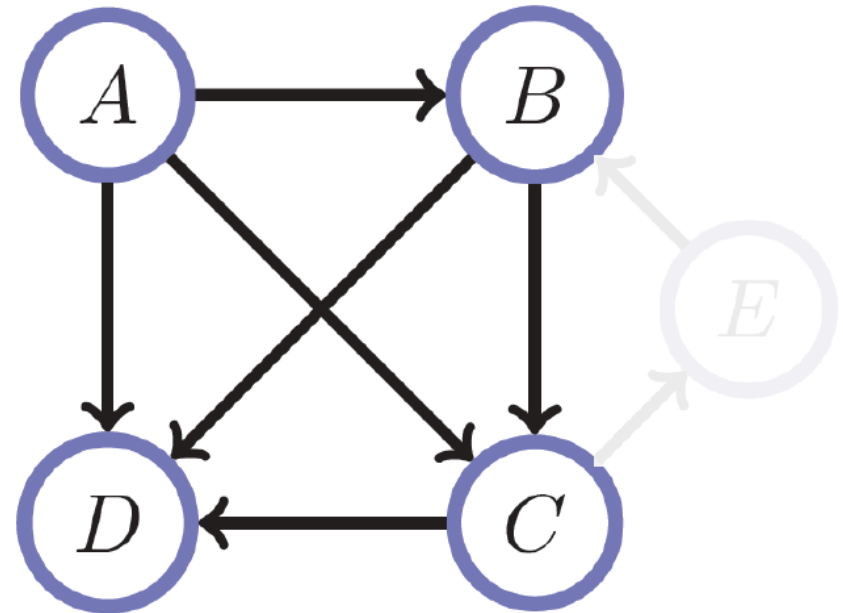
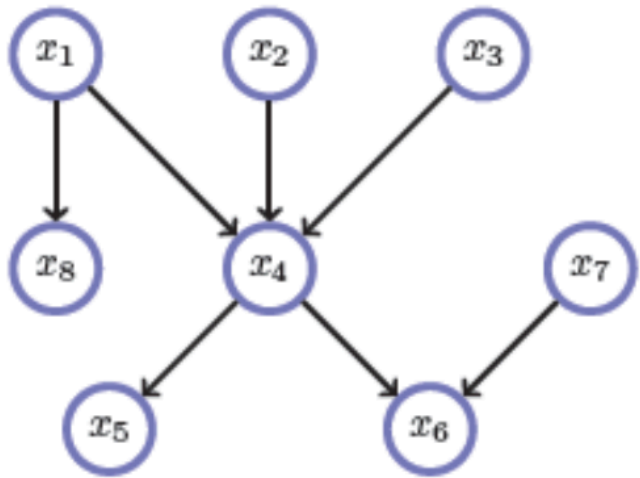


Special graphs

- Trees: undirected graph, no cycles
- Spanning tree: Same set of vertices, but subset of edges, connected and no cycles



Directed acyclic graphs (DAGs)



Joint distribution

- 3 variables
 - Intelligence (I)
 - Difficulty (D)
 - Grade (G)
- Independent parameters?

I	D	G	Prob.
i^0	d^0	g^1	0.126
i^0	d^0	g^2	0.168
i^0	d^0	g^3	0.126
i^0	d^1	g^1	0.009
i^0	d^1	g^2	0.045
i^0	d^1	g^3	0.126
i^1	d^0	g^1	0.252
i^1	d^0	g^2	0.0224
i^1	d^0	g^3	0.0056
i^1	d^1	g^1	0.06
i^1	d^1	g^2	0.036
i^1	d^1	g^3	0.024

Conditioning

- Condition on g^1

I	D	G	Prob.
i^0	d^0	g^1	0.126
i^0	d^0	g^2	0.168
i^0	d^0	g^3	0.126
i^0	d^1	g^1	0.009
i^0	d^1	g^2	0.045
i^0	d^1	g^3	0.126
i^1	d^0	g^1	0.252
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i^1	d^0	g^3	0.0056
i^1	d^1	g^1	0.06
i^1	d^1	g^2	0.036
i^1	d^1	g^3	0.024

Conditioning

- $P(Y = y \mid X = x)$
- Informally,
 - What do you believe about $Y=y$ when I tell you $X=x$?
- $P(\text{France wins a football tournament in 2023}) ?$
- What if I tell you:
 - France almost won the world cup 2022
 - Hasn't had catastrophic results since 😊

Conditioning: Reduction

- Condition on g^1

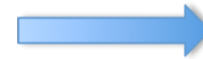
I	D	G	Prob.
i^0	d^0	g^1	0.126
i^0	d^1	g^1	0.009
i^1	d^0	g^1	0.252
i^1	d^1	g^1	0.06

Conditioning: Renormalization

I	D	G	Prob.
i^0	d^0	g^1	0.126
i^0	d^1	g^1	0.009
i^1	d^0	g^1	0.252
i^1	d^1	g^1	0.06

$P(\mathbf{I}, \mathbf{D}, g^1)$

Unnormalized measure



I	D	Prob.
i^0	d^0	0.282
i^0	d^1	0.02
i^1	d^0	0.564
i^1	d^1	0.134

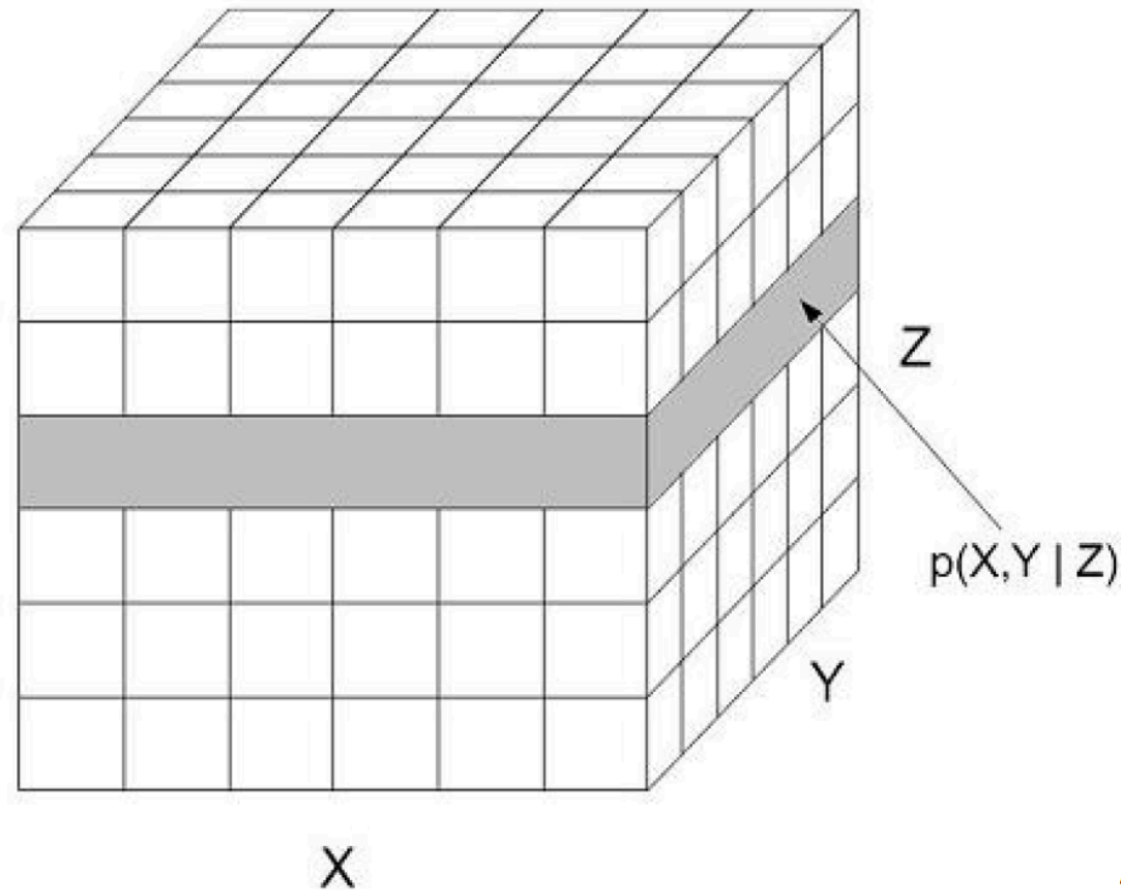
$P(\mathbf{I}, \mathbf{D} \mid g^1)$

Conditional probability distribution

- Example $P(G \mid I, D)$

	g^1	g^2	g^3
i^0, d^0	0.3	0.4	0.3
i^0, d^1	0.05	0.25	0.7
i^1, d^0	0.9	0.08	0.02
i^1, d^1	0.5	0.3	0.2

Conditional probability distribution



$$p(x, y | Z = z) = \frac{p(x, y, z)}{p(z)}$$

Marginalization

$P(I,D)$

I	D	Prob.
i^0	d^0	0.282
i^0	d^1	0.02
i^1	d^0	0.564
i^1	d^1	0.134

Marginalize I

D	Prob.
d^0	0.846
d^1	0.154

Marginalization

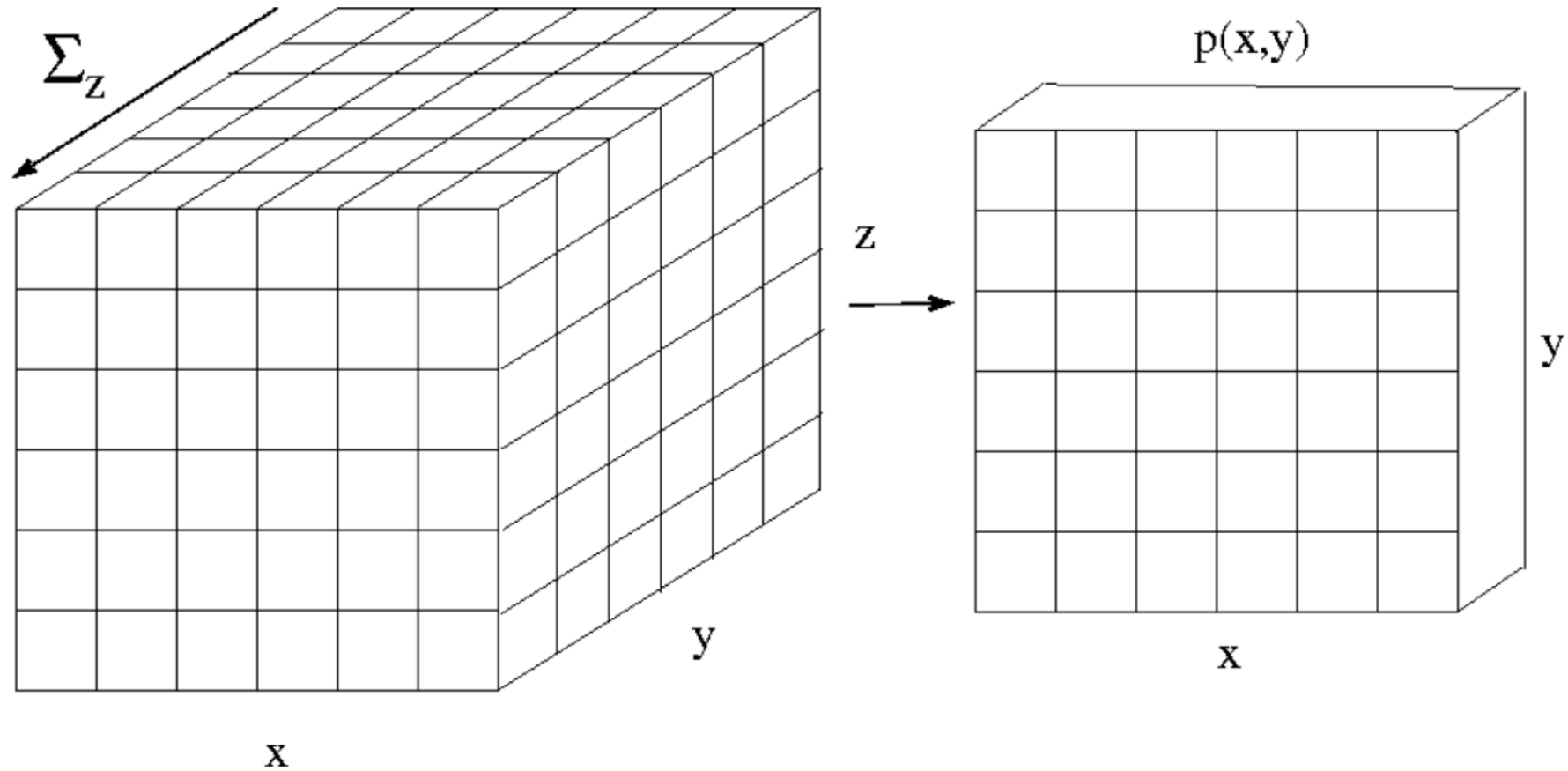
- Events

- $P(A) = P(A \text{ and } B) + P(A \text{ and not } B)$

- Random variables

- $P(X = x) = \sum_y P(X = x, Y = y)$

Marginalization



$$p(x, y) = \sum_{z \in \mathcal{Z}} p(x, y, z)$$

$$p(x) = \sum_{y \in \mathcal{Y}} p(x, y)$$

Slide courtesy: Erik Sudderth

Factors

- A factor $\Phi(Y_1, \dots, Y_k)$

$$\Phi: \text{Val}(Y_1, \dots, Y_k) \rightarrow \mathbb{R}$$

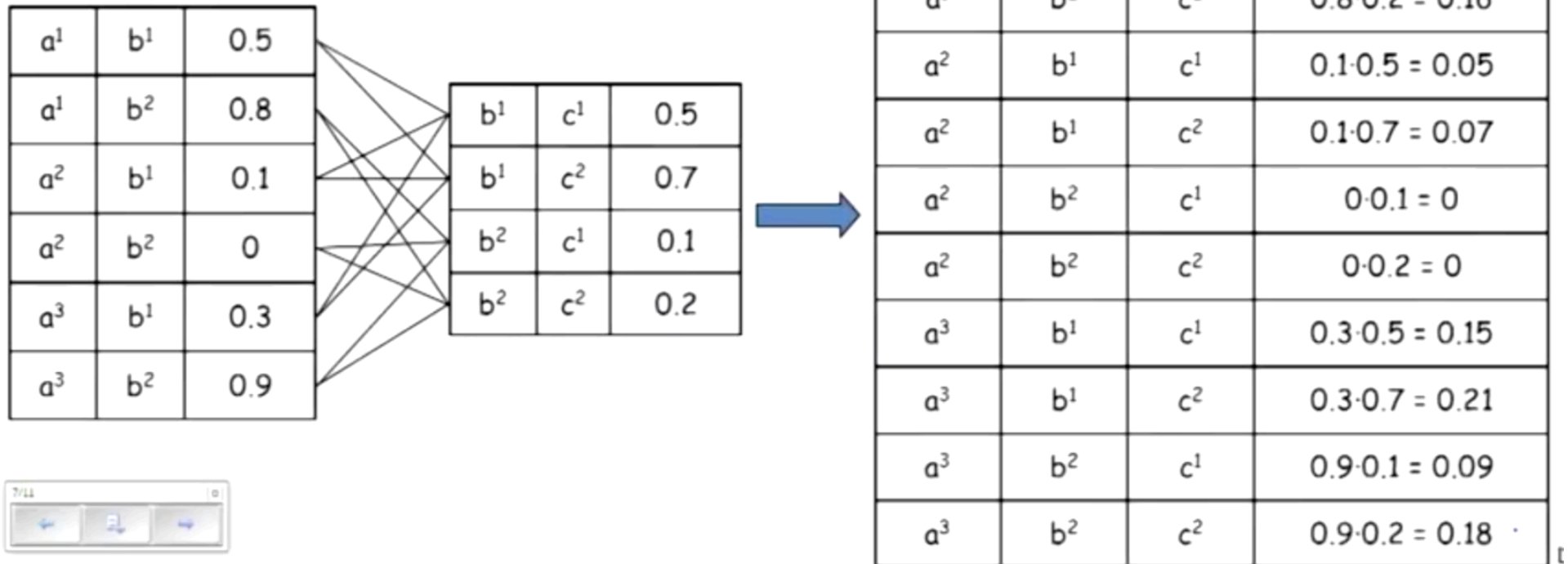
- Scope = $\{Y_1, \dots, Y_k\}$

General factors

- Not necessarily for probabilities

A	B	ϕ
a^0	b^0	30
a^0	b^1	5
a^1	b^0	1
a^1	b^1	10

Factor product



Factor marginalization

The diagram illustrates the process of factor marginalization. It shows a large table on the left representing a joint distribution over variables a , b , and c . The rows are grouped by the value of a (a^1 , a^2 , a^3). The columns are a , b , c , and a numerical value. Lines connect the rows of the large table to the rows of a smaller table on the right, which represents the marginal distribution over a and c . The b variable has been marginalized out, and the numerical values are the sum of the corresponding rows in the large table.

a^1	b^1	c^1	0.25
a^1	b^1	c^2	0.35
a^1	b^2	c^1	0.08
a^1	b^2	c^2	0.16
a^2	b^1	c^1	0.05
a^2	b^1	c^2	0.07
a^2	b^2	c^1	0
a^2	b^2	c^2	0
a^3	b^1	c^1	0.15
a^3	b^1	c^2	0.21
a^3	b^2	c^1	0.09
a^3	b^2	c^2	0.18

a^1	c^1	0.33
a^1	c^2	0.51
a^2	c^1	0.05
a^2	c^2	0.07
a^3	c^1	0.24
a^3	c^2	0.39

Factor reduction

a^1	b^1	c^1	0.25
a^1	b^1	c^2	0.35
a^1	b^2	c^1	0.08
a^1	b^2	c^2	0.16
a^2	b^1	c^1	0.05
a^2	b^1	c^2	0.07
a^2	b^2	c^1	0
a^2	b^2	c^2	0
a^3	b^1	c^1	0.15
a^3	b^1	c^2	0.21
a^3	b^2	c^1	0.09
a^3	b^2	c^2	0.18

a^1	b^1	c^1	0.25
a^1	b^2	c^1	0.08
a^2	b^1	c^1	0.05
a^2	b^2	c^1	0
a^3	b^1	c^1	0.15
a^3	b^2	c^1	0.09

Why factors ?

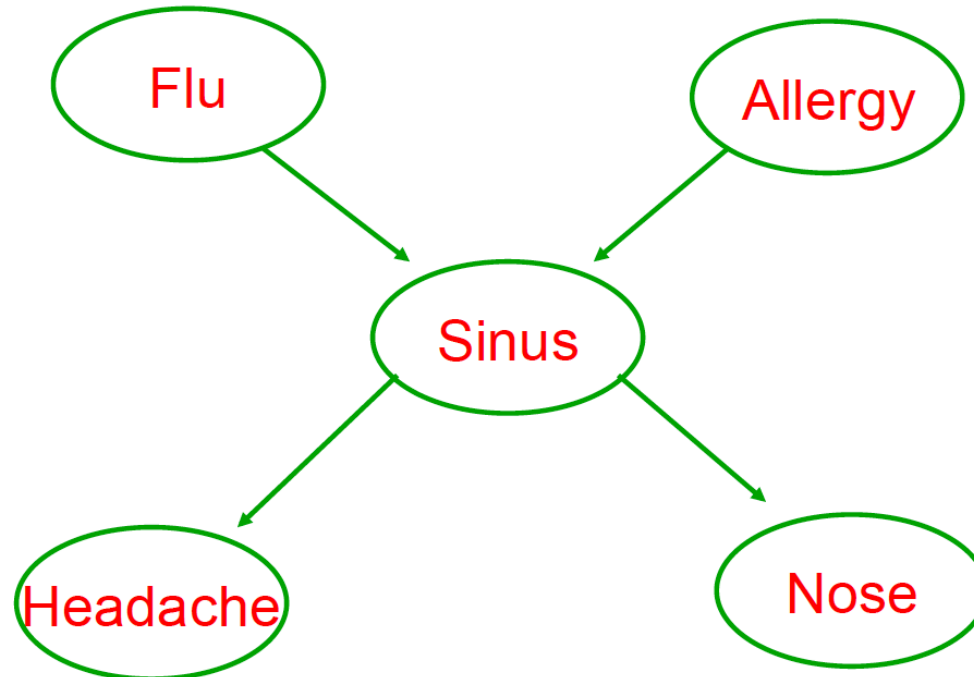
- Building blocks for defining distributions in high-dimensional spaces
- Set of basic operations for manipulating these distributions

Bayesian Networks

- DAGs
 - nodes represent variables in the Bayesian sense
 - edges represent conditional dependencies
- Example
 - Suppose that we know the following:
 - The flu causes sinus inflammation
 - Allergies cause sinus inflammation
 - Sinus inflammation causes a runny nose
 - Sinus inflammation causes headaches
 - How are these connected ?

Bayesian Networks

- Example



Bayesian Networks

- A general Bayes net
 - Set of random variables
 - DAG: encodes independence assumptions
 - Conditional probability trees
 - Joint distribution

$$P(Y_1, \dots, Y_n) = \prod_{i=1}^n P(Y_i \mid \text{Pa}_{Y_i})$$

Bayesian Networks

- A general Bayes net
 - How many parameters ?
 - Discrete variables Y_1, \dots, Y_n
 - Graph: Defines parents of Y_i , i.e., (Pa_{Y_i})
 - CPTs: $P(Y_i | Pa_{Y_i})$

Markov nets

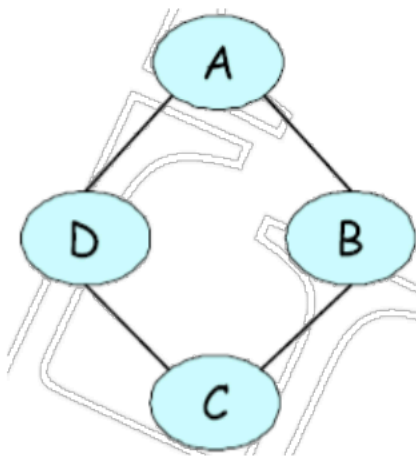
- Set of random variables
- Undirected graph
 - Encodes independence assumptions
- Factors

Comparison to Bayesian Nets ?

Pairwise MRFs

- Composed of pairwise factors
 - A function of two variables
 - Can also have unary terms

- Example



$\phi_1[A, B]$			$\phi_2[B, C]$			$\phi_3[C, D]$			$\phi_4[D, A]$		
a^0	b^0	30	b^0	c^0	100	c^0	d^0	1	d^0	a^0	100
a^0	b^1	5	b^0	c^1	1	c^0	d^1	100	d^0	a^1	1
a^1	b^0	1	b^1	c^0	1	c^1	d^0	100	d^1	a^0	1
a^1	b^1	10	b^1	c^1	100	c^1	d^1	1	d^1	a^1	100

Markov Nets: Computing probabilities

- Can only compute ratio of probabilities directly

$\phi_1[A, B]$			$\phi_2[B, C]$			$\phi_3[C, D]$			$\phi_4[D, A]$		
a^0	b^0	30	b^0	c^0	100	c^0	d^0	1	d^0	a^0	100
a^0	b^1	5	b^0	c^1	1	c^0	d^1	100	d^0	a^1	1
a^1	b^0	1	b^1	c^0	1	c^1	d^0	100	d^1	a^0	1
a^1	b^1	10	b^1	c^1	100	c^1	d^1	1	d^1	a^1	100

- Need to normalize with a **partition function**
 - Hard ! (sum over all possible assignments)
- In Bayesian Nets, can do by multiplying CPTs

Markov nets \leftrightarrow Factorization

- Given an undirected graph H over variables $Y = \{Y_1, \dots, Y_n\}$
- A distribution P factorizes over H if there exist
 - Subsets of variables $S^i \subseteq Y$ s.t. S^i are fully-connected in H
 - Non-negative potentials (factors) $\Phi_1(S^1), \dots, \Phi_m(S^m)$: clique potentials
 - Such that

$$P(Y_1, \dots, Y_n) = \frac{1}{Z} \prod_{i=1}^m \Phi_i(S^i)$$

Conditional Markov Random Fields

- Also known as: Markov networks, undirected graphical models, MRFs
- Note: Not making a distinction between CRFs and MRFs
- $\mathbf{X} \in \mathcal{X}$: observed random variables
- $\mathbf{Y} = (Y_1, \dots, Y_n) \in \mathcal{Y}$: output random variables
- \mathbf{Y}_c are subset of variables for clique $c \subseteq \{1, \dots, n\}$
- Define a factored probability distribution

$$P(\mathbf{Y} | \mathbf{X}) = \frac{1}{Z(\mathbf{X})} \prod_c \psi_c(\mathbf{Y}_c; \mathbf{X})$$

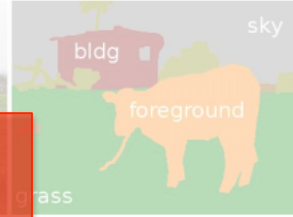
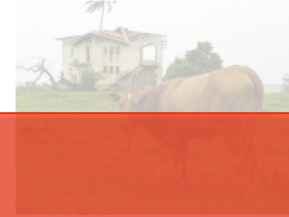
Partition function = $\sum_{\mathbf{Y} \in \mathcal{Y}} \prod_c \psi_c(\mathbf{Y}_c; \mathbf{X})$ **Exponential number of configurations !**

MRFs / CRFs

- Several applications, e.g., computer vision



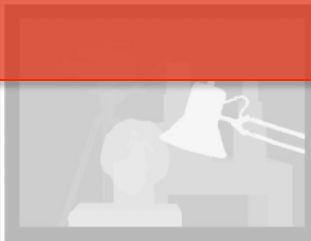
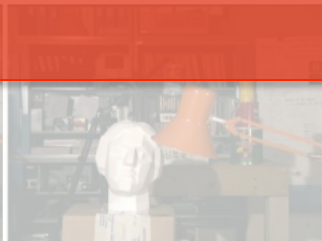
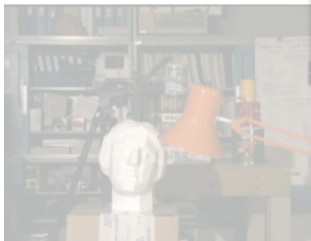
Interactive figure-ground segmentation [Boykov and Jolly, 2001; Boykov and Funka-Lea, 2003]



Surface context [Hoiem et al., 2005]

Semantic labeling [He et al., 2004; Shotton et al., 2006; Gould et al., 2005]

Low-level vision problems



Stereo matching [Kolmogorov and Zabih, 2001; Scharstein and Szeliski, 2002]

Image denoising [Felzenszwalb and Huttenlocher 2004]

MRFs / CRFs

- Several applications, e.g., computer vision

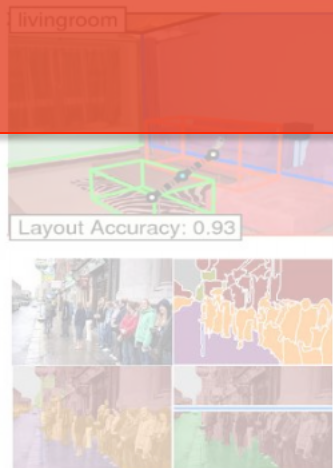
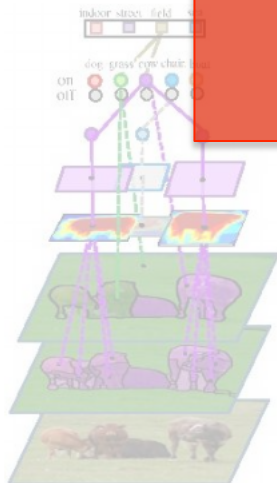


Object detection [Felzenszwalb et al., 2008]



Human pose estimation [Lipton and Black, 2015; Ramakrishna et al., 2012]

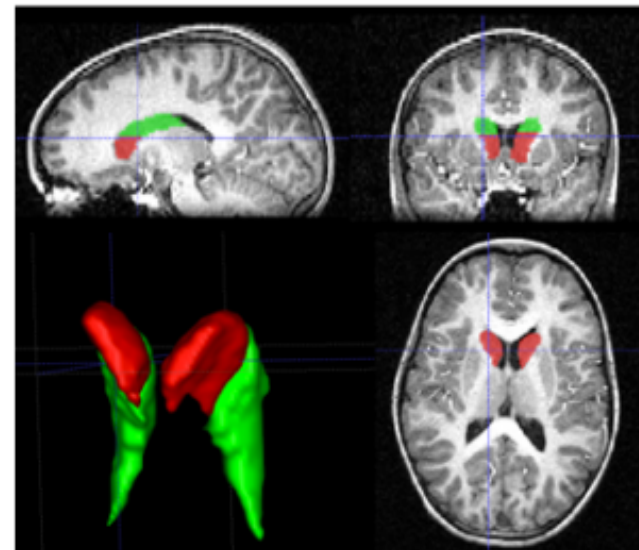
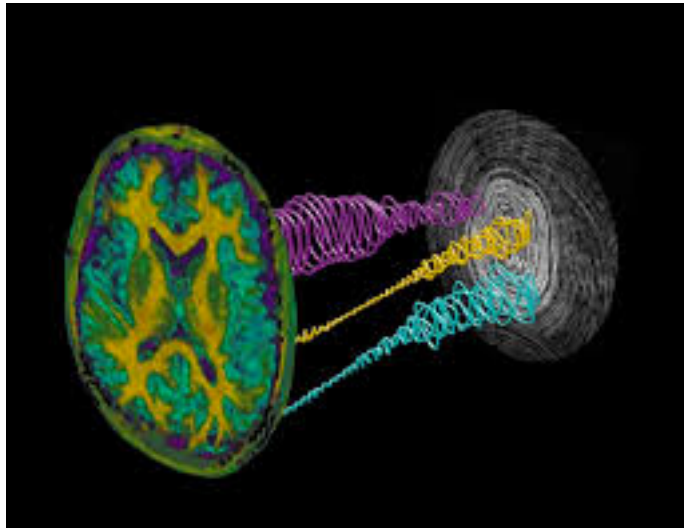
High-level vision problems



Scene understanding
[Fouhey et al., 2014; Ladicky et al., 2010;
Xiao et al., 2013; Yao et al., 2012]

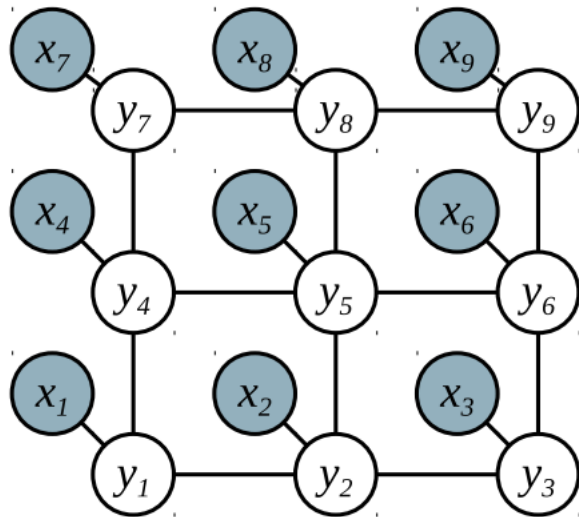
MRFs / CRFs

- Several applications, e.g., medical imaging

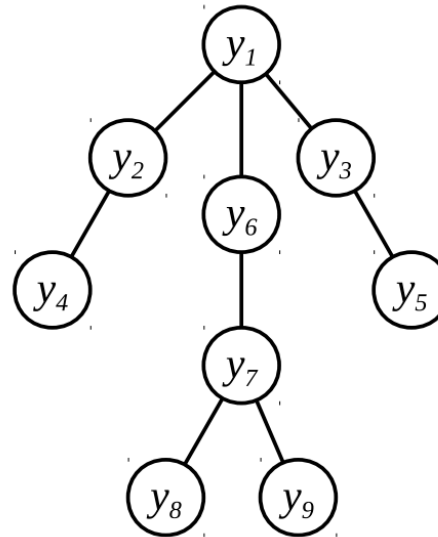


MRFs / CRFs

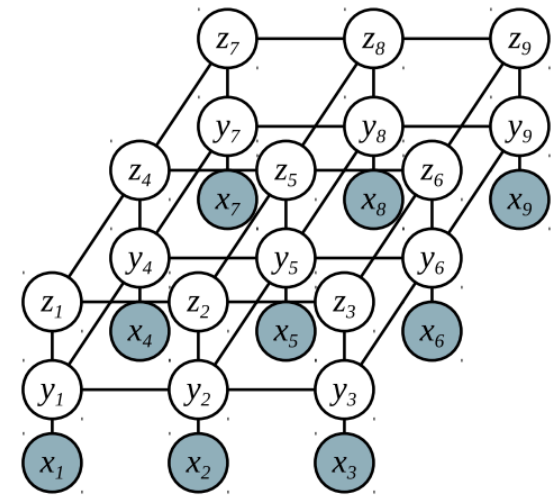
- Inherent in all these problems are graphical models



Pixel labeling



Object detection
Pose estimation



Scene understanding

Maximum a posteriori (MAP) inference

$$\begin{aligned}\mathbf{y}^* &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} P(\mathbf{y} \mid \mathbf{x}) \\ &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} \frac{1}{Z(\mathbf{x})} \prod_c \psi_c(\mathbf{Y}_c; \mathbf{X}) \\ &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} \log \left(\frac{1}{Z(\mathbf{x})} \prod_c \psi_c(\mathbf{Y}_c; \mathbf{X}) \right) \\ &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} \sum_c \log \psi_c(\mathbf{Y}_c; \mathbf{X}) - \log Z(\mathbf{X}) \\ &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} \sum_c \log \psi_c(\mathbf{Y}_c; \mathbf{X}) \rightarrow -E(\mathbf{Y}; \mathbf{X})\end{aligned}$$

Maximum a posteriori (MAP) inference

$$\begin{aligned}\mathbf{y}^* &= \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} P(\mathbf{y} \mid \mathbf{x}) = \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} \sum_c \log \Psi_c(\mathbf{Y}_c; \mathbf{X}) \\ &= \operatorname{argmin}_{\mathbf{y} \in \mathcal{Y}} E(\mathbf{y}; \mathbf{x})\end{aligned}$$

MAP inference \Leftrightarrow Energy minimization

The energy function is $E(\mathbf{Y}; \mathbf{X}) = \sum_c \psi_c(\mathbf{Y}_c; \mathbf{X})$

where $\psi_c(\cdot) = -\log \Psi_c(\cdot)$

 Clique potential

Clique potentials

- Defines a mapping from an assignment of random variables to a real number

$$\psi_c : \mathcal{Y}_c \times \mathcal{X} \rightarrow \mathbb{R}$$

- Encodes a preference for assignments to the random variables (lower is better)
- Parameterized as $\psi_c(\mathbf{y}_c; \mathbf{x}) = \mathbf{w}_c^T \phi_c(\mathbf{y}_c; \mathbf{x})$

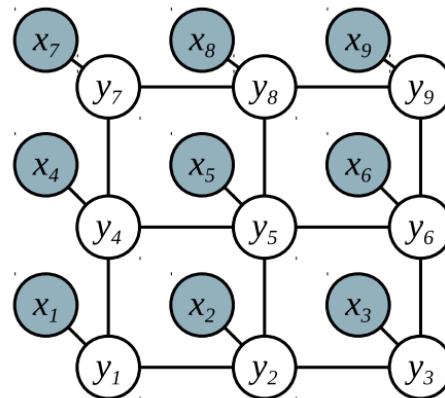


Parameters

Clique potentials

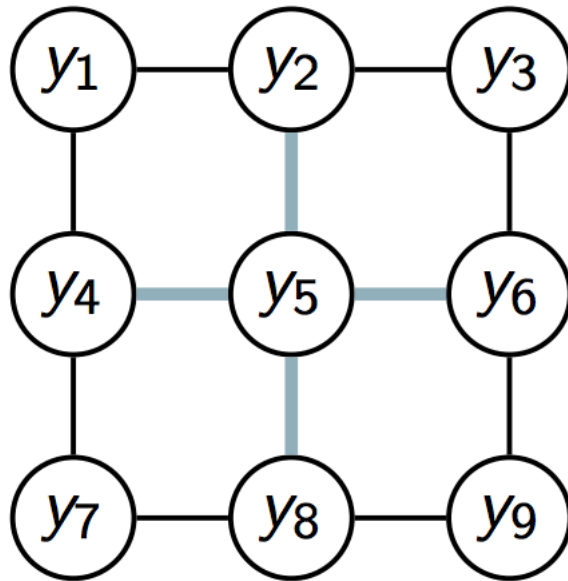
- Arity

$$\begin{aligned} E(\mathbf{y}; \mathbf{x}) &= \sum_c \psi_c(\mathbf{y}_c; \mathbf{x}) \\ &= \underbrace{\sum_{i \in \mathcal{V}} \psi_i^U(y_i; \mathbf{x})}_{\text{unary}} + \underbrace{\sum_{ij \in \mathcal{E}} \psi_{ij}^P(y_i, y_j; \mathbf{x})}_{\text{pairwise}} + \underbrace{\sum_{c \in \mathcal{C}} \psi_c^H(\mathbf{y}_c; \mathbf{x})}_{\text{higher-order}}. \end{aligned}$$

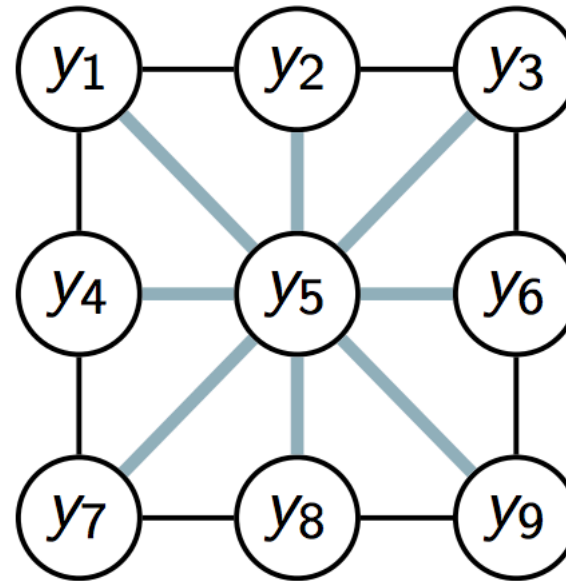


Clique potentials

- Arity

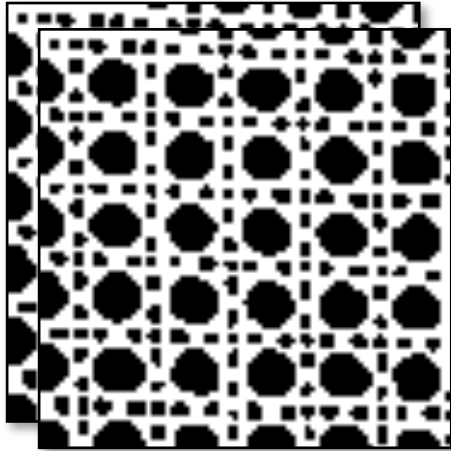


4-connected, \mathcal{N}_4

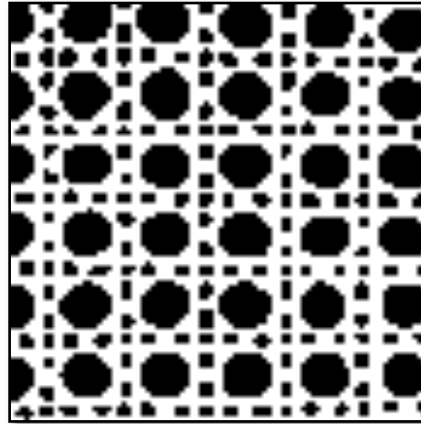


8-connected, \mathcal{N}_8

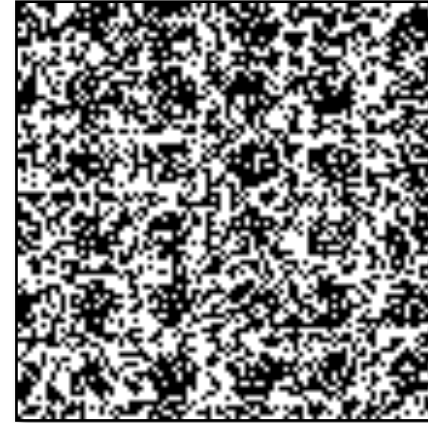
Reason 1: Texture modelling



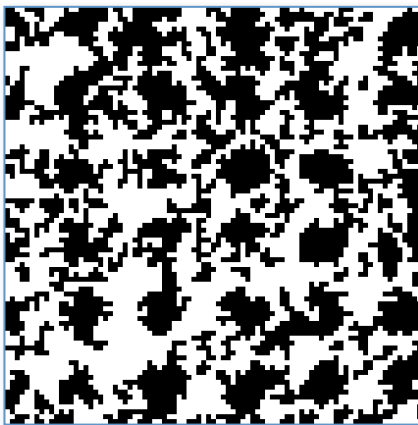
Training images



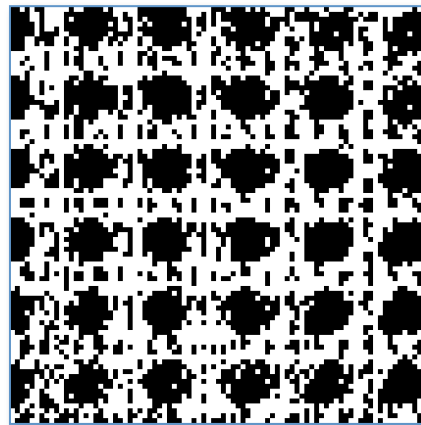
Test image



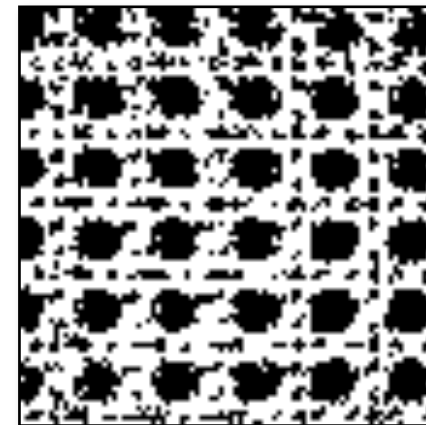
Test image (60% Noise)



Result MRF
4-connected
(neighbours)



Result MRF
4-connected

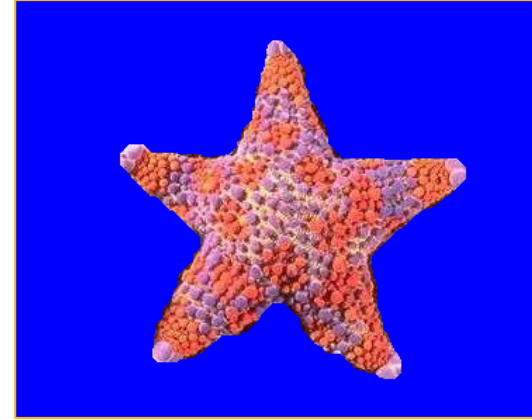


Result MRF
9-connected
(7 attractive; 2 repulsive)

Reason2: Discretization artefacts



4-connected
Euclidean



8-connected
Euclidean

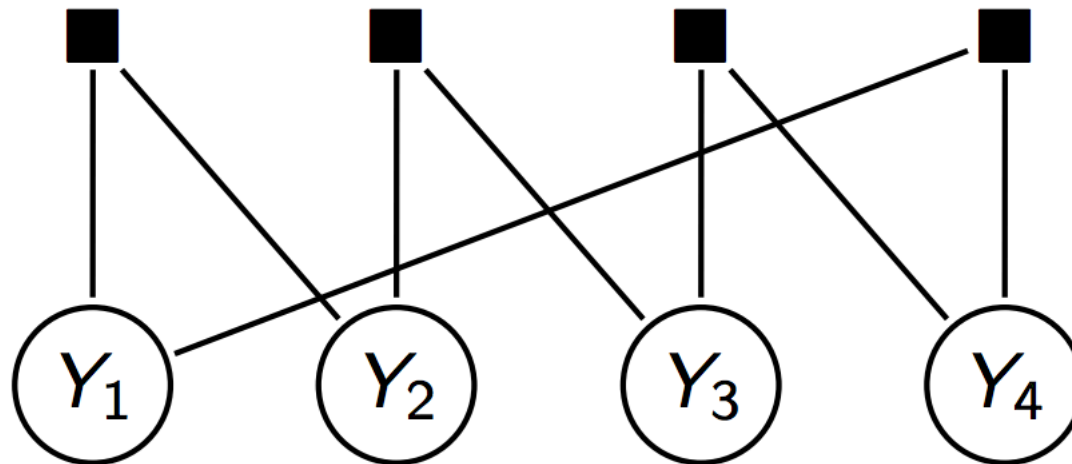
higher-connectivity can model
true Euclidean length

[Boykov et al. '03; '05]

Graphical representation

- Example

$$E(\mathbf{y}) = \psi(y_1, y_2) + \psi(y_2, y_3) + \psi(y_3, y_4) + \psi(y_4, y_1)$$

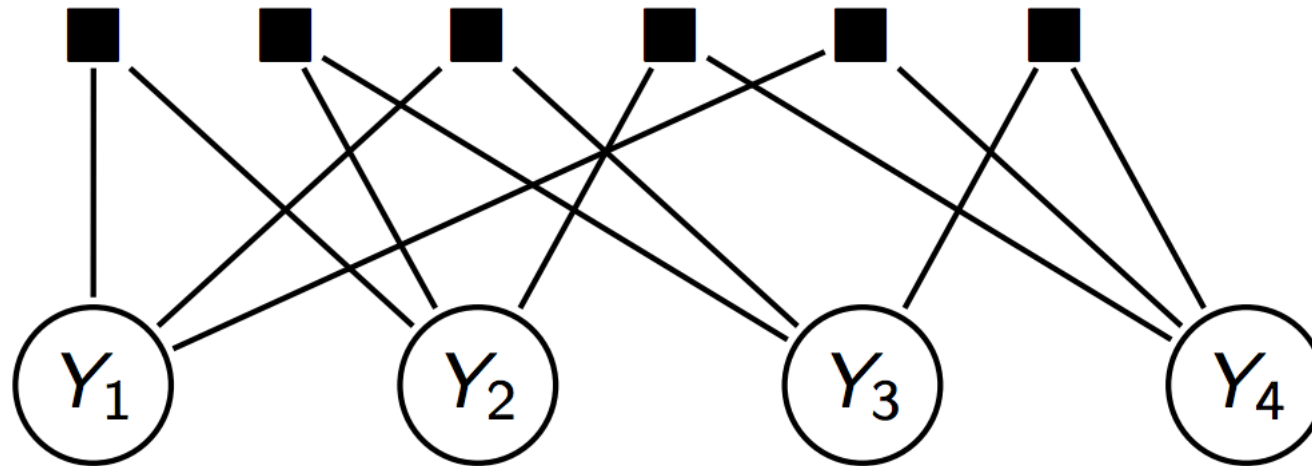


factor graph

Graphical representation

- Example

$$E(\mathbf{y}) = \sum_{i,j} \psi(y_i, y_j)$$

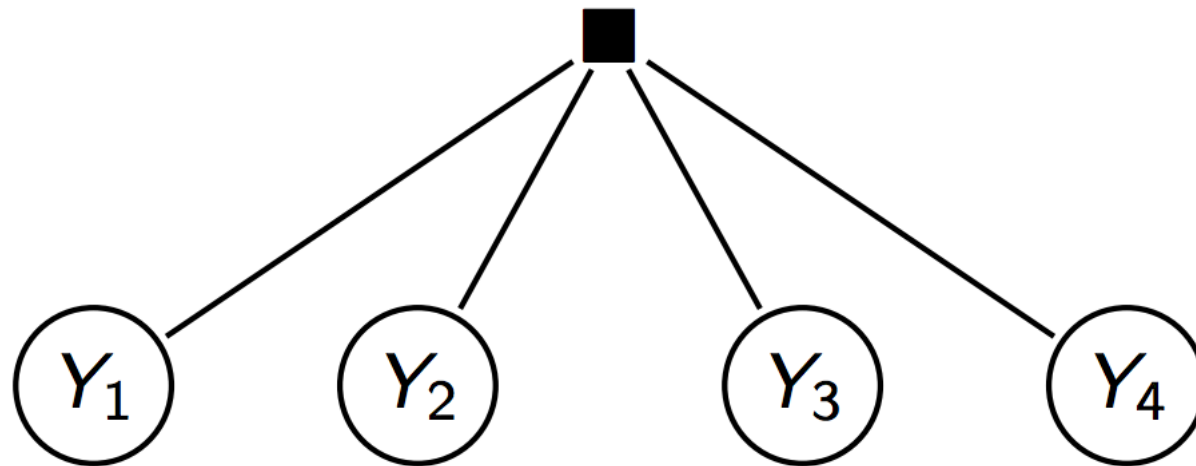


factor graph

Graphical representation

- Example

$$E(\mathbf{y}) = \psi(y_1, y_2, y_3, y_4)$$



factor graph

A Computer Vision Application

Binary Image Segmentation



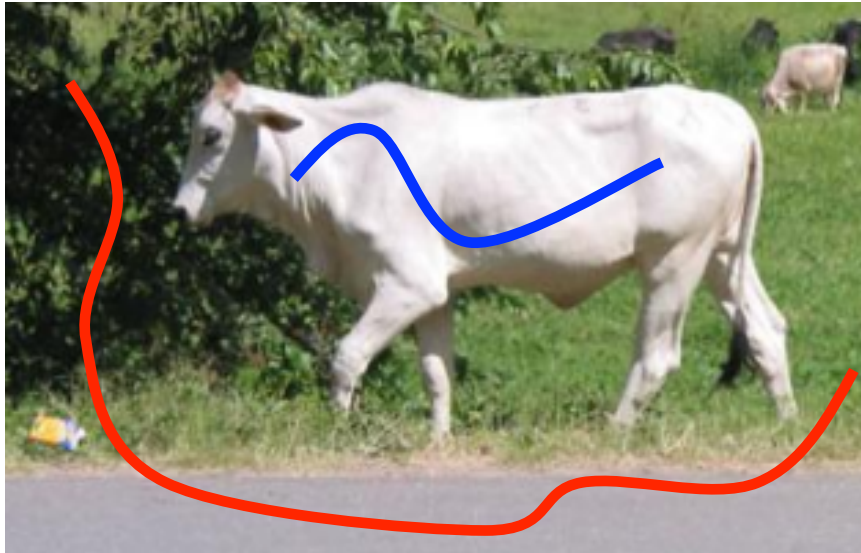
How ?

Cost function Models *our* knowledge about natural images

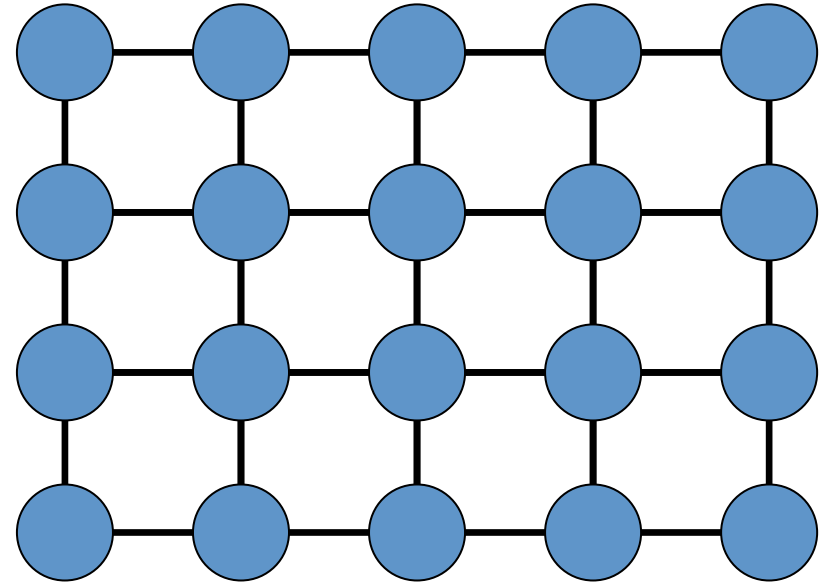
Optimize cost function to obtain the segmentation

A Computer Vision Application

Binary Image Segmentation



Object - white, Background - green/grey



Graph $G = (V, E)$

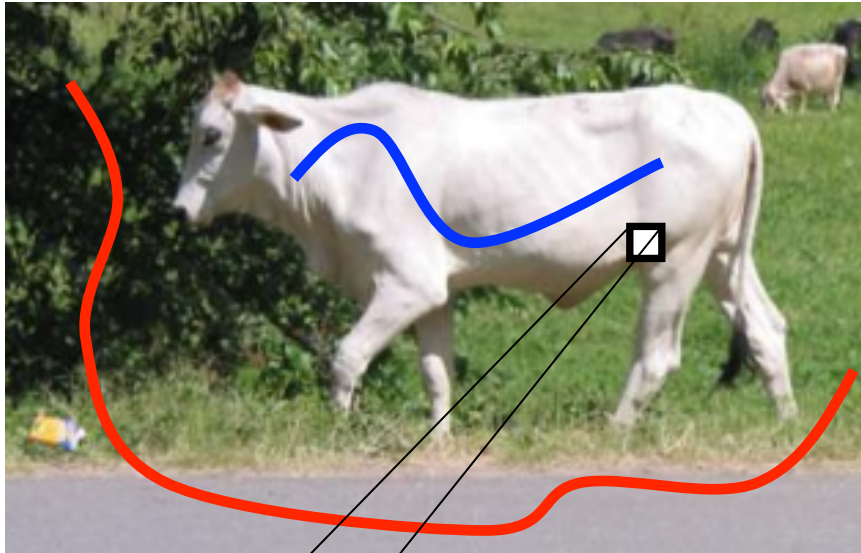
Each vertex corresponds to a pixel

Edges define a 4-neighbourhood *grid* graph

Assign a label to each vertex from $L = \{\text{obj}, \text{bkg}\}$

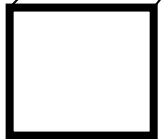
A Computer Vision Application

Binary Image Segmentation

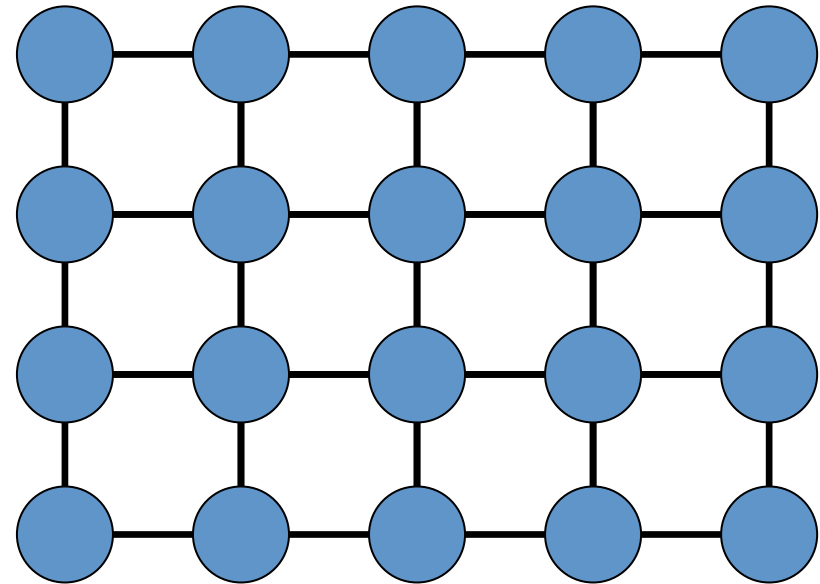


Object - white, Background - green/grey

Cost of a labelling $f : V \rightarrow L$



Cost of label 'obj' low Cost of label 'bkg' high

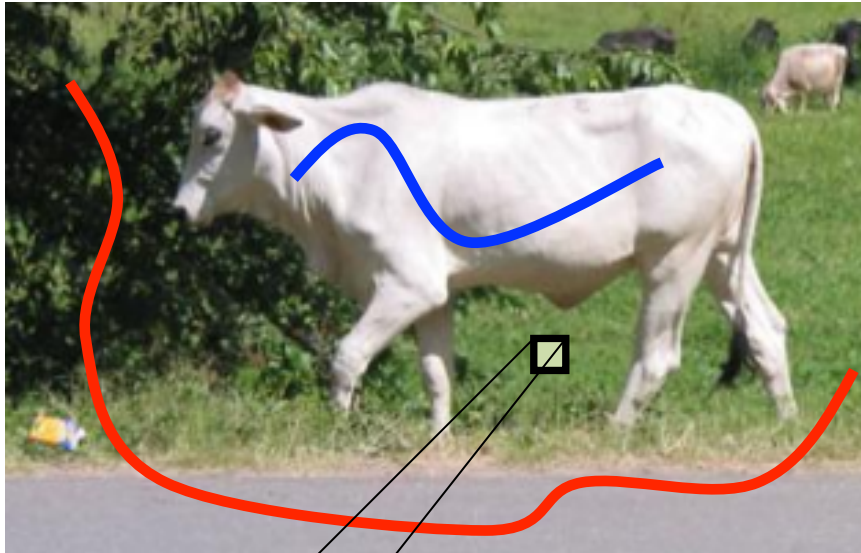


Graph $G = (V, E)$

Per Vertex Cost

A Computer Vision Application

Binary Image Segmentation

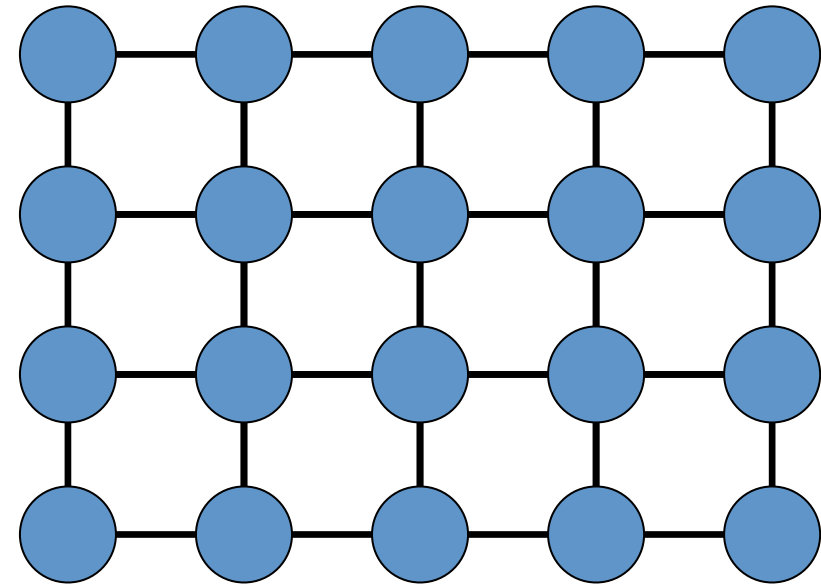


Object - white, Background - green/grey

Cost of a labelling $f : V \rightarrow L$



Cost of label 'obj' high Cost of label 'bkg' low



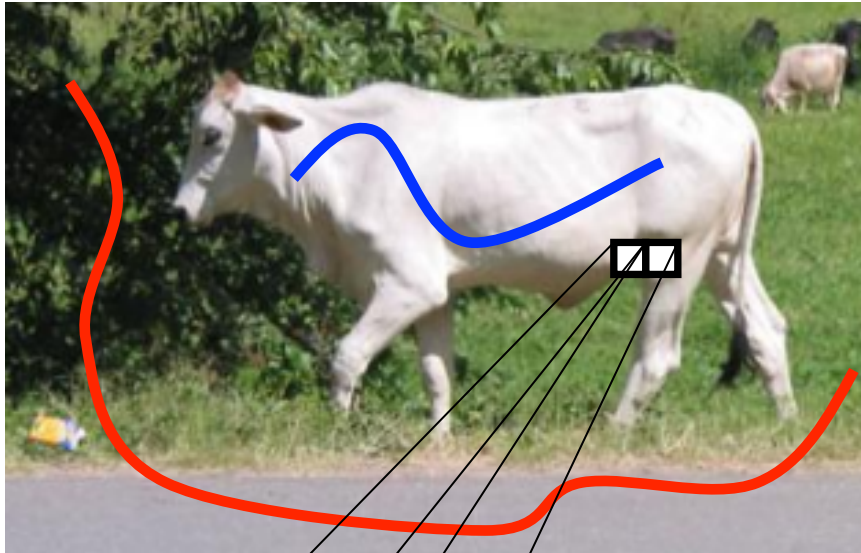
Graph $G = (V, E)$

Per Vertex Cost

UNARY COST

A Computer Vision Application

Binary Image Segmentation



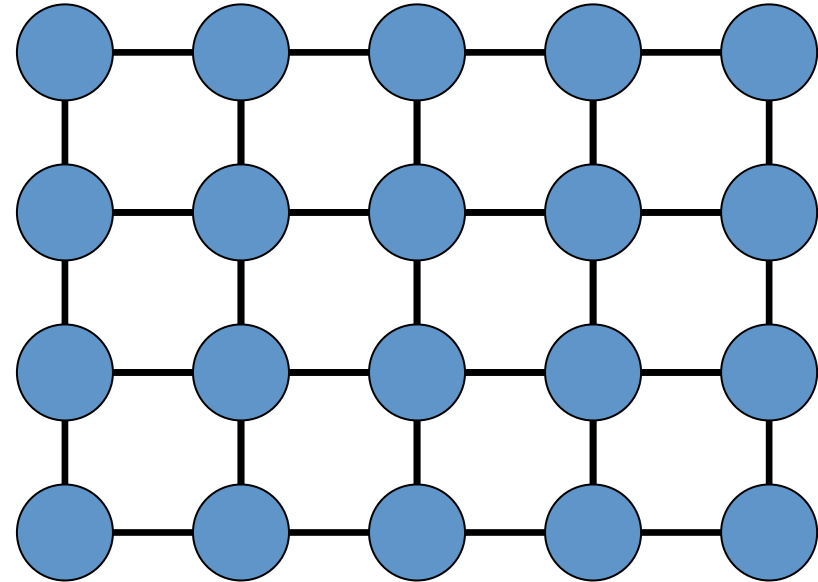
Object - white, Background - green/grey

Cost of a labelling $f : V \rightarrow L$



Cost of same label low

Cost of different labels high

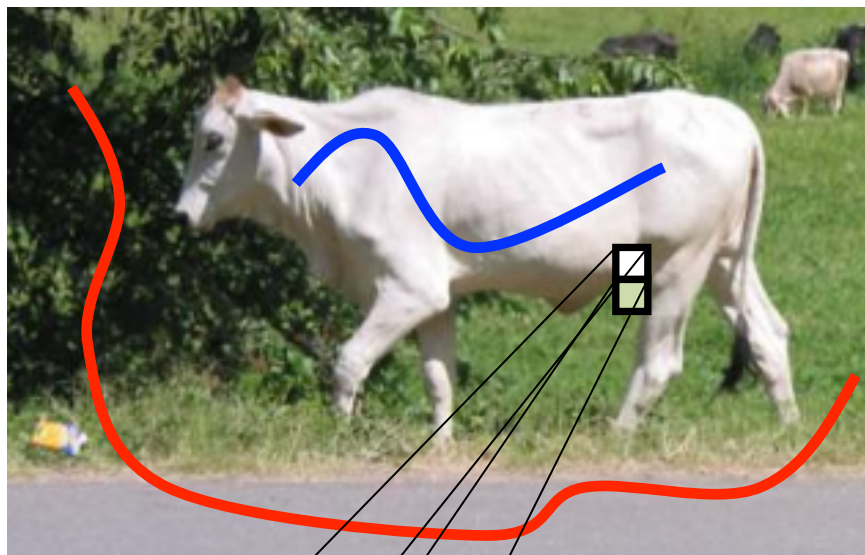


Graph $G = (V, E)$

Per Edge Cost

A Computer Vision Application

Binary Image Segmentation



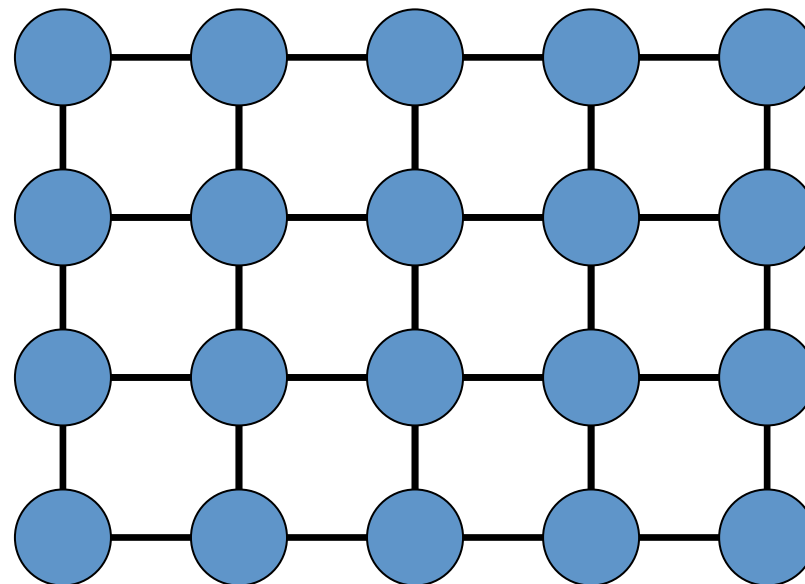
Object - white, Background - green/grey

Cost of a labelling $f : V \rightarrow L$



Cost of same label high

Cost of different labels low



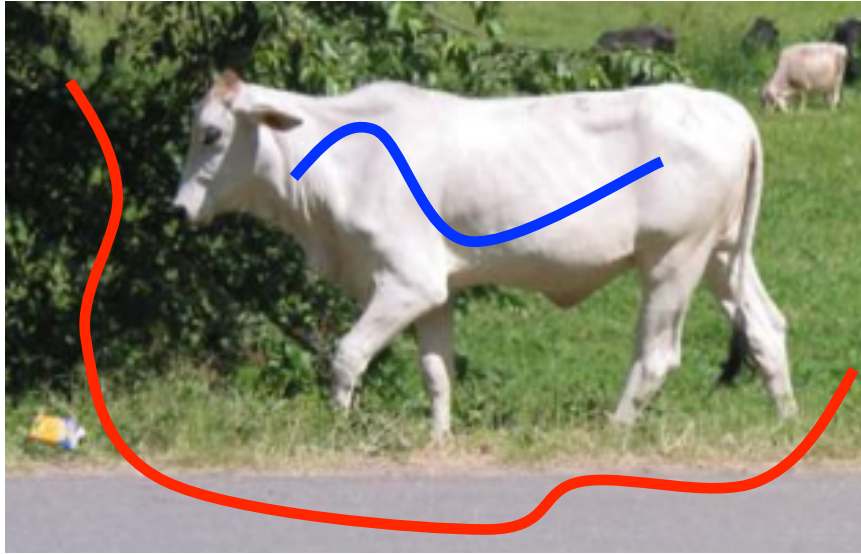
Graph $G = (V, E)$

Per Edge Cost

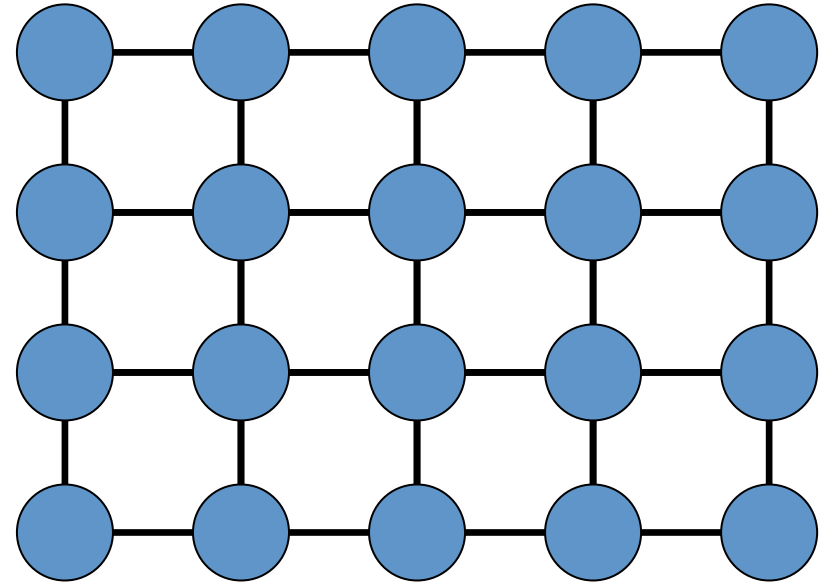
**PAIRWISE
COST**

A Computer Vision Application

Binary Image Segmentation



Object - white, Background - green/grey

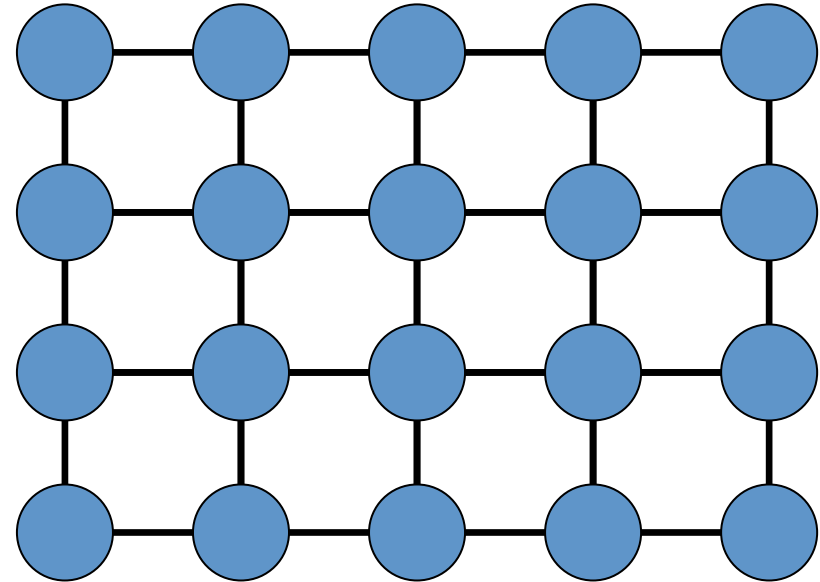


Graph $G = (V, E)$

Problem: Find the labelling with minimum cost f^*

A Computer Vision Application

Binary Image Segmentation



Graph $G = (V, E)$

Problem: Find the labelling with minimum cost f^*

Another Computer Vision Application

Stereo Correspondence



Disparity Map

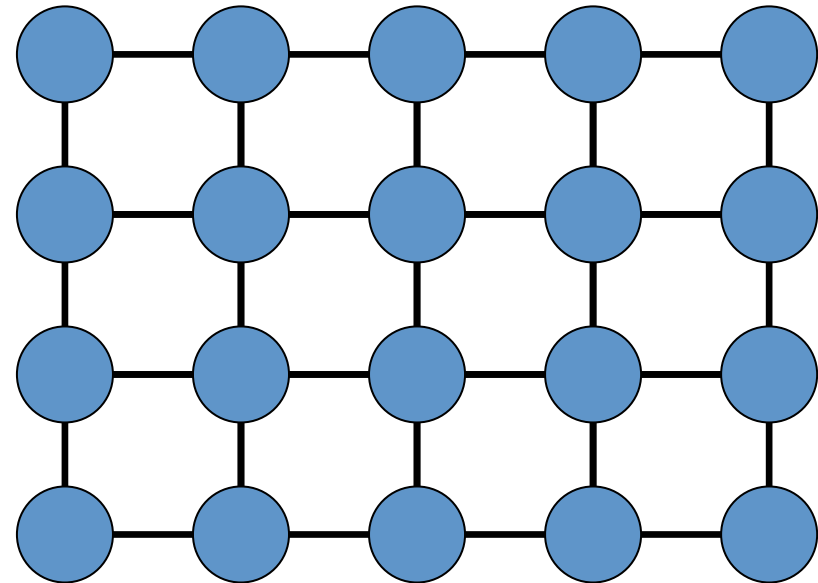


How ?

Minimizing a cost function

Another Computer Vision Application

Stereo Correspondence



Graph $G = (V, E)$

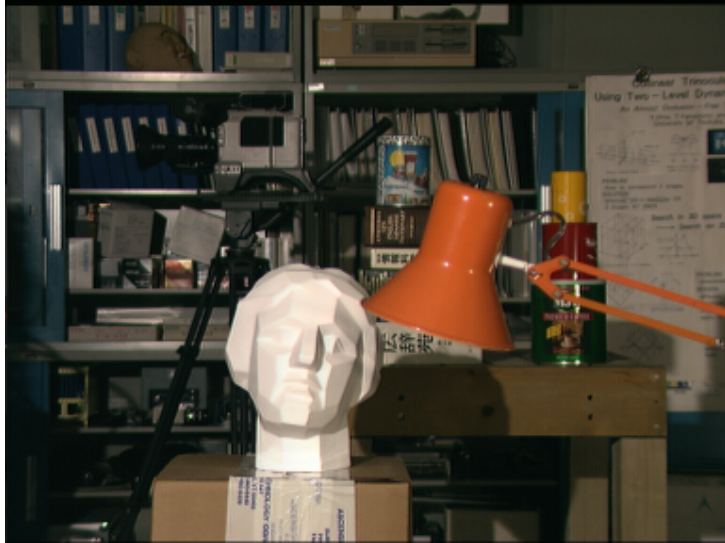
Vertex corresponds to a pixel

Edges define grid graph

$L = \{\text{disparities}\}$

Another Computer Vision Application

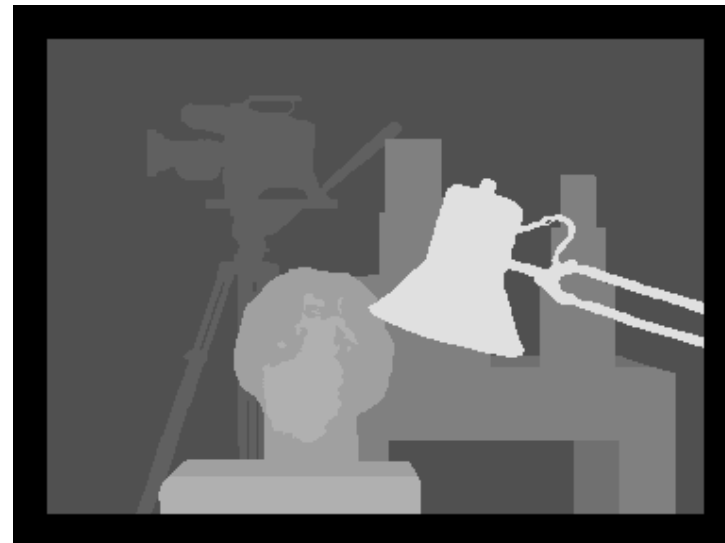
Stereo Correspondence



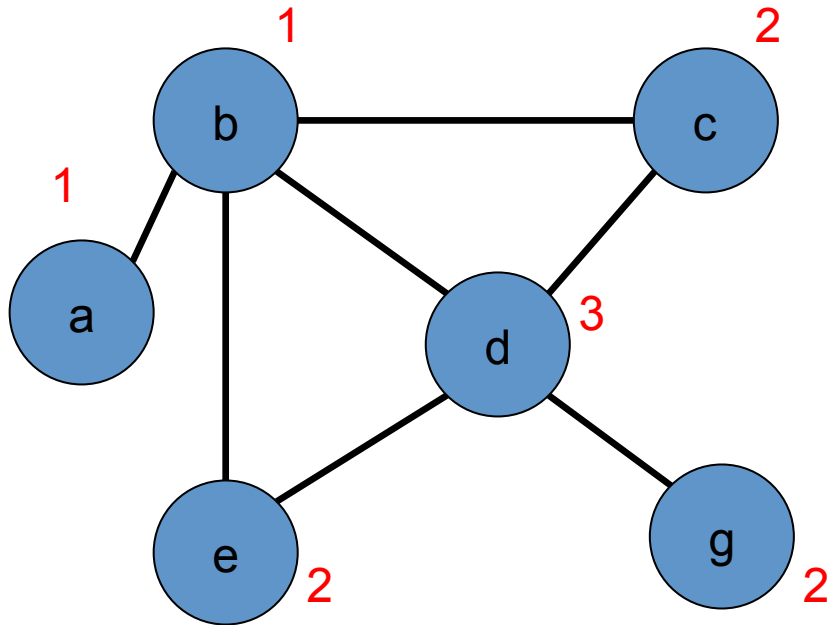
Cost of labelling f :

Unary cost + Pairwise Cost

Find minimum cost f^*



The General Problem



Graph $G = (V, E)$

Discrete label set $L = \{1, 2, \dots, h\}$

Assign a label to each vertex

$f: V \rightarrow L$

Cost of a labelling $Q(f)$

Unary Cost

Pairwise Cost

Find $f^* = \arg \min Q(f)$

Overview

- Basics: problem formulation
 - Energy Function
 - MAP Estimation
 - Computing min-marginals
 - Reparameterization
- Solutions
 - Belief Propagation and related methods
 - Graph cuts

Remainder of today's lecture

- Belief propagation
- TRW
- Graph cuts

Belief Propagation

A Computer Vision Application

Binary Image Segmentation



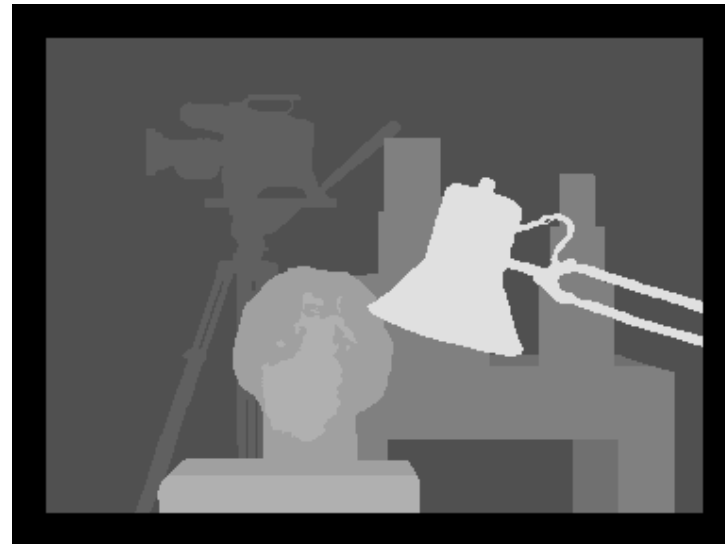
How ?

Cost function Models *our* knowledge about natural images

Optimize cost function to obtain the segmentation

Another Computer Vision Application

Stereo Correspondence

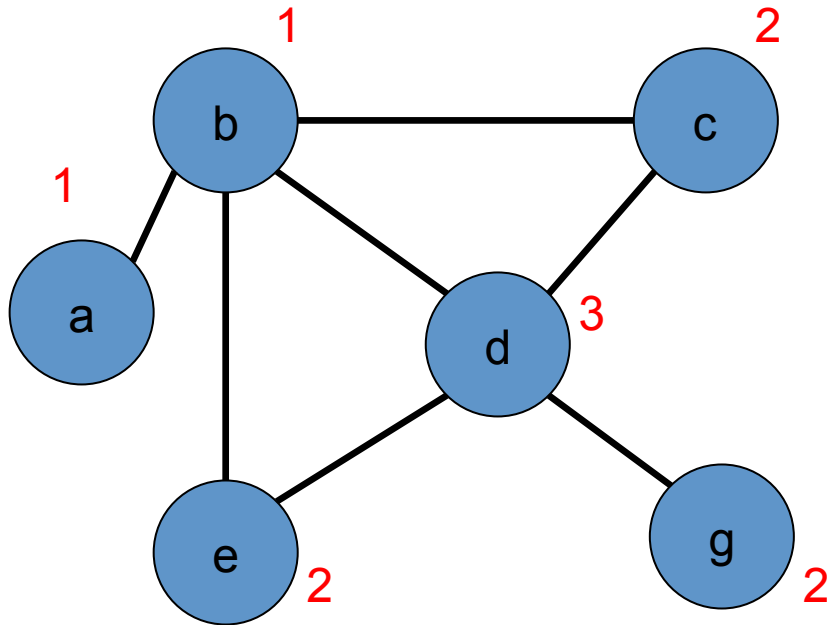


Disparity Map

How ?

Minimizing a cost function

The General Problem



Graph $G = (V, E)$

Discrete label set $L = \{1, 2, \dots, h\}$

Assign a label to each vertex

$f: V \rightarrow L$

Cost of a labelling $Q(f)$

Unary Cost

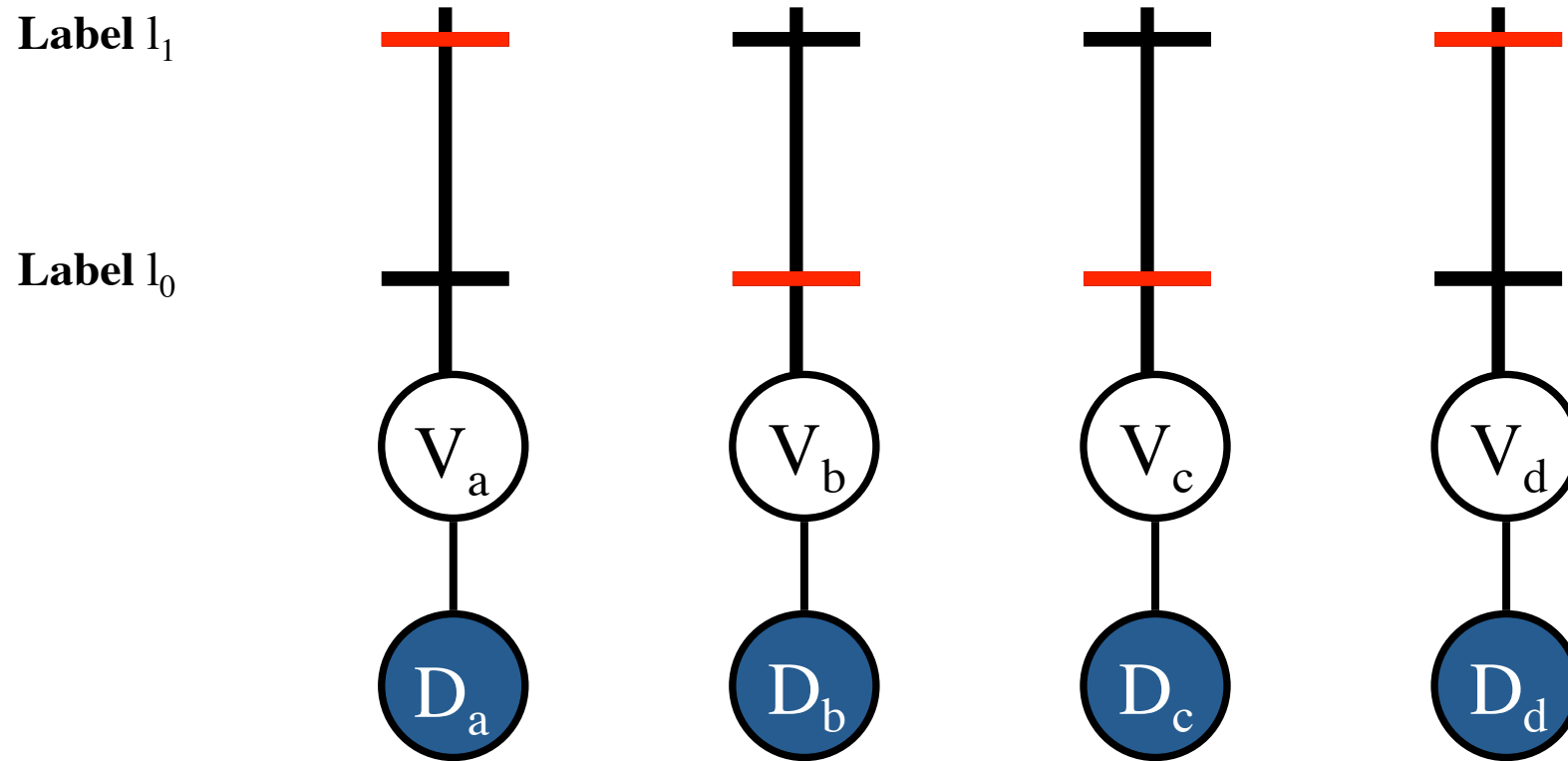
Pairwise Cost

Find $f^* = \arg \min Q(f)$

Overview

- Basics: problem formulation
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Energy Function

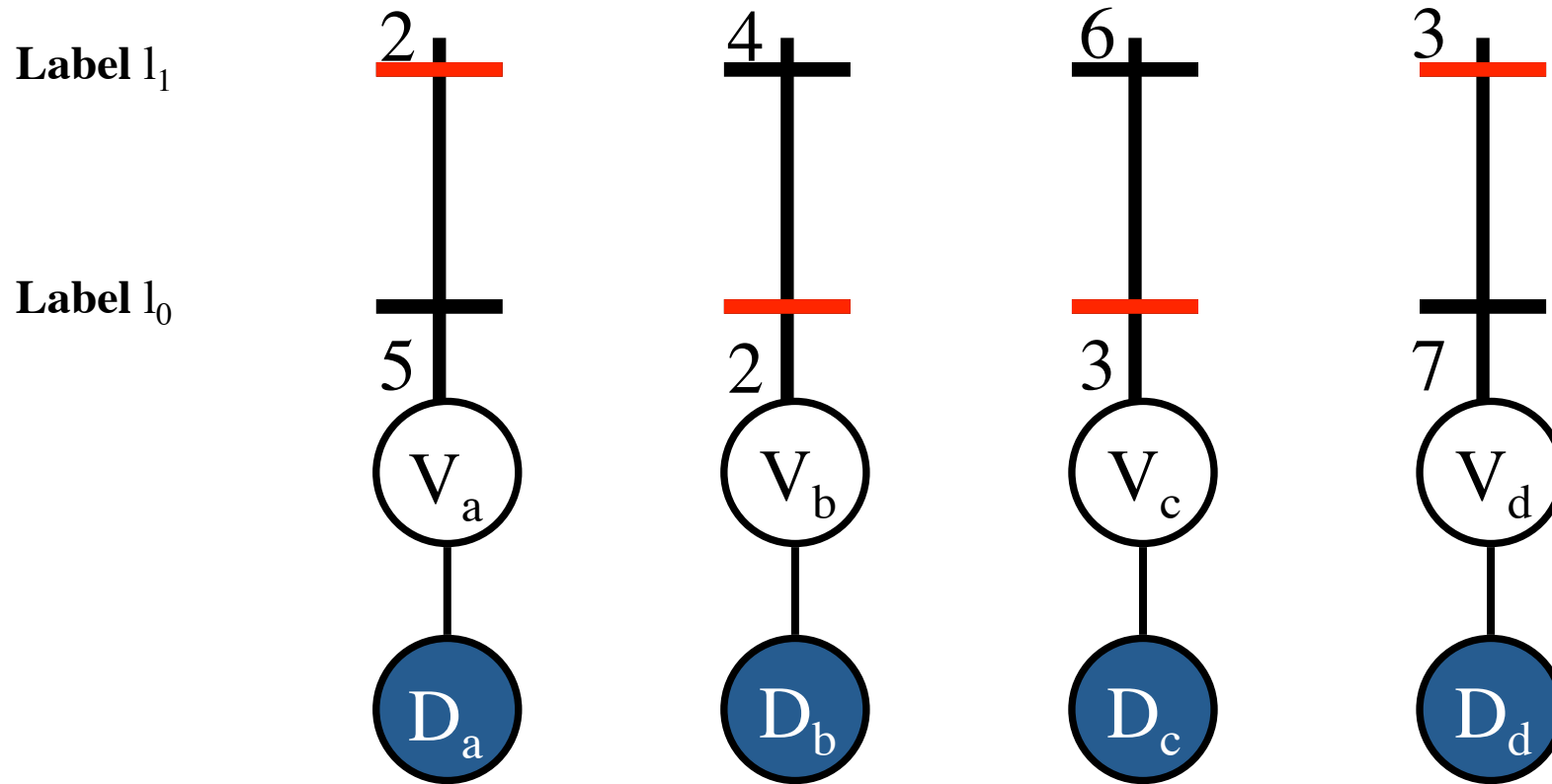


Random Variables $V = \{V_a, V_b, \dots\}$

Labels $L = \{l_0, l_1, \dots\}$ Data D

Labelling $f: \{a, b, \dots\} \rightarrow \{0, 1, \dots\}$

Energy Function



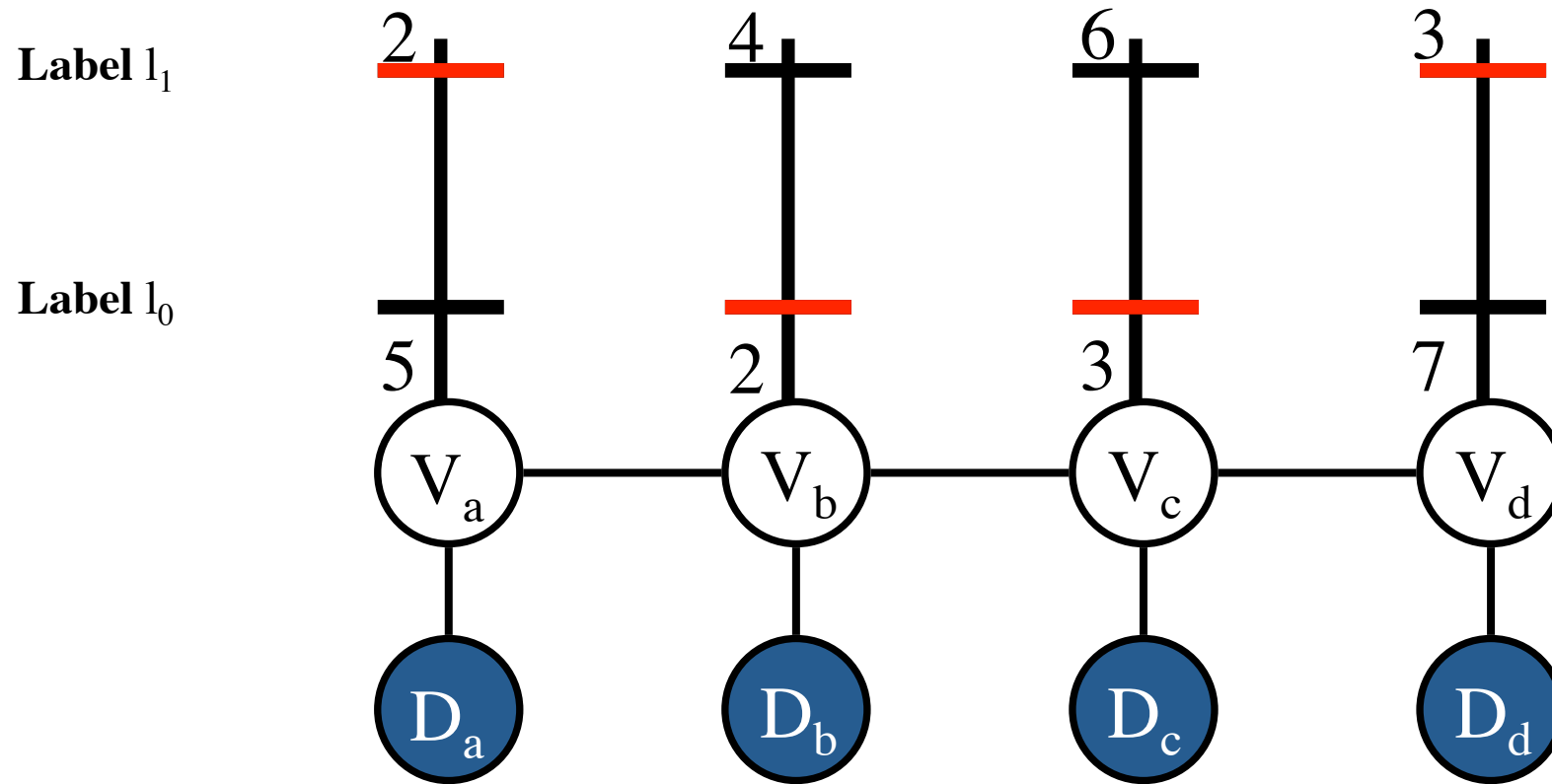
$$Q(f) = \sum_a \theta_{a;f(a)}$$

Unary Potential

Easy to minimize

Neighbourhood

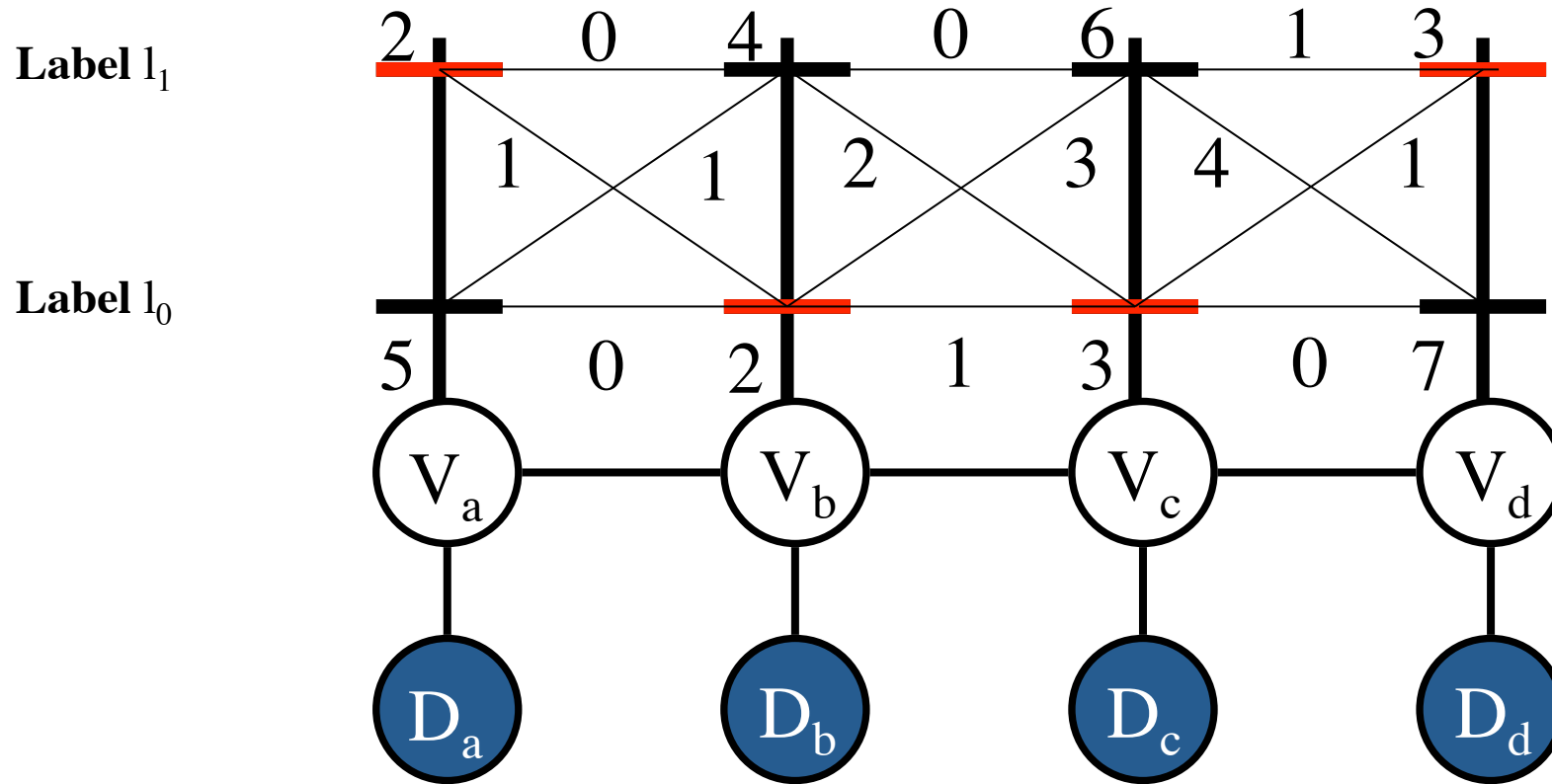
Energy Function



$E : (a,b) \in E$ iff V_a and V_b are neighbours

$$E = \{ (a,b) , (b,c) , (c,d) \}$$

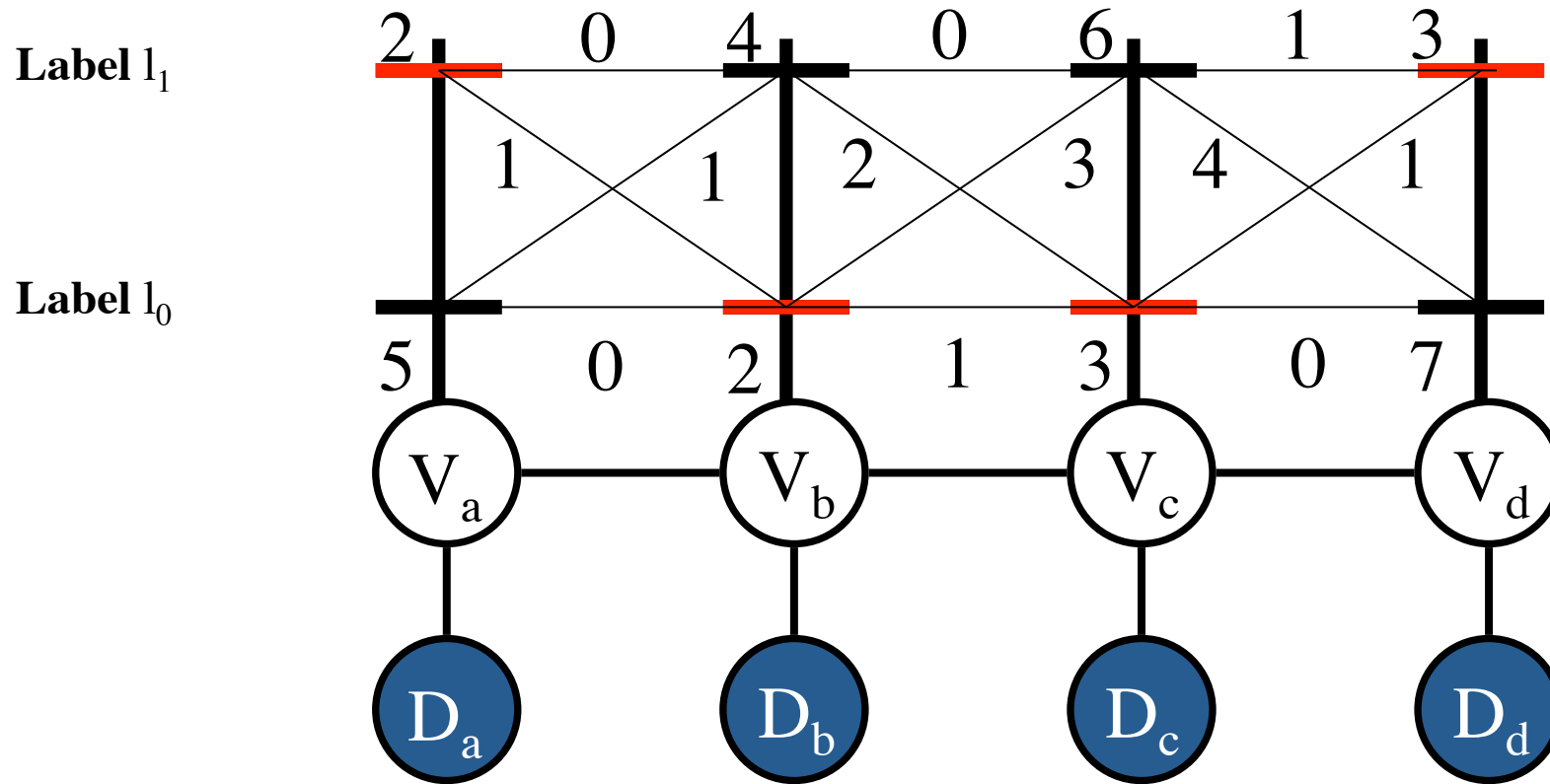
Energy Function



Pairwise Potential

$$Q(f) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

Energy Function



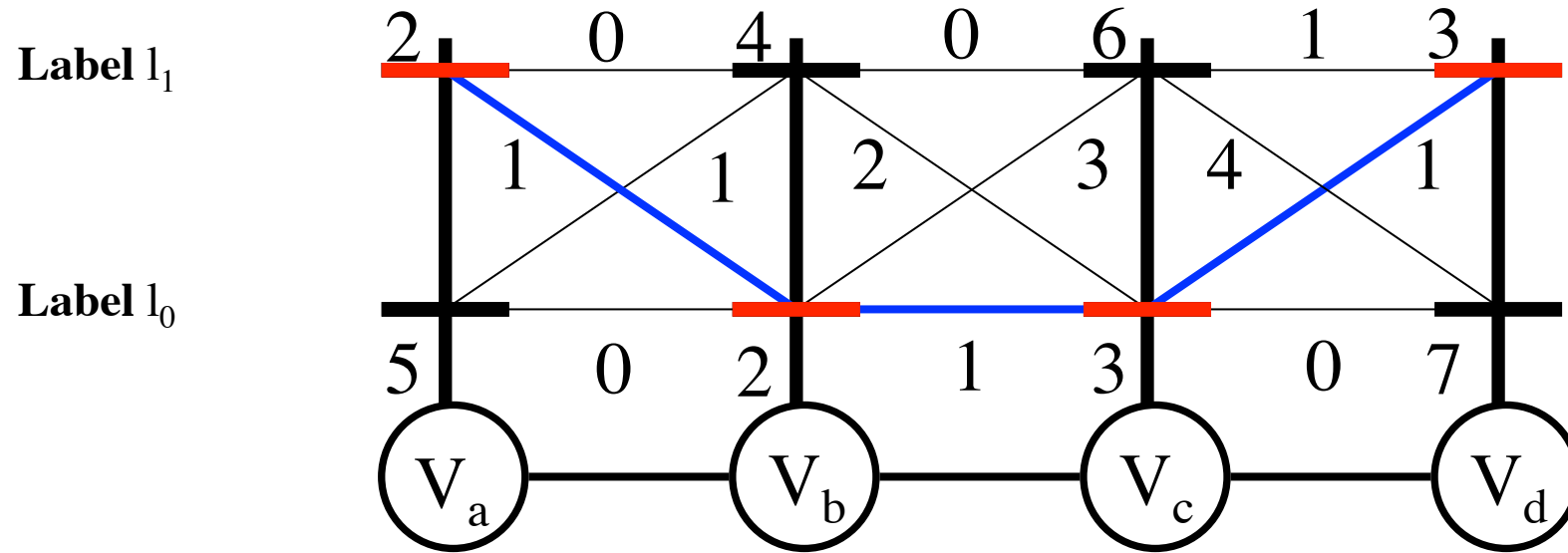
$$Q(f; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

Parameter

Overview

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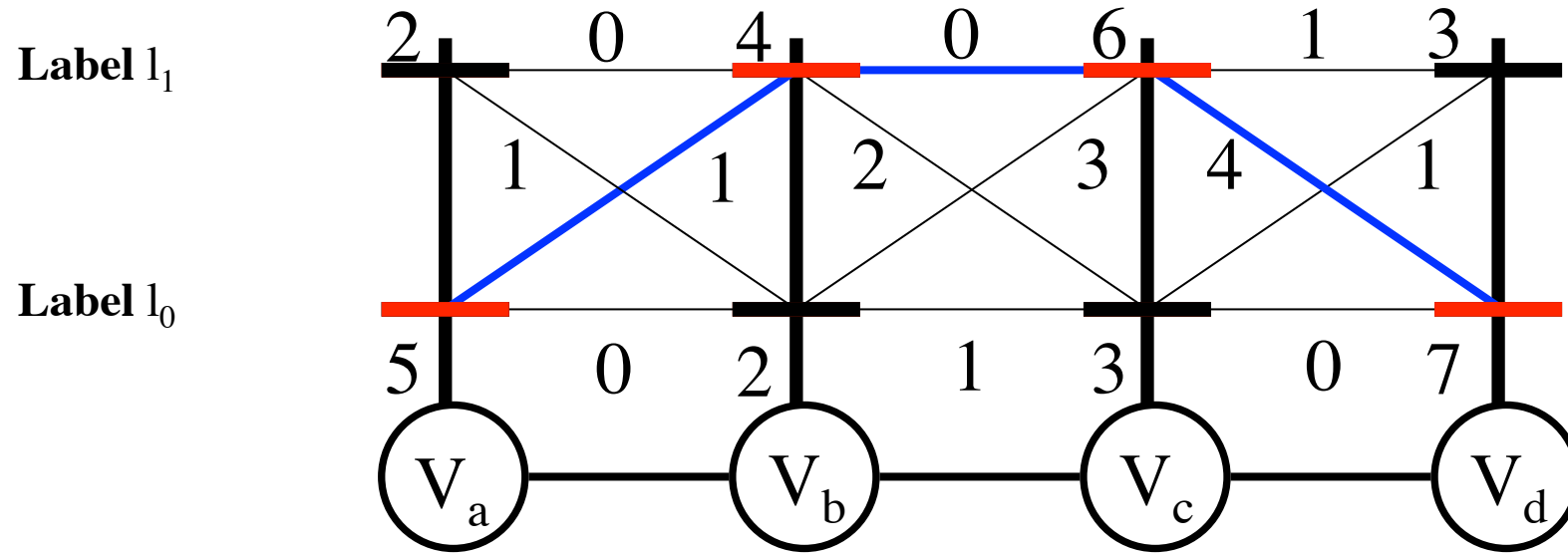
MAP Estimation



$$Q(f; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

$$2 + 1 + 2 + 1 + 3 + 1 + 3 = 13$$

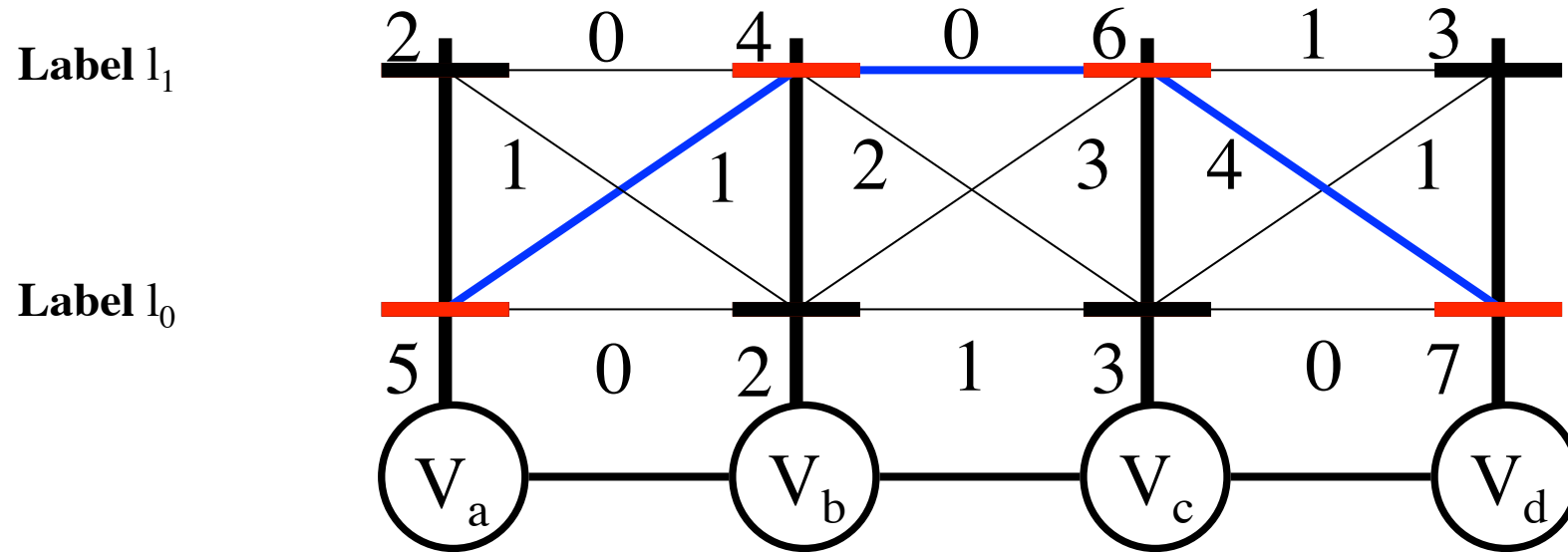
MAP Estimation



$$Q(f; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

$$5 + 1 + 4 + 0 + 6 + 4 + 7 = 27$$

MAP Estimation



$$q^* = \min Q(\mathbf{f}; \theta) = Q(\mathbf{f}^*; \theta)$$

$$Q(\mathbf{f}; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

$$\mathbf{f}^* = \arg \min Q(\mathbf{f}; \theta)$$

Equivalent to maximizing the associated probability

MAP Estimation

16 possible labellings

$$f^* = \{1, 0, 0, 1\}$$

$$q^* = 13$$

f(a)	f(b)	f(c)	f(d)	Q(f; θ)
0	0	0	0	18
0	0	0	1	15
0	0	1	0	27
0	0	1	1	20
0	1	0	0	22
0	1	0	1	19
0	1	1	0	27
0	1	1	1	20

f(a)	f(b)	f(c)	f(d)	Q(f; θ)
1	0	0	0	16
1	0	0	1	13
1	0	1	0	25
1	0	1	1	18
1	1	0	0	18
1	1	0	1	15
1	1	1	0	23
1	1	1	1	16

Computational Complexity

Segmentation

$2^{|V|}$



$|V|$ = number of pixels ≈ 153600

Can we do better than brute-force?

MAP Estimation is NP-hard !!

MAP Inference / Energy Minimization

- Computing the assignment minimizing the energy in NP-hard in general

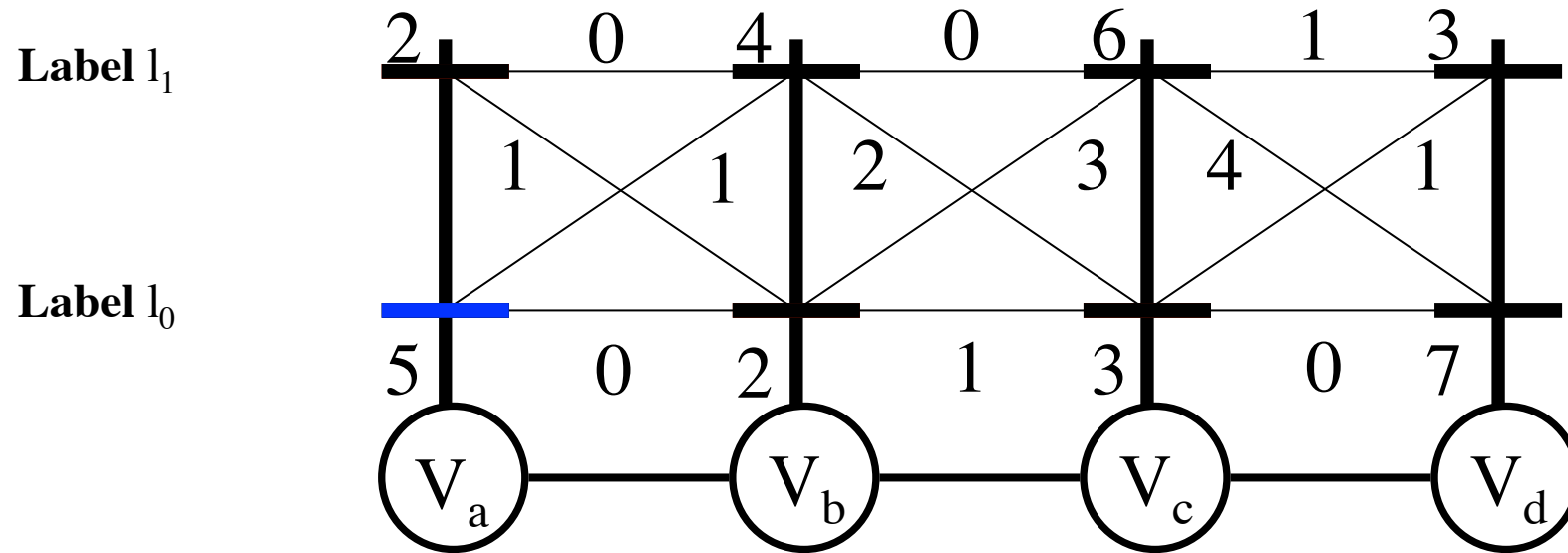
$$\operatorname{argmin}_{\mathbf{y} \in \mathcal{Y}} E(\mathbf{y}; \mathbf{x}) = \operatorname{argmax}_{\mathbf{y} \in \mathcal{Y}} P(\mathbf{y} | \mathbf{x})$$

- Exact inference is possible in some cases, e.g.,
 - Low treewidth graphs \rightarrow message-passing
 - Submodular potentials \rightarrow graph cuts
- Efficient approximate inference algorithms exist
 - Message passing on general graphs
 - Move-making algorithms
 - Relaxation algorithms

Overview

- Basics: problem formulation
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Min-Marginals



Not a marginal (no summation)

$$f^* = \arg \min Q(f; \theta) \quad \text{such that } f(a) = i$$

Min-marginal $q_{a,i}$

Min-Marginals

16 possible labellings

$$Q_{a;0} = 15$$

f(a)	f(b)	f(c)	f(d)	Q(f; θ)
0	0	0	0	18
0	0	0	1	15
0	0	1	0	27
0	0	1	1	20
0	1	0	0	22
0	1	0	1	19
0	1	1	0	27
0	1	1	1	20

f(a)	f(b)	f(c)	f(d)	Q(f; θ)
1	0	0	0	16
1	0	0	1	13
1	0	1	0	25
1	0	1	1	18
1	1	0	0	18
1	1	0	1	15
1	1	1	0	23
1	1	1	1	16

Min-Marginals

16 possible labellings

$$Q_{a;1} = 13$$

f(a)	f(b)	f(c)	f(d)	Q(f; θ)
0	0	0	0	18
0	0	0	1	15
0	0	1	0	27
0	0	1	1	20
0	1	0	0	22
0	1	0	1	19
0	1	1	0	27
0	1	1	1	20

f(a)	f(b)	f(c)	f(d)	Q(f; θ)
1	0	0	0	16
1	0	0	1	13
1	0	1	0	25
1	0	1	1	18
1	1	0	0	18
1	1	0	1	15
1	1	1	0	23
1	1	1	1	16

Min-Marginals and MAP

- Minimum min-marginal of any variable = energy of MAP labelling

$$\min_i q_{a;i}$$

$$\min_i (\min_f Q(f; \theta) \text{ such that } f(a) = i)$$

V_a has to take one label

$$\min_f Q(f; \theta)$$

Summary

Energy Function

$$Q(\mathbf{f}; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

MAP Estimation

$$\mathbf{f}^* = \arg \min Q(\mathbf{f}; \theta)$$

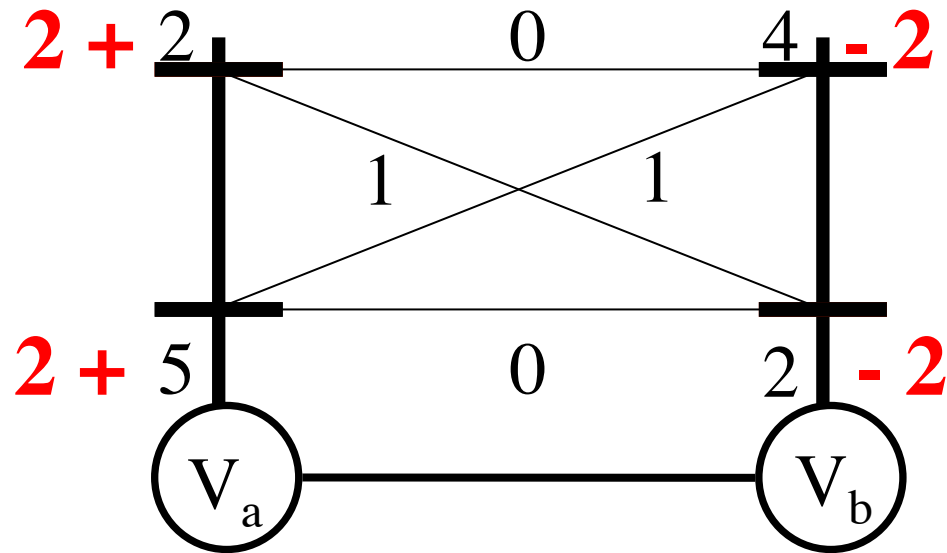
Min-marginals

$$q_{a;i} = \min Q(\mathbf{f}; \theta) \quad \text{s.t. } f(a) = i$$

Overview

- Basics: problem formulation
 - Energy Function
 - MAP Estimation
 - Computing min-marginals
 - Reparameterization
- Solutions
 - Belief Propagation and related methods
 - Graph cuts

Reparameterization



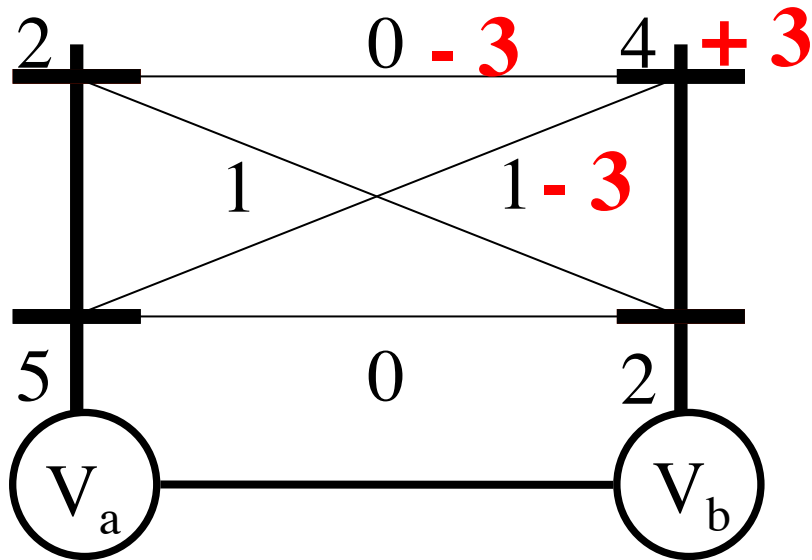
f(a)	f(b)	Q(f; θ)
0	0	7 + 2 - 2
0	1	10 + 2 - 2
1	0	5 + 2 - 2
1	1	6 + 2 - 2

Add a constant to all $\theta_{a;i}$

Subtract that constant from all $\theta_{b;k}$

$$Q(f; \theta') = Q(f; \theta)$$

Reparameterization



f(a)	f(b)	Q(f; θ)
0	0	7
0	1	10 - 3 + 3
1	0	5
1	1	6 - 3 + 3

Add a constant to one $\theta_{b;k}$

Subtract that constant from $\theta_{ab;ik}$ for all 'i'

$$Q(f; \theta') = Q(f; \theta)$$

Reparameterization

θ' is a reparameterization of θ , iff

$$Q(f; \theta') = Q(f; \theta), \text{ for all } f \quad \theta' \equiv \theta$$

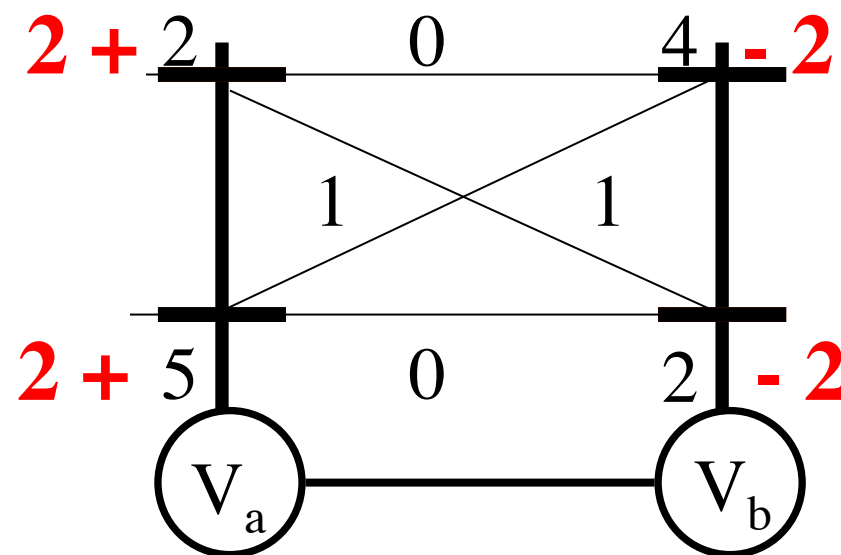
Equivalently

$$\theta'_{a;i} = \theta_{a;i} + M_{ba;i}$$

$$\theta'_{b;k} = \theta_{b;k} + M_{ab;k}$$

$$\theta'_{ab;ik} = \theta_{ab;ik} - M_{ab;k} - M_{ba;i}$$

Kolmogorov, PAMI, 2006



Recap

MAP Estimation

$$\mathbf{f}^* = \arg \min \mathbf{Q}(\mathbf{f}; \theta)$$

$$Q(\mathbf{f}; \theta) = \sum_a \theta_{a;f(a)} + \sum_{(a,b)} \theta_{ab;f(a)f(b)}$$

Min-marginals

$$q_{a;i} = \min \mathbf{Q}(\mathbf{f}; \theta) \quad \text{s.t. } f(a) = i$$

Reparameterization

$$Q(\mathbf{f}; \theta') = Q(\mathbf{f}; \theta), \text{ for all } \mathbf{f} \quad \theta' \equiv \theta$$

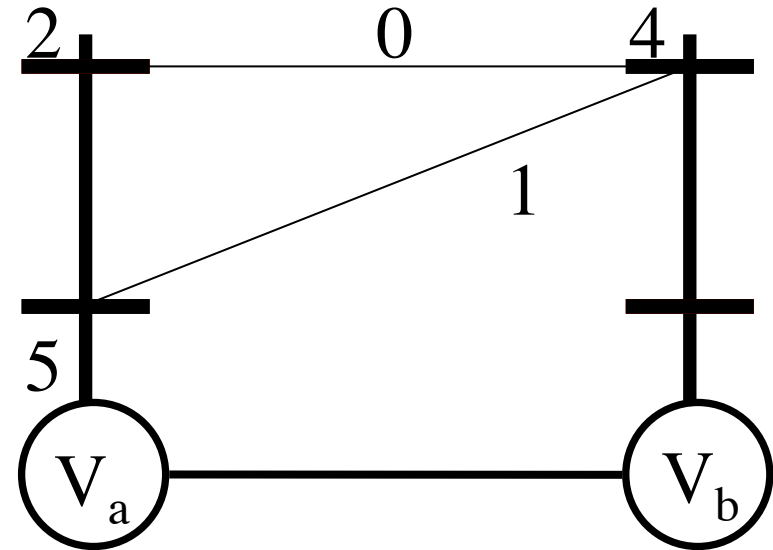
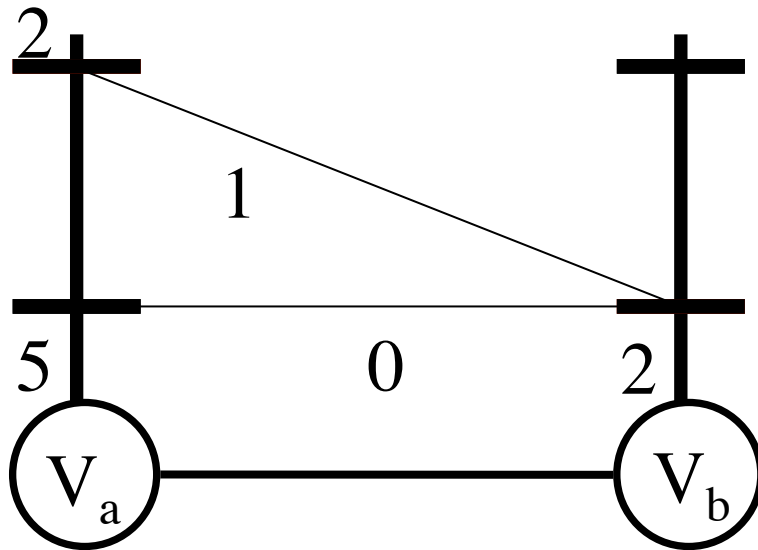
Overview

- Basics: problem formulation
 - Energy Function
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 - Reparameterization
- Solutions
 - Belief Propagation and related methods
 - Graph cuts

Belief Propagation

- Remember, some MAP problems are easy
- Belief Propagation gives exact MAP for chains
- Exact MAP for trees
- Clever Reparameterization

Two Variables

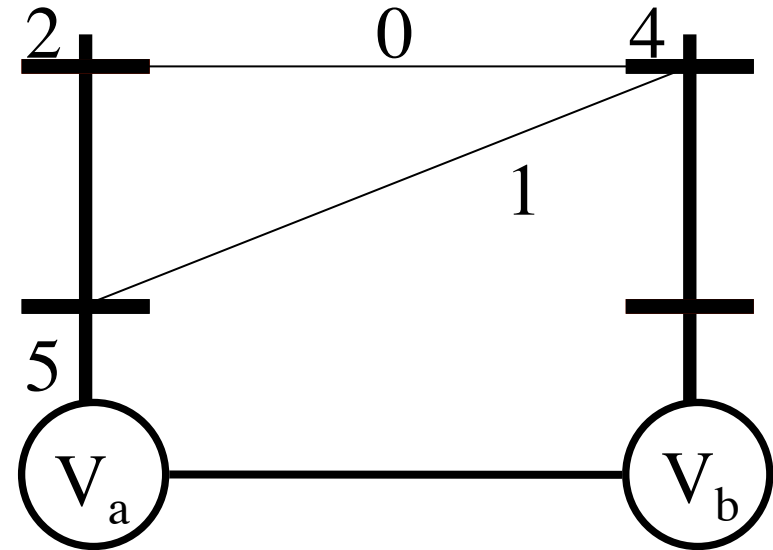
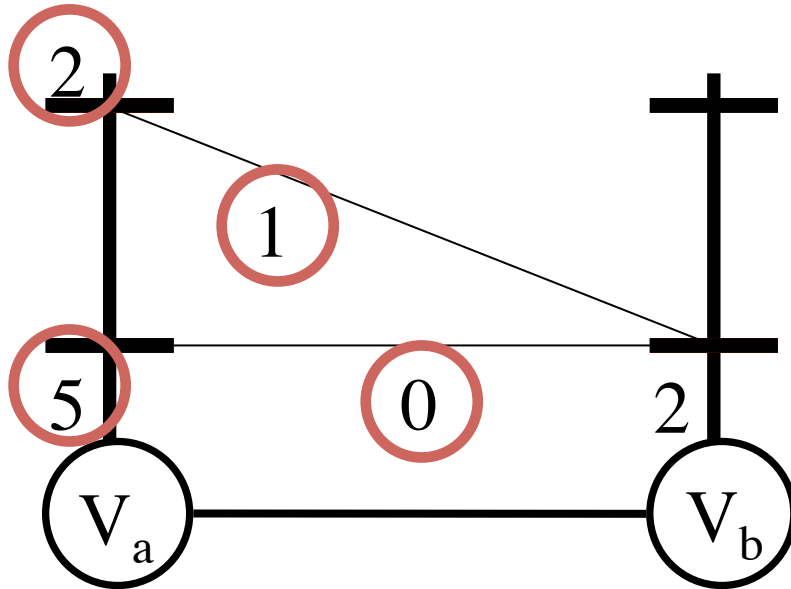


Add a constant to one $\theta_{b;k}$

Subtract that constant from $\theta_{ab;ik}$ for all 'i'

Choose the *right* constant $\theta'_{b;k} = q_{b;k}$

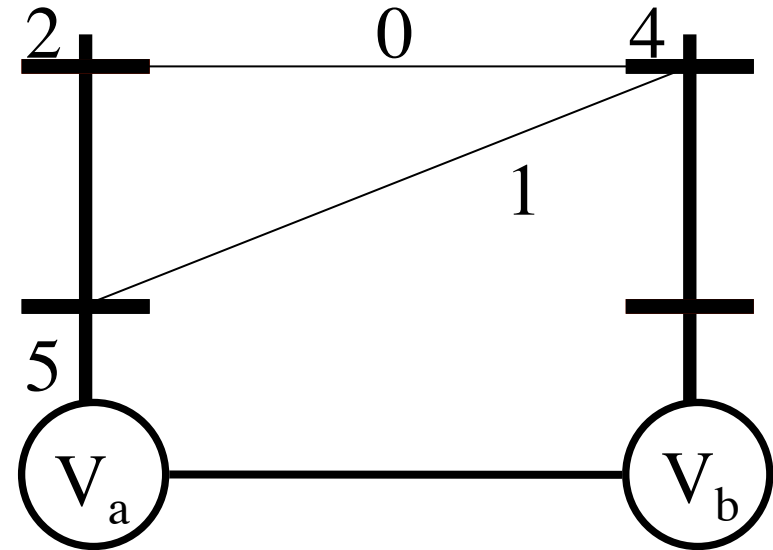
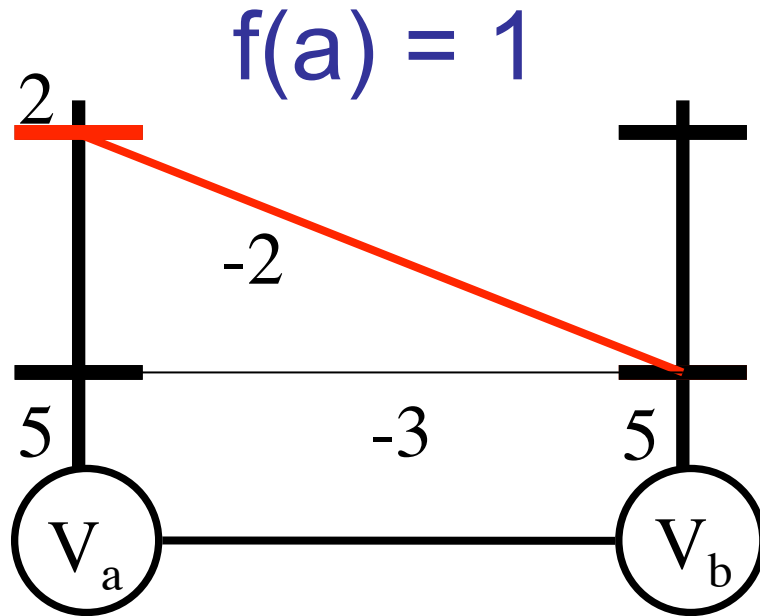
Two Variables



$$M_{ab;0} = \min \begin{cases} \theta_{a;0} + \theta_{ab;00} = 5 + 0 \\ \theta_{a;1} + \theta_{ab;10} = 2 + 1 \end{cases}$$

Choose the *right* constant $\theta'_{b;k} = q_{b;k}$

Two Variables

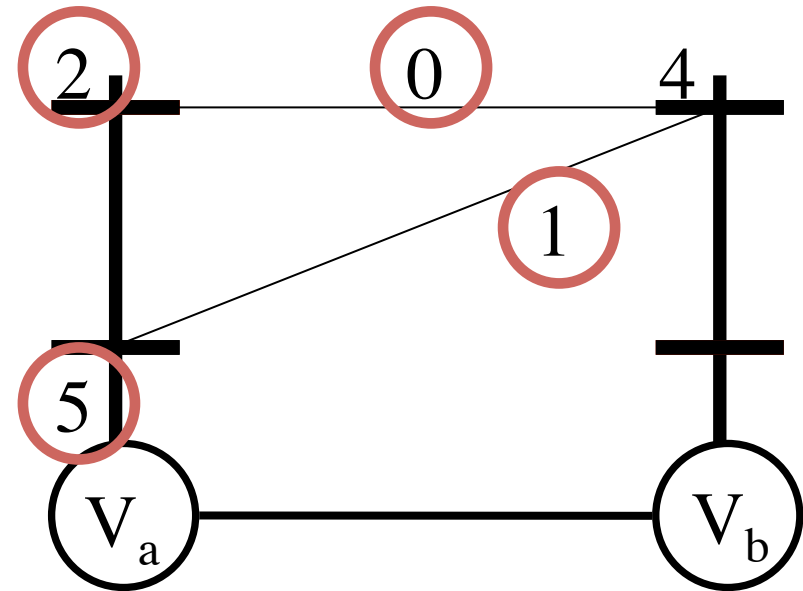
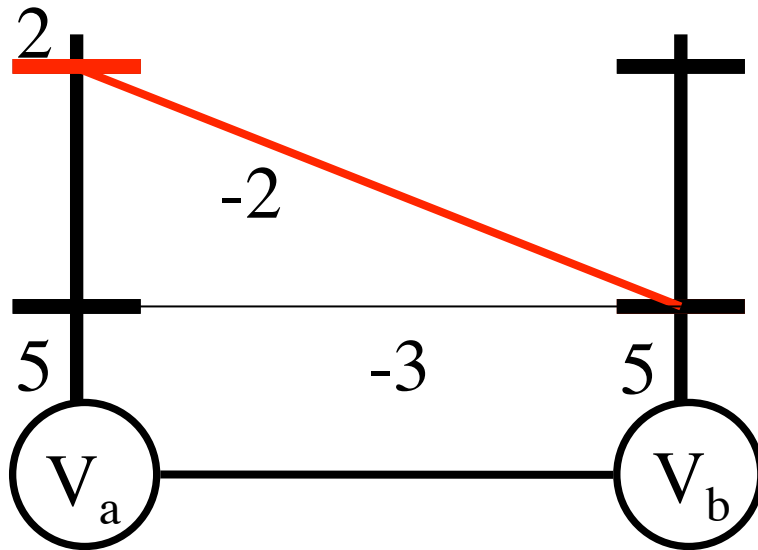


$$\theta'_{b;0} = q_{b;0}$$

Potentials along the red path add up to 0

Choose the **right** constant $\theta'_{b;k} = q_{b;k}$

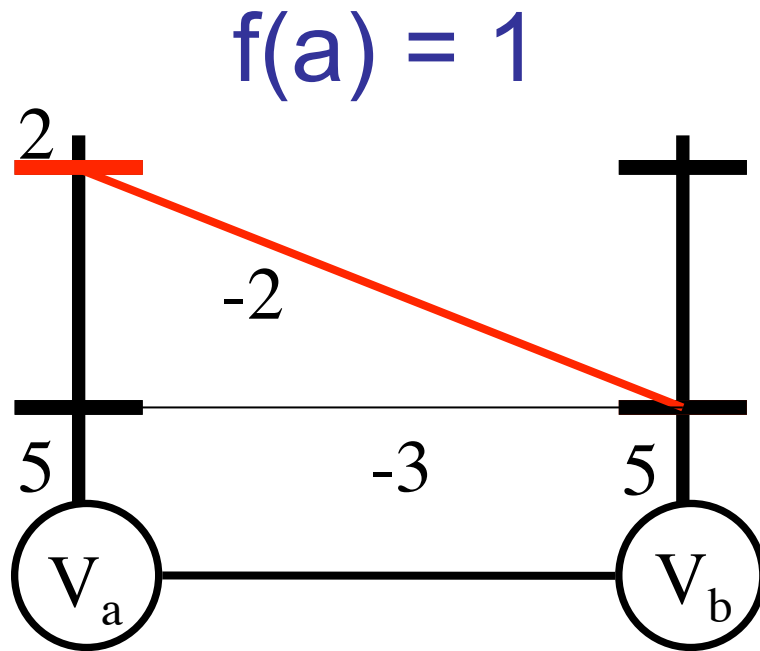
Two Variables



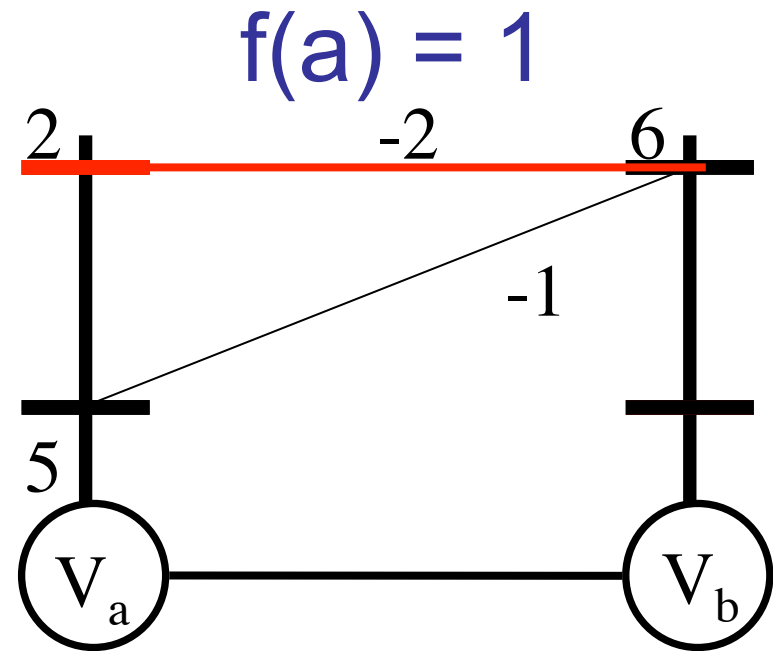
$$M_{ab;1} = \min \begin{cases} \theta_{a;0} + \theta_{ab;01} = 5 + 1 \\ \theta_{a;1} + \theta_{ab;11} = 2 + 0 \end{cases}$$

Choose the *right* constant $\theta'_{b;k} = q_{b;k}$

Two Variables



$$\theta'_{b;0} = q_{b;0}$$

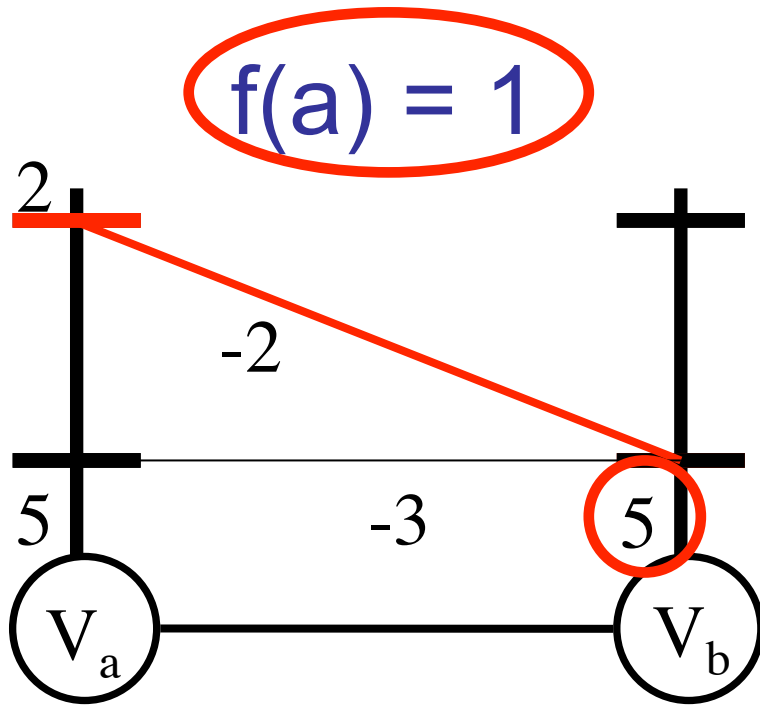


$$\theta'_{b;1} = q_{b;1}$$

Minimum of min-marginals = MAP estimate

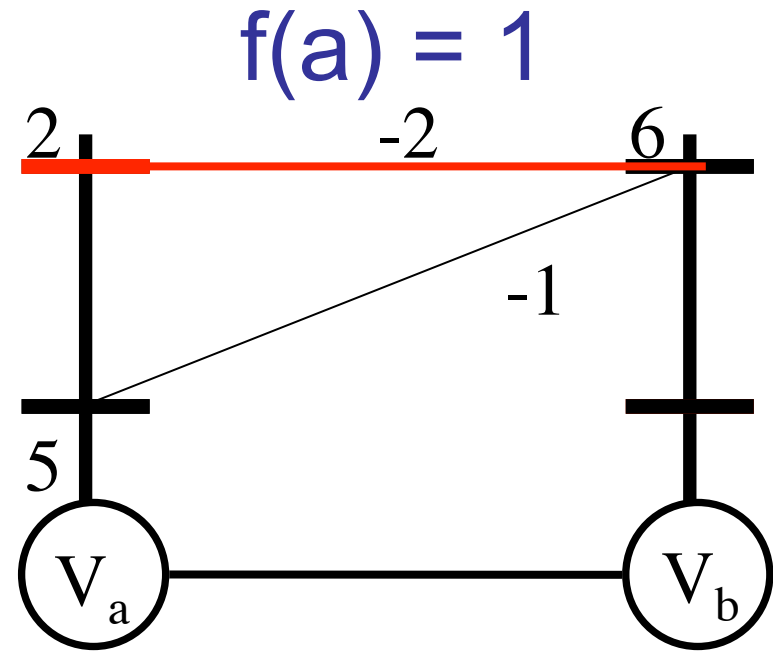
Choose the **right** constant $\theta'_{b;k} = q_{b;k}$

Two Variables



$$\theta'_{b;0} = q_{b;0}$$

$$f^*(b) = 0 \quad f^*(a) = 1$$

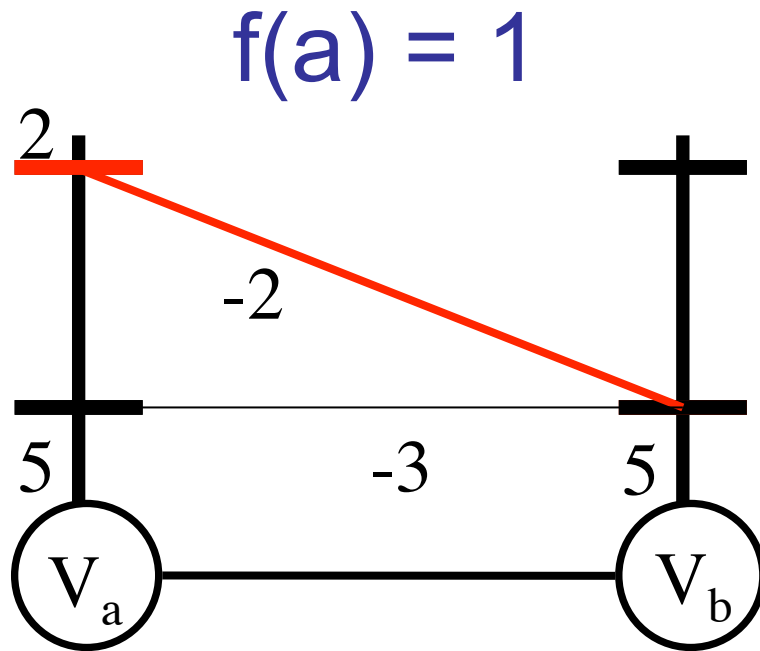


$$\theta'_{b;1} = q_{b;1}$$

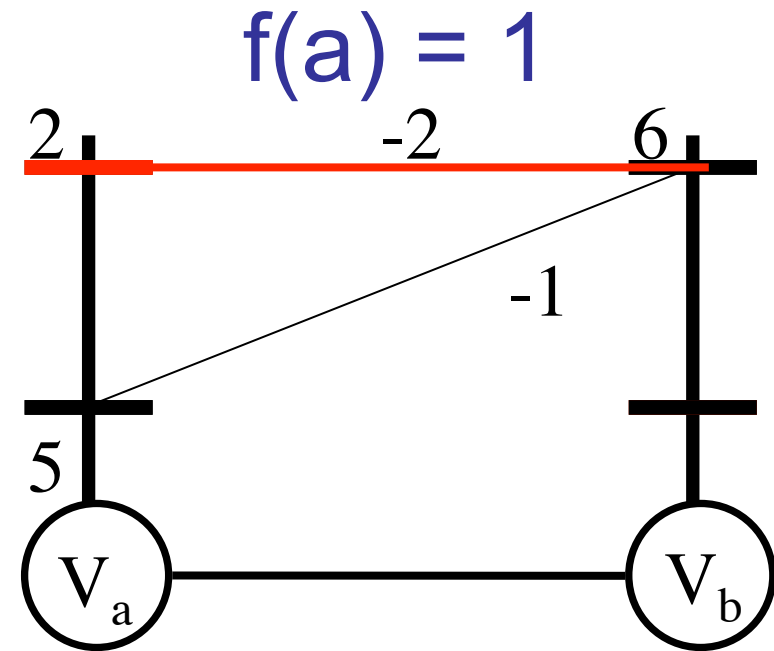
Choose the **right** constant

$$\theta'_{b;k} = q_{b;k}$$

Two Variables



$$\theta'_{b;0} = q_{b;0}$$



$$\theta'_{b;1} = q_{b;1}$$

We get all the min-marginals of V_b

Choose the *right* constant $\theta'_{b;k} = q_{b;k}$

Recap

We only need to know two sets of equations

General form of Reparameterization

$$\theta'_{a;i} = \theta_{a;i} + M_{ba;i} \quad \theta'_{b;k} = \theta_{b;k} + M_{ab;k}$$

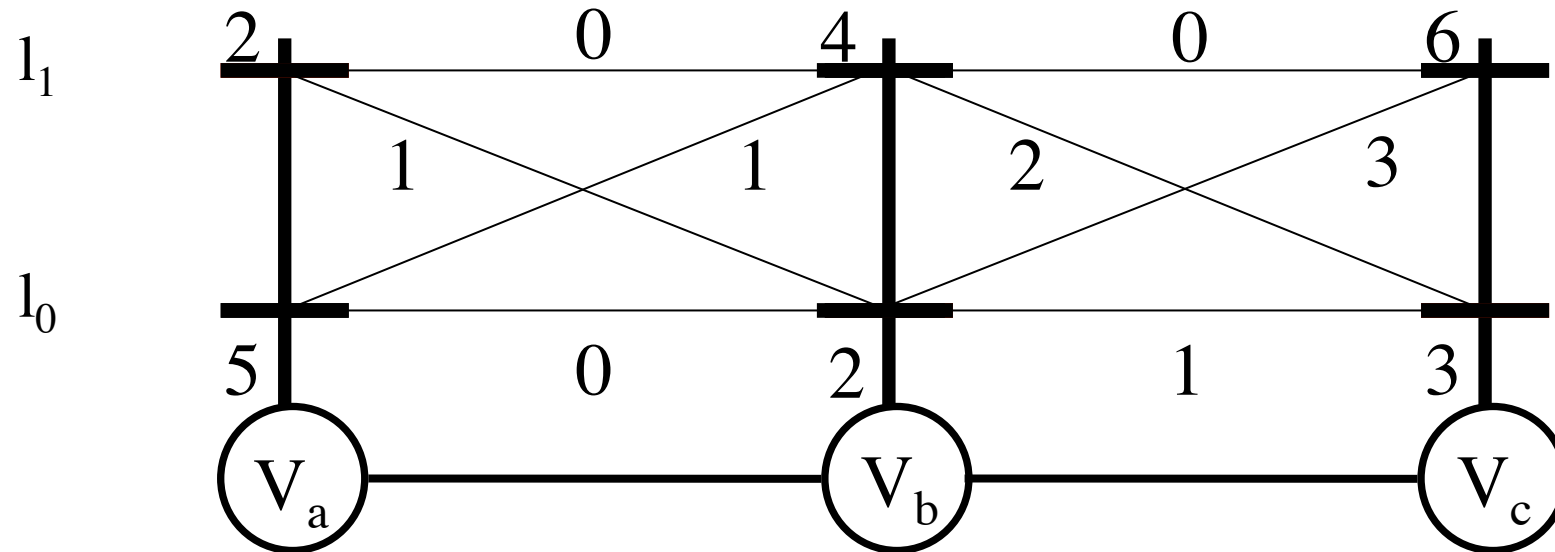
$$\theta'_{ab;ik} = \theta_{ab;ik} - M_{ab;k} - M_{ba;i}$$

Reparameterization of (a,b) in Belief Propagation

$$M_{ab;k} = \min_i \{ \theta_{a;i} + \theta_{ab;ik} \}$$

$$M_{ba;i} = 0$$

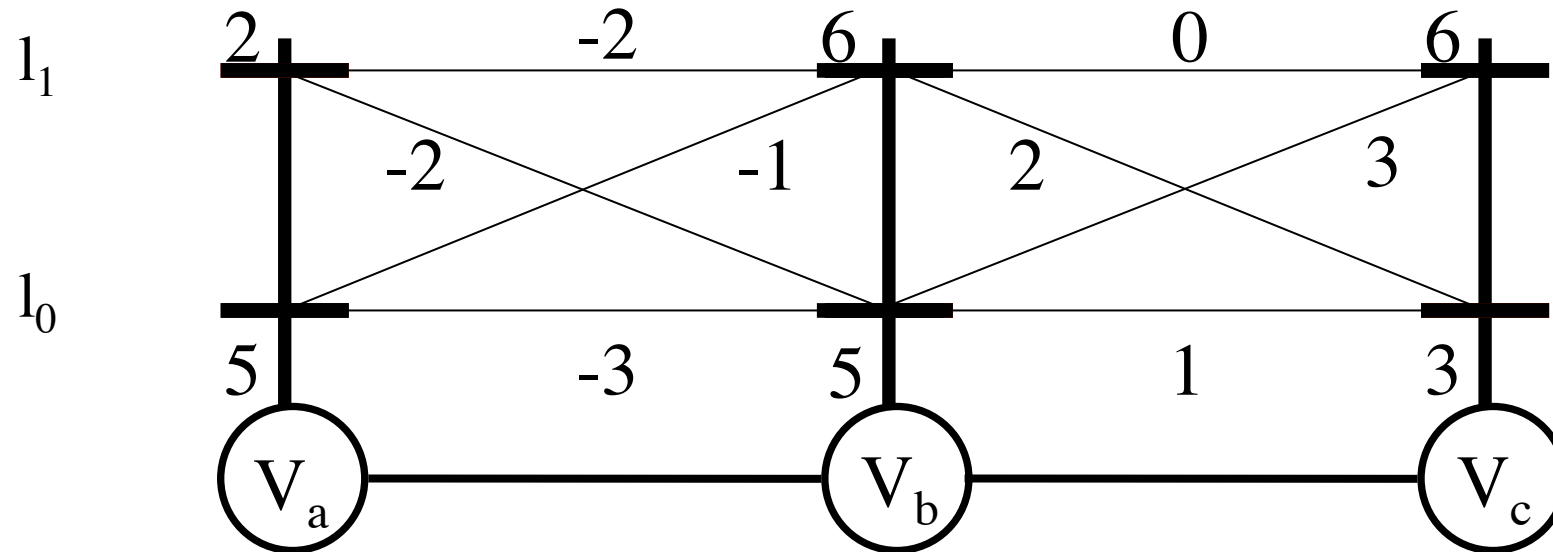
Three Variables



Reparameterize the edge (a,b) as before

Three Variables

$$f(a) = 1$$

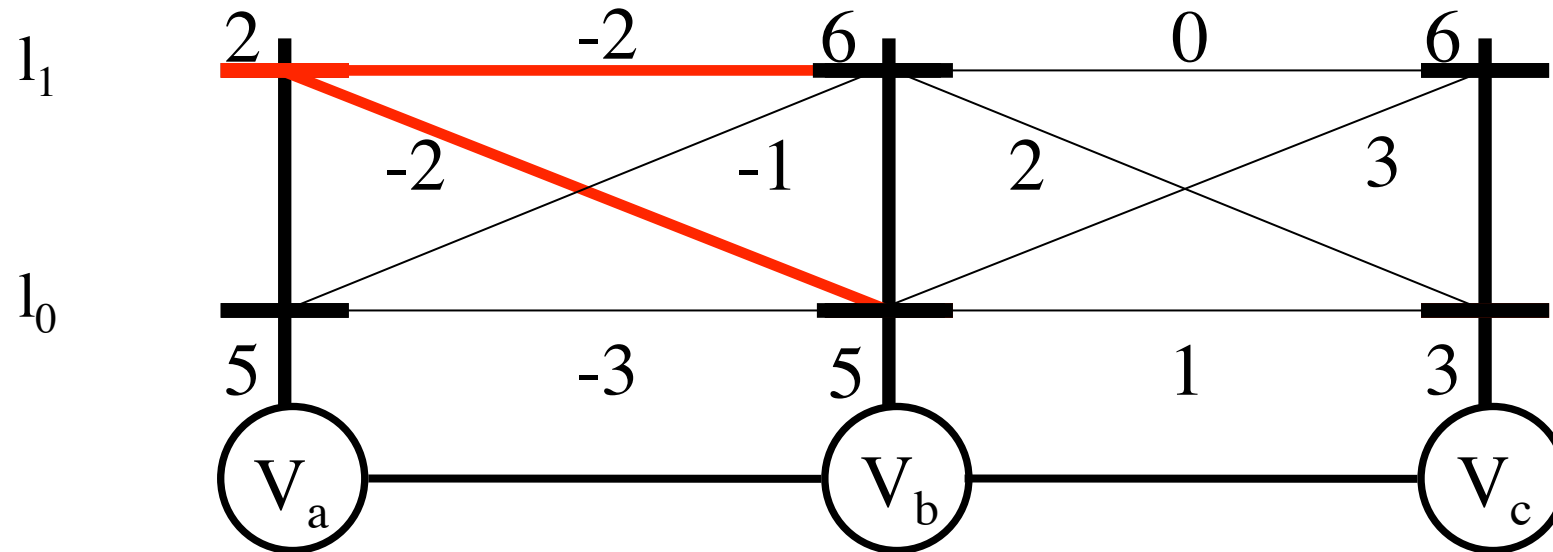


$$f(a) = 1$$

Reparameterize the edge (a,b) as before

Three Variables

$$f(a) = 1$$



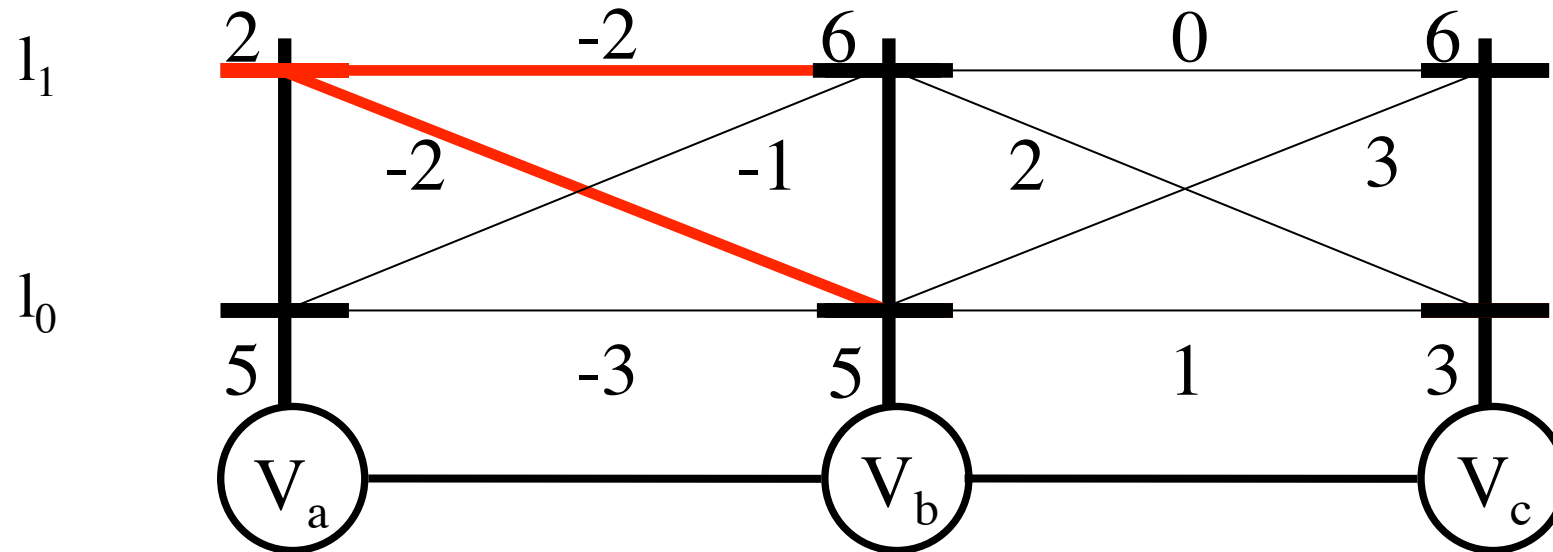
$$f(a) = 1$$

Reparameterize the edge (a,b) as before

Potentials along the red path add up to 0

Three Variables

$$f(a) = 1$$

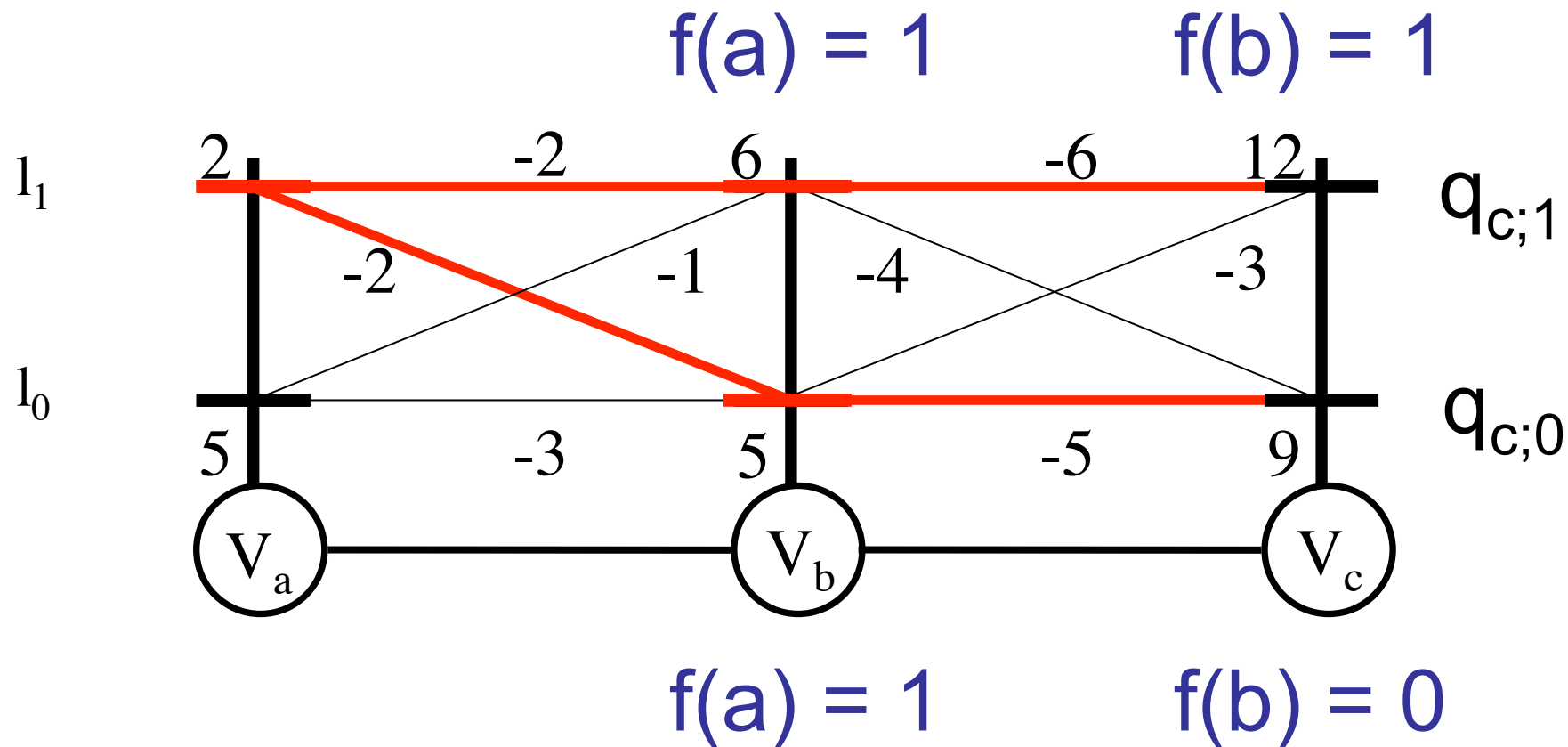


$$f(a) = 1$$

Reparameterize the edge (b,c) as before

Potentials along the red path add up to 0

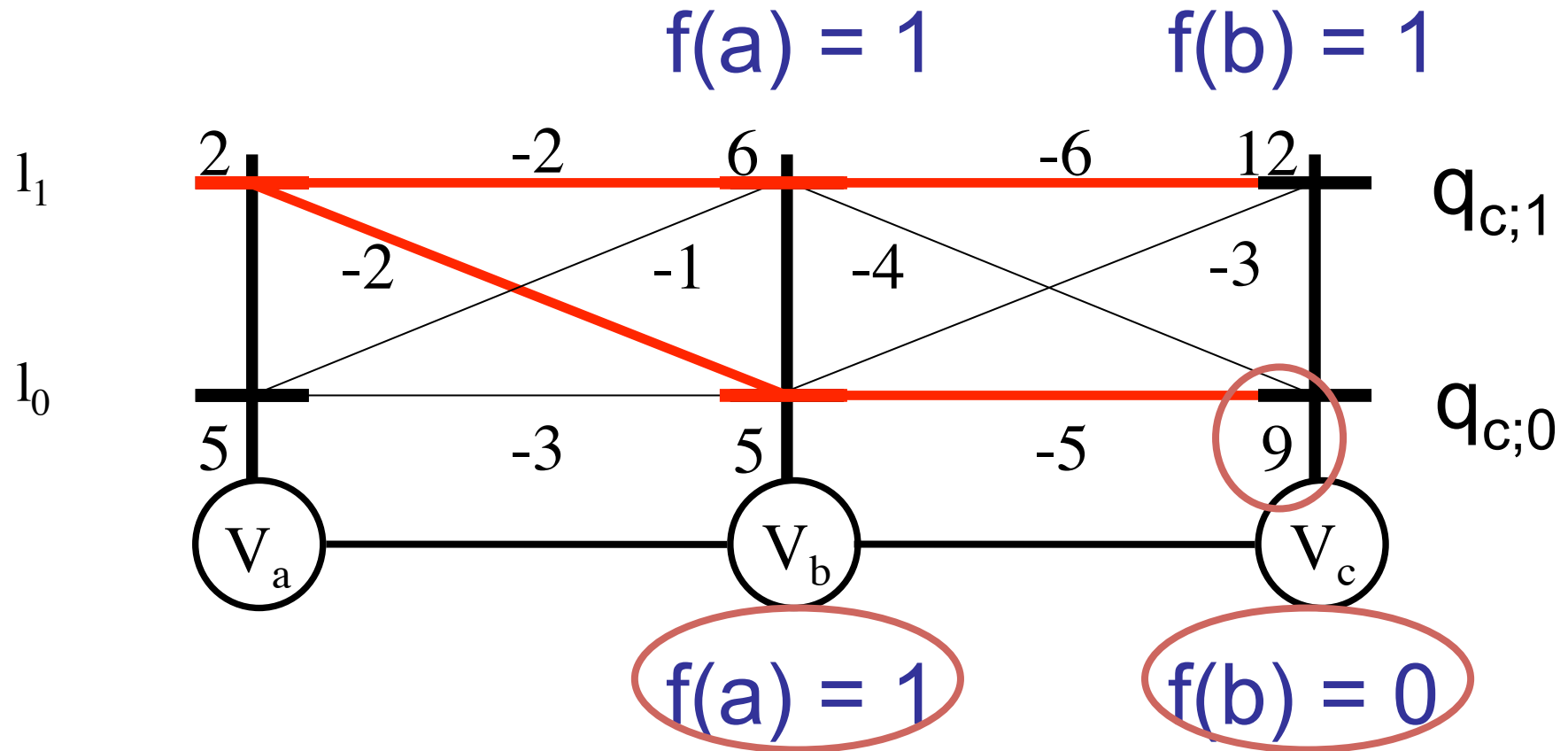
Three Variables



Reparameterize the edge (b,c) as before

Potentials along the red path add up to 0

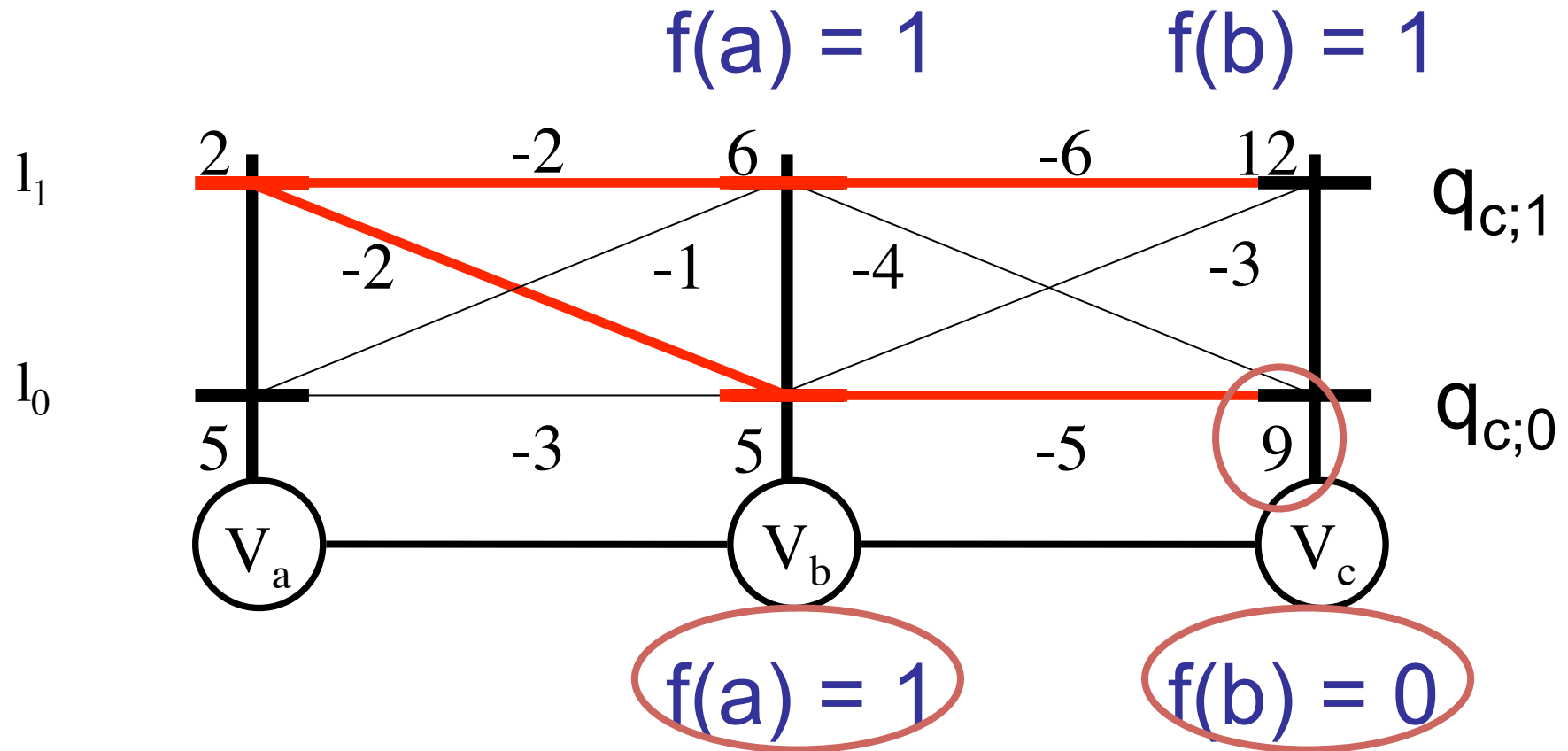
Three Variables



$$f^*(c) = 0 \quad f^*(b) = 0 \quad f^*(a) = 1$$

Generalizes to any length chain

Three Variables



$f^*(c) = 0 \quad f^*(b) = 0 \quad f^*(a) = 1$

Only Dynamic Programming

Why Dynamic Programming?

3 variables \equiv 2 variables + book-keeping

n variables \equiv (n-1) variables + book-keeping

Start from left, go to right

Reparameterize current edge (a,b)

$$M_{ab;k} = \min_i \{ \theta_{a;i} + \theta_{ab;ik} \}$$

$$\theta'_{b;k} = \theta_{b;k} + M_{ab;k} \quad \theta'_{ab;ik} = \theta_{ab;ik} - M_{ab;k}$$

Repeat

Why Dynamic Programming?

Messages Message Passing

Why stop at dynamic programming?

Start from left, go to right

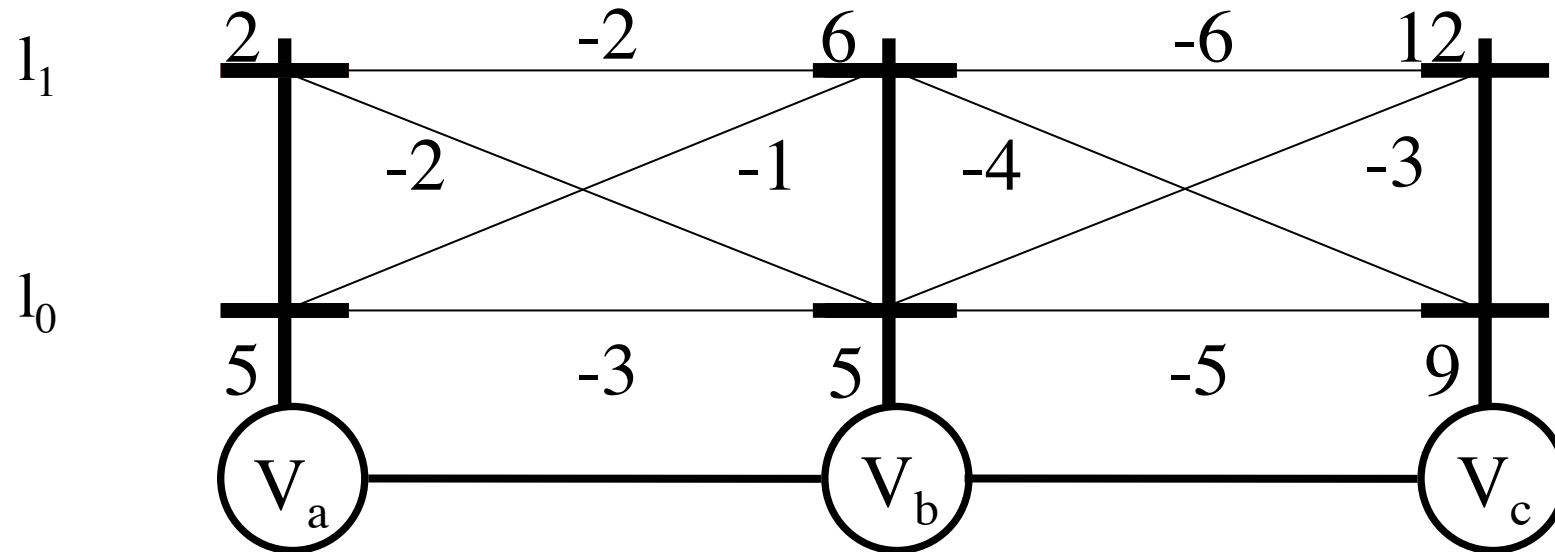
Reparameterize current edge (a,b)

$$M_{ab;k} = \min_i \{ \theta_{a;i} + \theta_{ab;ik} \}$$

$$\theta'_{b;k} = \theta_{b;k} + M_{ab;k} \quad \theta'_{ab;ik} = \theta_{ab;ik} - M_{ab;k}$$

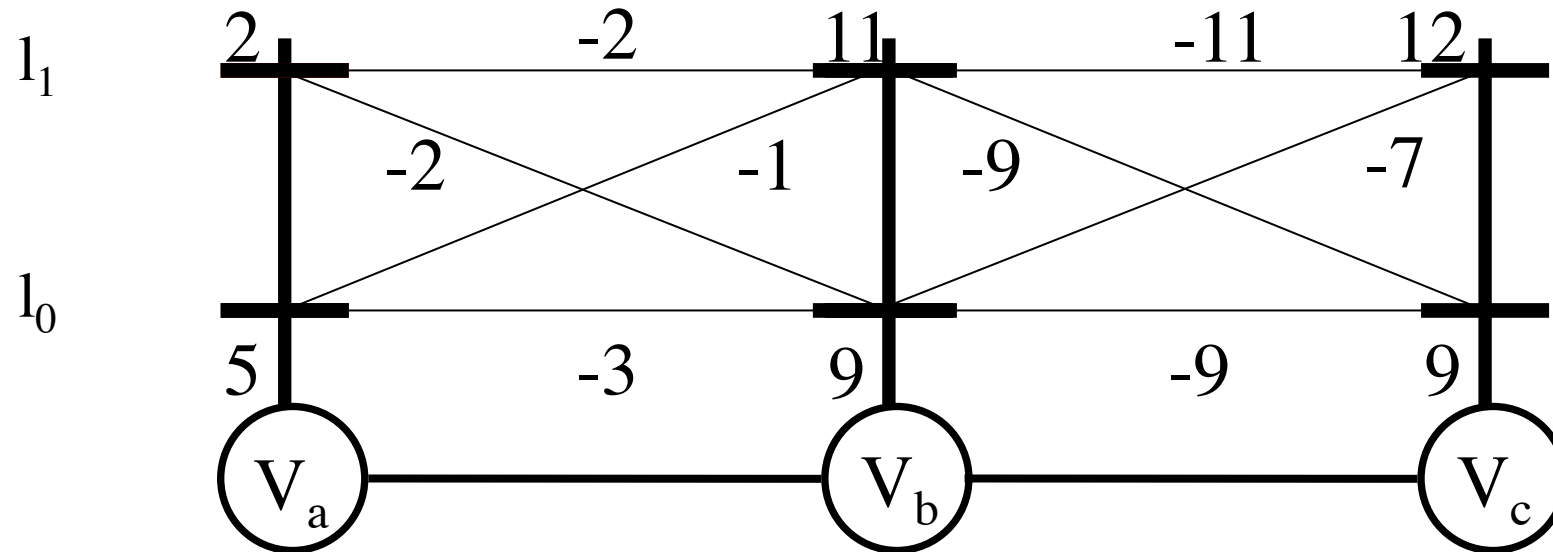
Repeat

Three Variables



Reparameterize the edge (c,b) as before

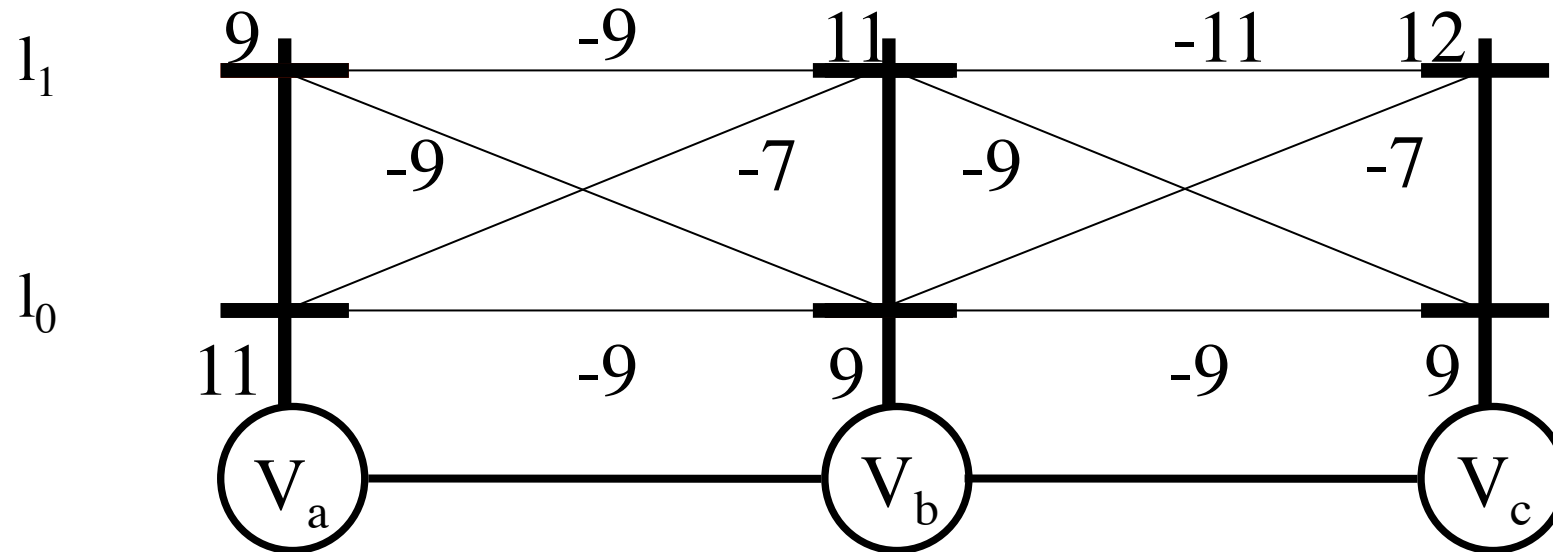
Three Variables



Reparameterize the edge (c,b) as before

$$\theta'_{b;i} = q_{b;i}$$

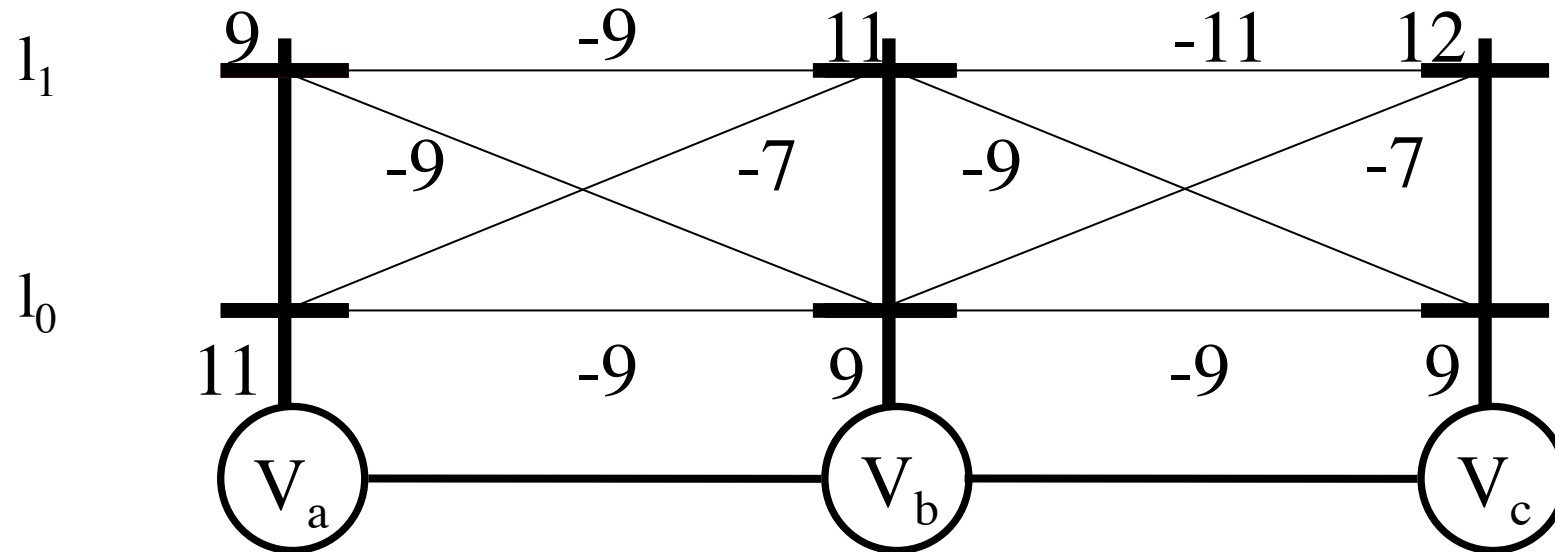
Three Variables



Reparameterize the edge (b,a) as before

$$\theta'_{a;i} = q_{a;i}$$

Three Variables



Forward Pass →

← Backward Pass

All min-marginals are computed

Chains



Reparameterize the edge (1,2)

Chains



Reparameterize the edge (2,3)

Chains



Reparameterize the edge (3,4)

Chains



Reparameterize the edge $(n-1, n)$

Min-marginals $e_n(i)$ for all labels

Belief Propagation on Chains

Start from left, go to right

Reparameterize current edge (a,b)

$$M_{ab;k} = \min_i \{ \theta_{a;i} + \theta_{ab;ik} \}$$

$$\theta'_{b;k} = \theta_{b;k} + M_{ab;k} \quad \theta'_{ab;ik} = \theta_{ab;ik} - M_{ab;k}$$

Repeat till the end of the chain

Start from right, go to left

Repeat till the end of the chain

Belief Propagation on Chains

- Generalizes to chains of any length
- A way of computing reparam constants
- Forward Pass - Start to End
 - MAP estimate
 - Min-marginals of final variable
- Backward Pass - End to start
 - All other min-marginals

Computational Complexity

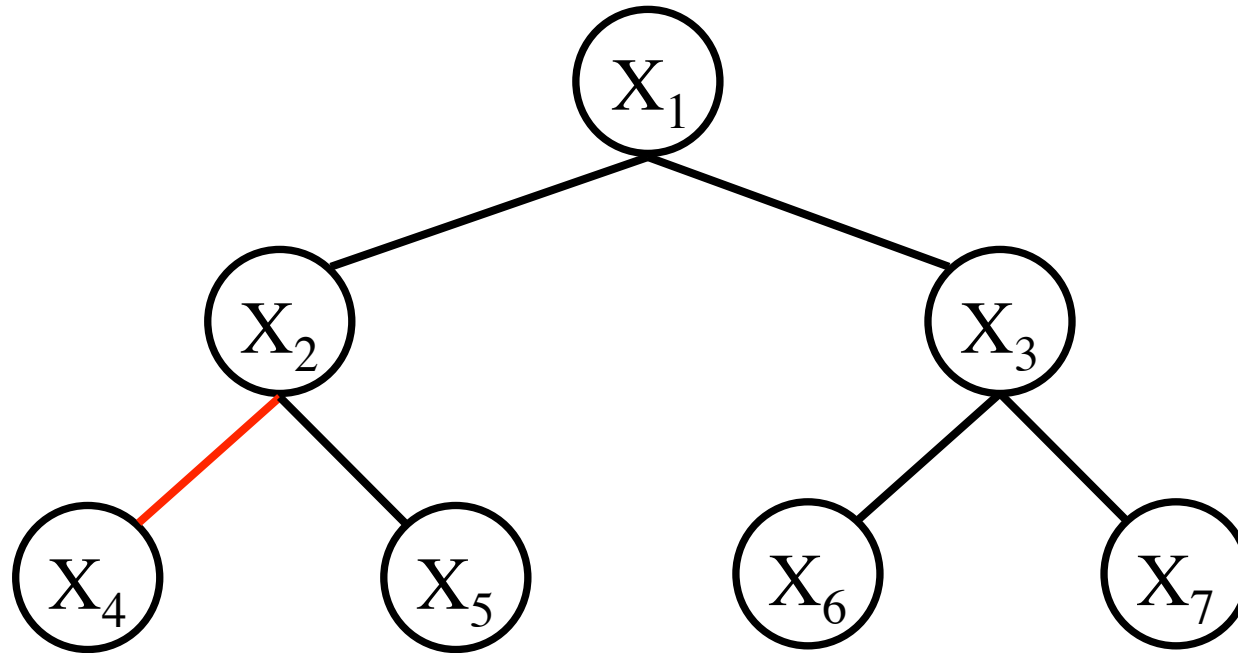
Number of reparameterization constants = $(n-1)h$

Complexity for each constant = $O(h)$

Total complexity = $O(nh^2)$

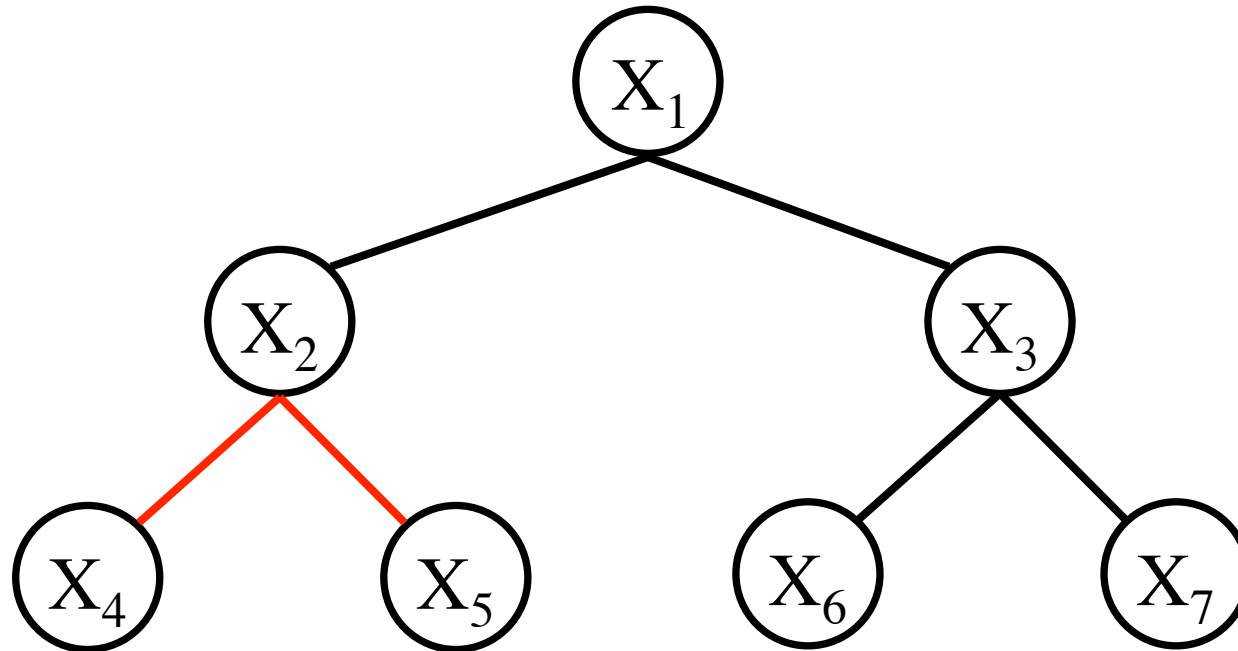
Better than brute-force $O(h^n)$

Trees



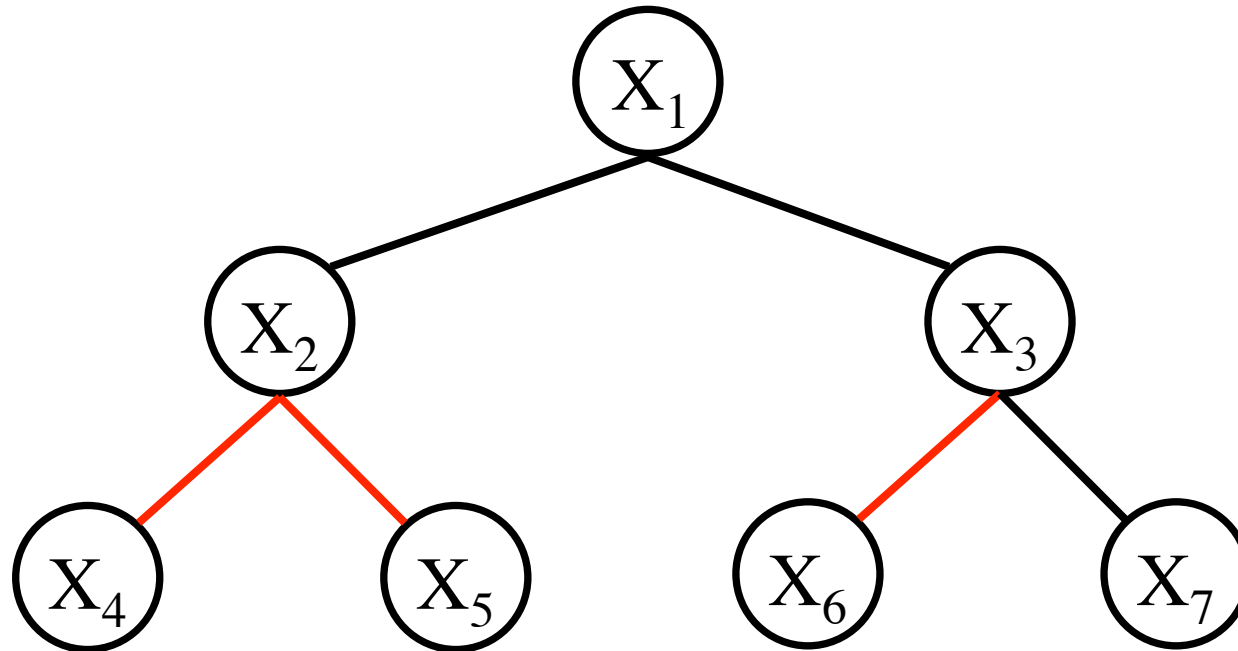
Reparameterize the edge (4,2)

Trees



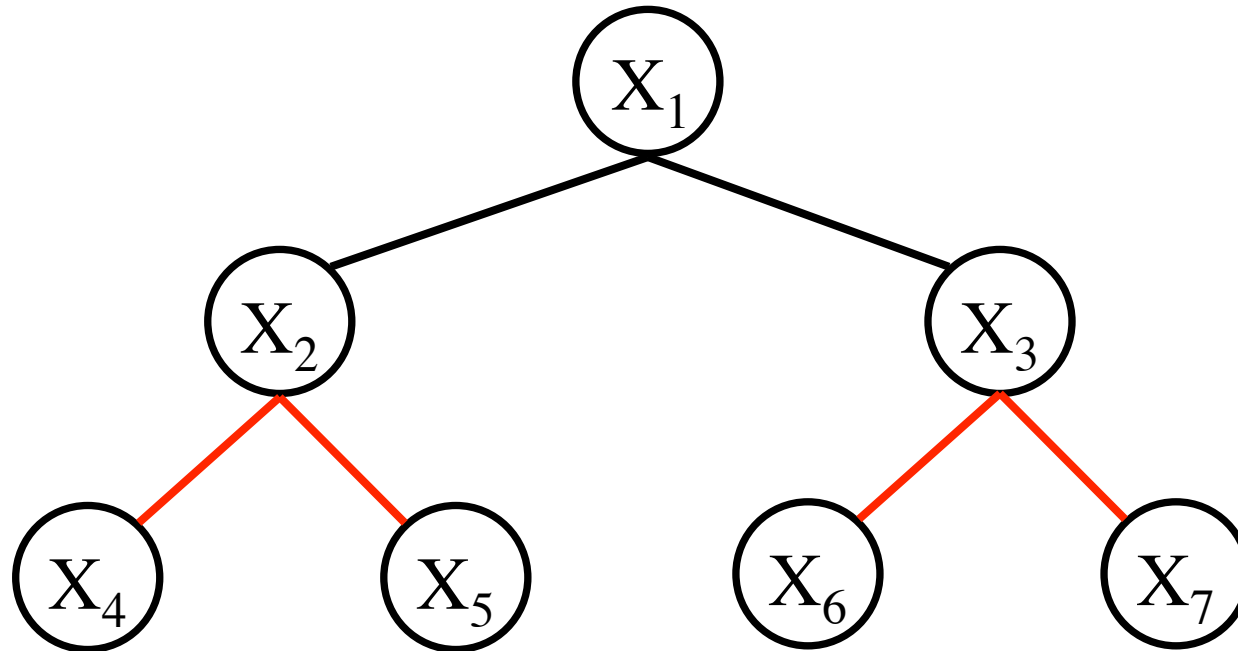
Reparameterize the edge (5,2)

Trees



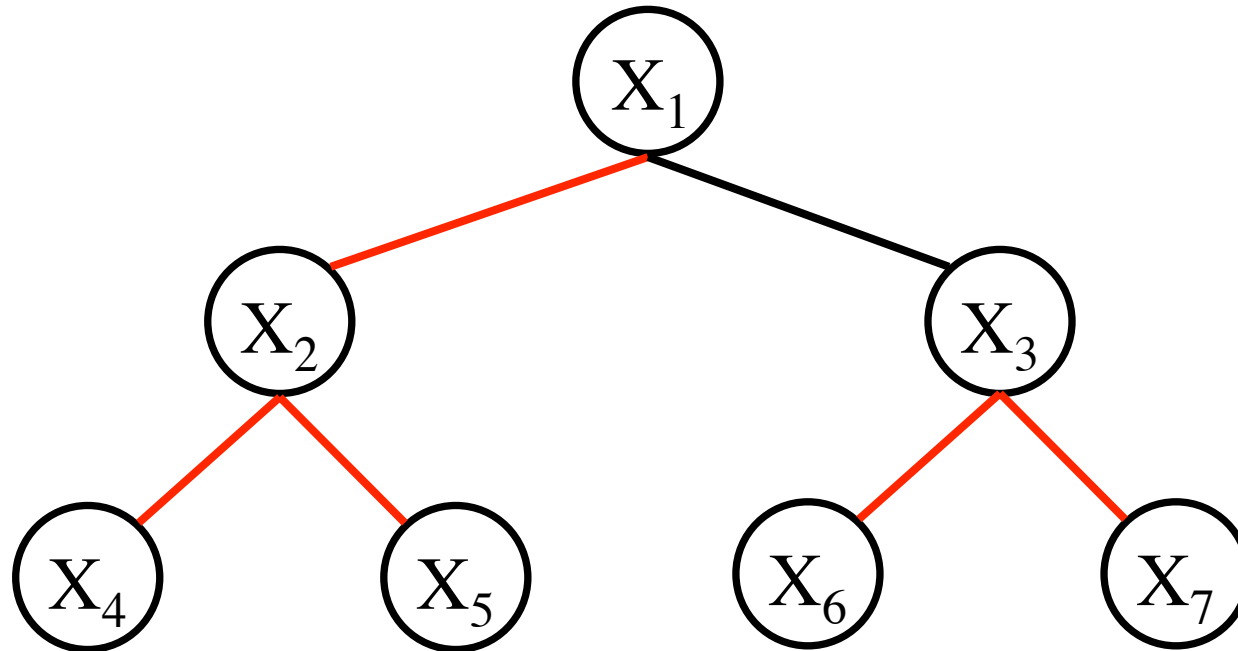
Reparameterize the edge (6,3)

Trees



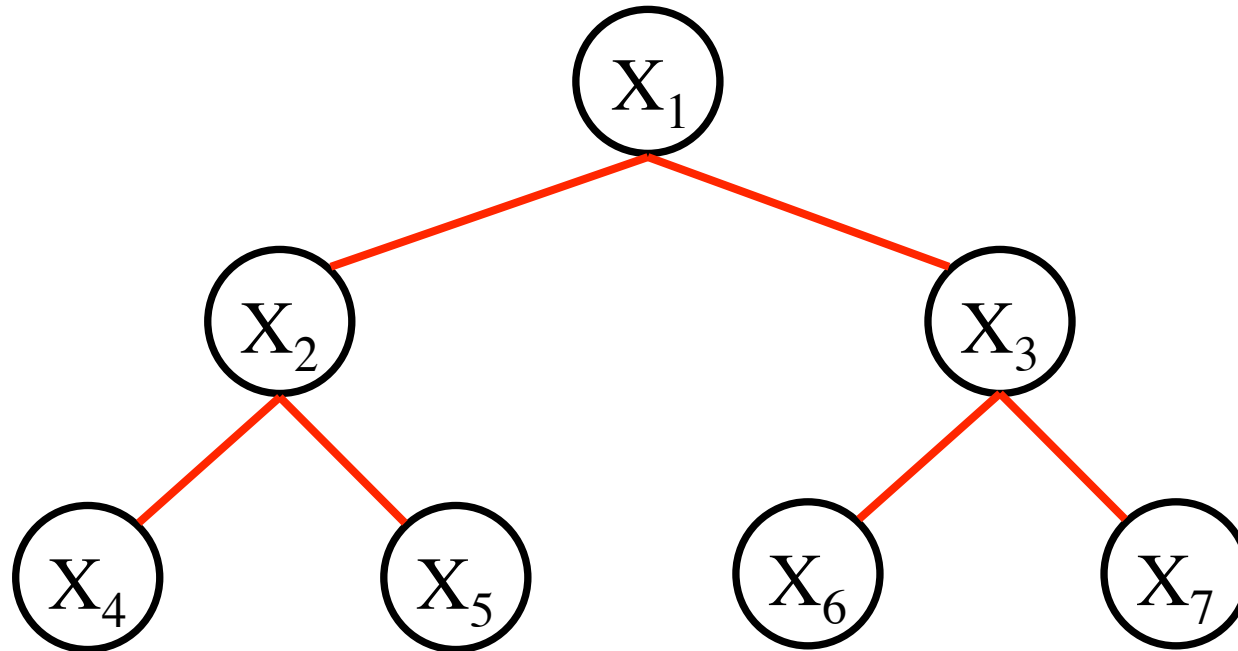
Reparameterize the edge (7,3)

Trees



Reparameterize the edge (2,1)

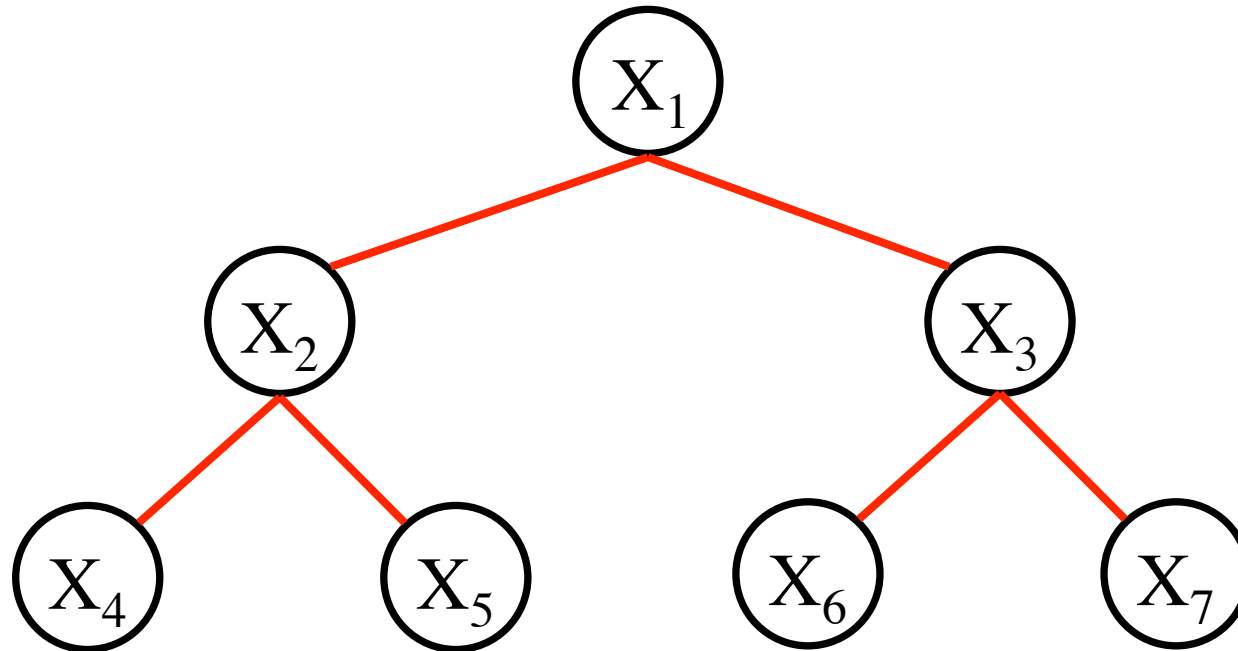
Trees



Reparameterize the edge (3,1)

Min-marginals $e_1(i)$ for all labels

Trees



Start from leaves and move towards root

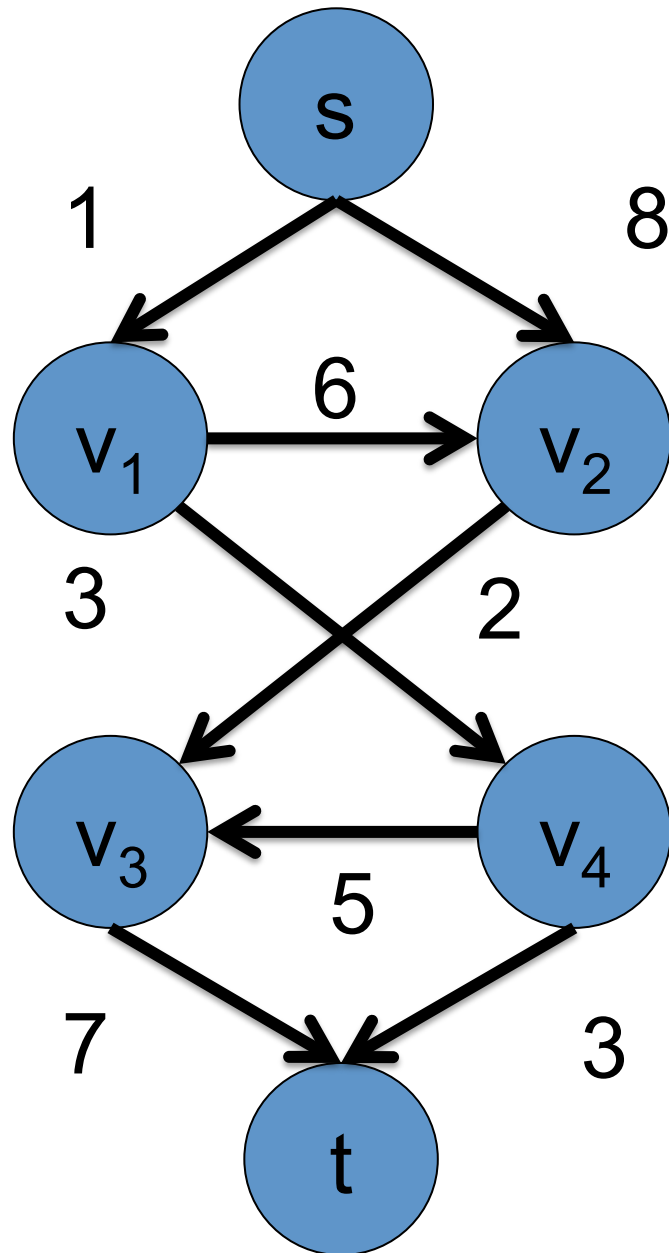
Pick the minimum of min-marginals

Backtrack to find the best labeling \mathbf{x}

Outline

- Preliminaries
 - **s-t Flow**
 - s-t Cut
 - Flows vs. Cuts
- Maximum Flow
- Algorithms
- Energy minimization with max flow/min cut

s-t Flow



Function flow: $A \rightarrow R$

Flow of arc \leq arc capacity

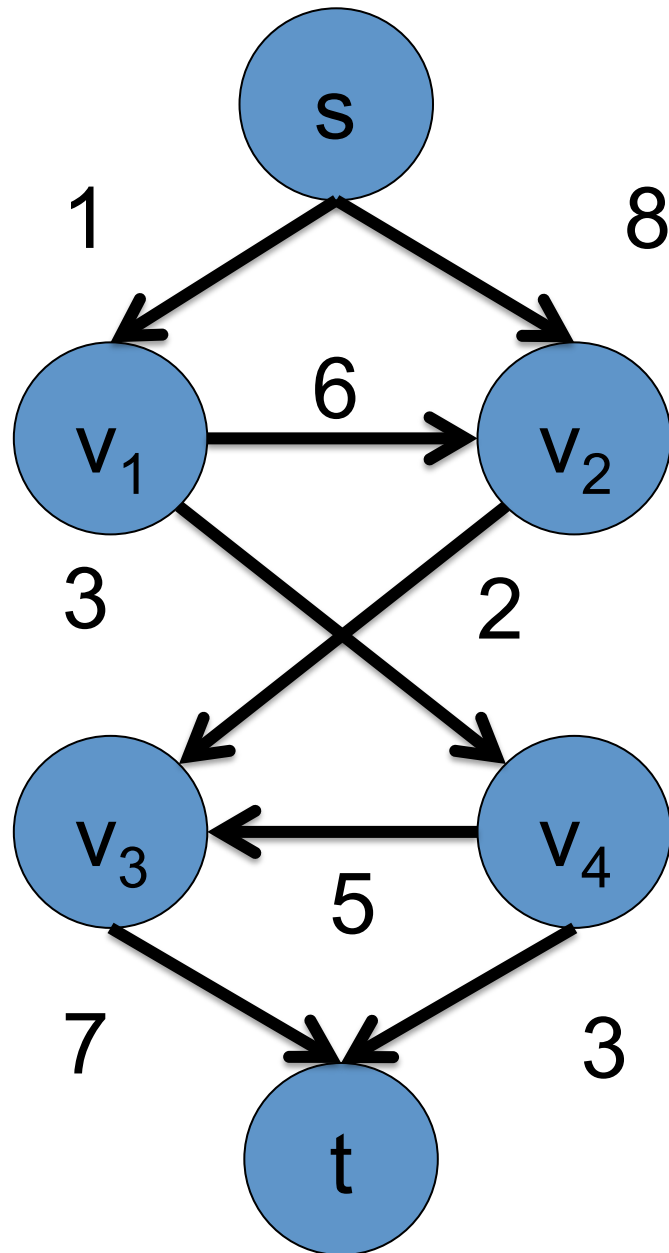
Flow is non-negative

For all vertex except s, t

Incoming flow

= Outgoing flow

s-t Flow



Function flow: $A \rightarrow R$

$\text{flow}(a) \leq c(a)$

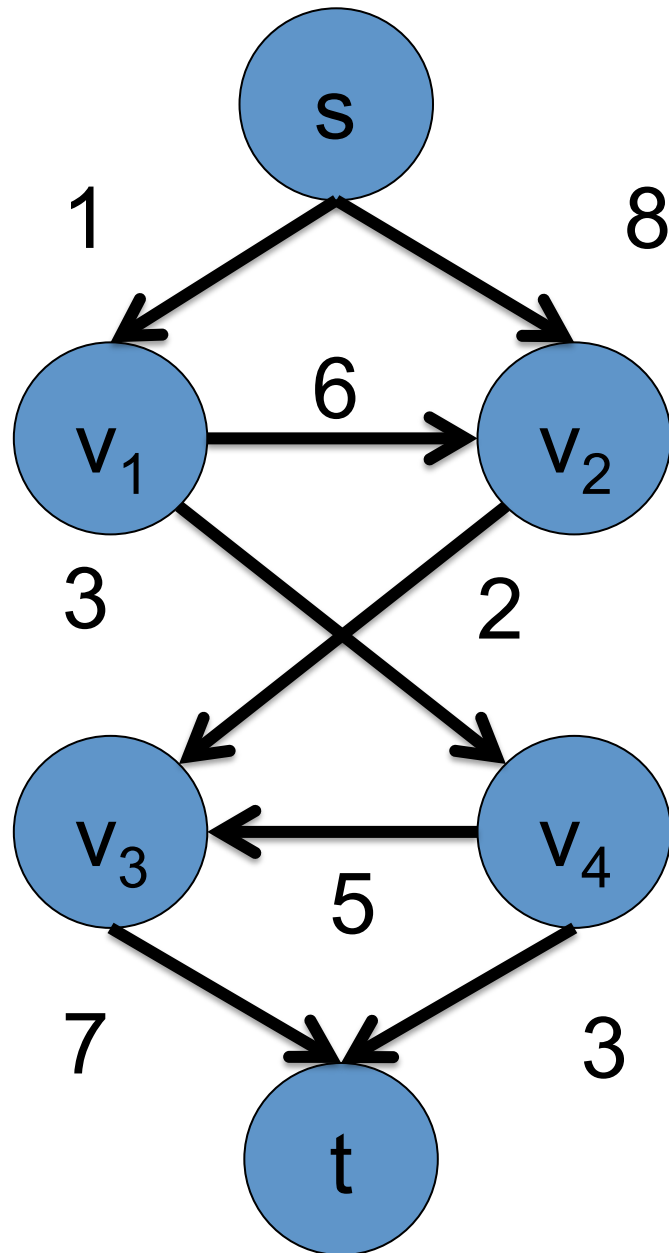
Flow is non-negative

For all vertex except s, t

Incoming flow

= Outgoing flow

s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

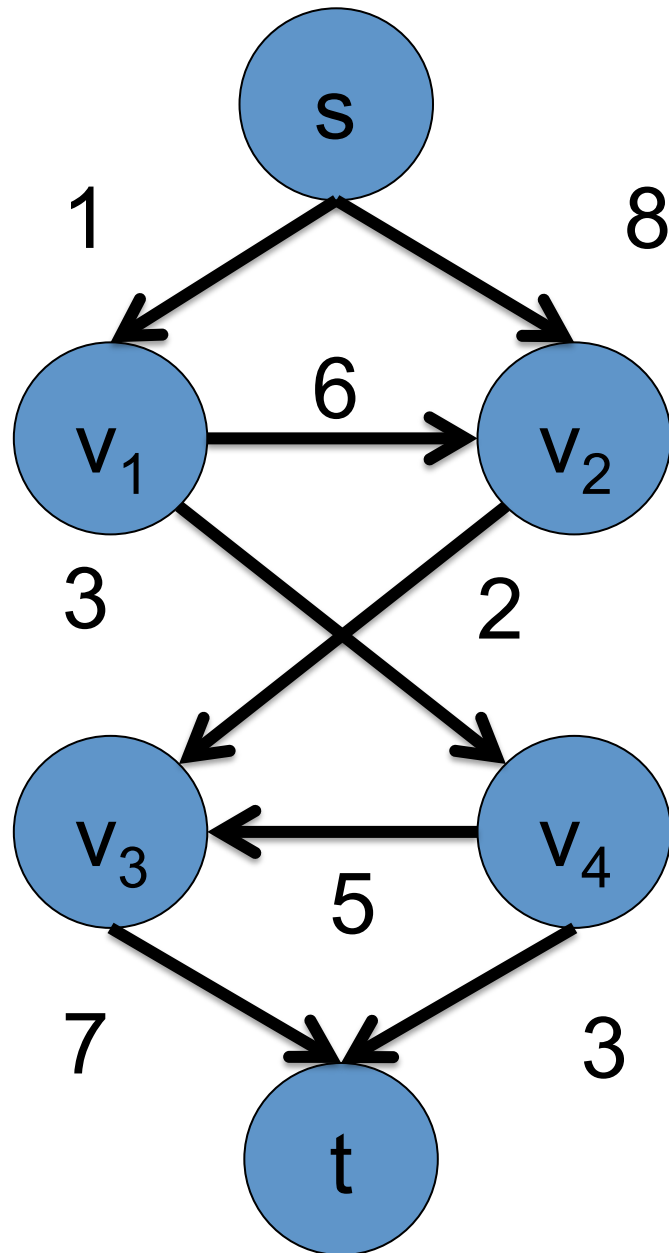
$$\text{flow}(a) \geq 0$$

For all vertex except s, t

Incoming flow

= Outgoing flow

s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

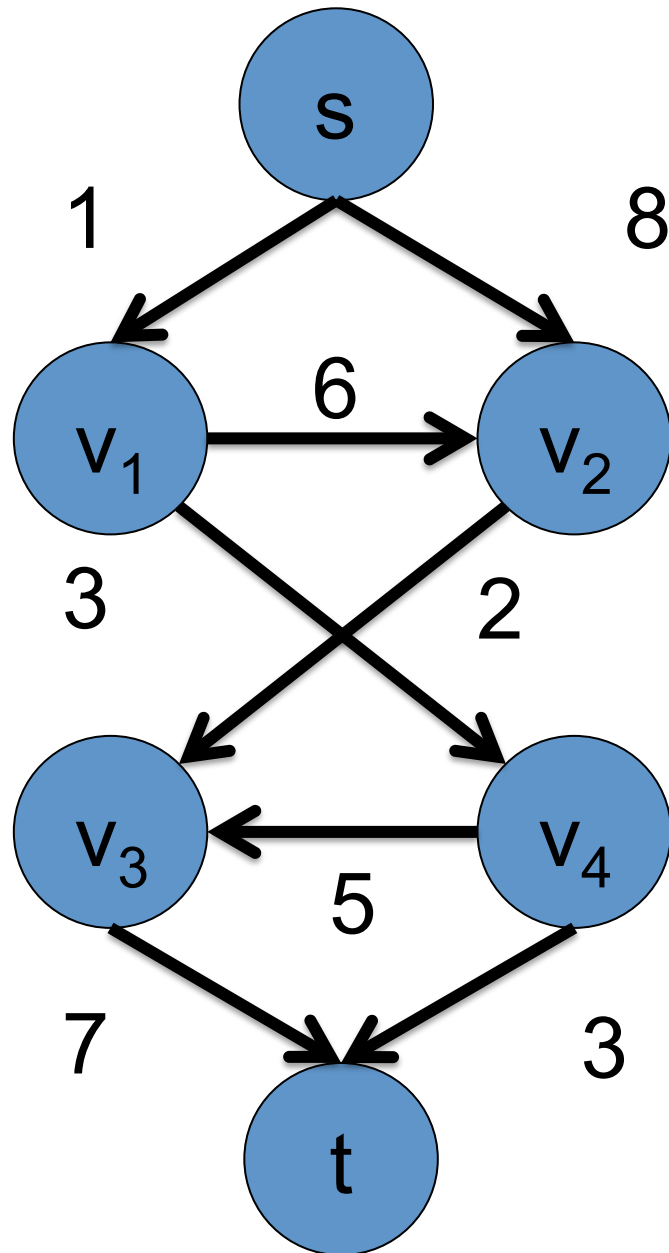
$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

Incoming flow

= Outgoing flow

s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

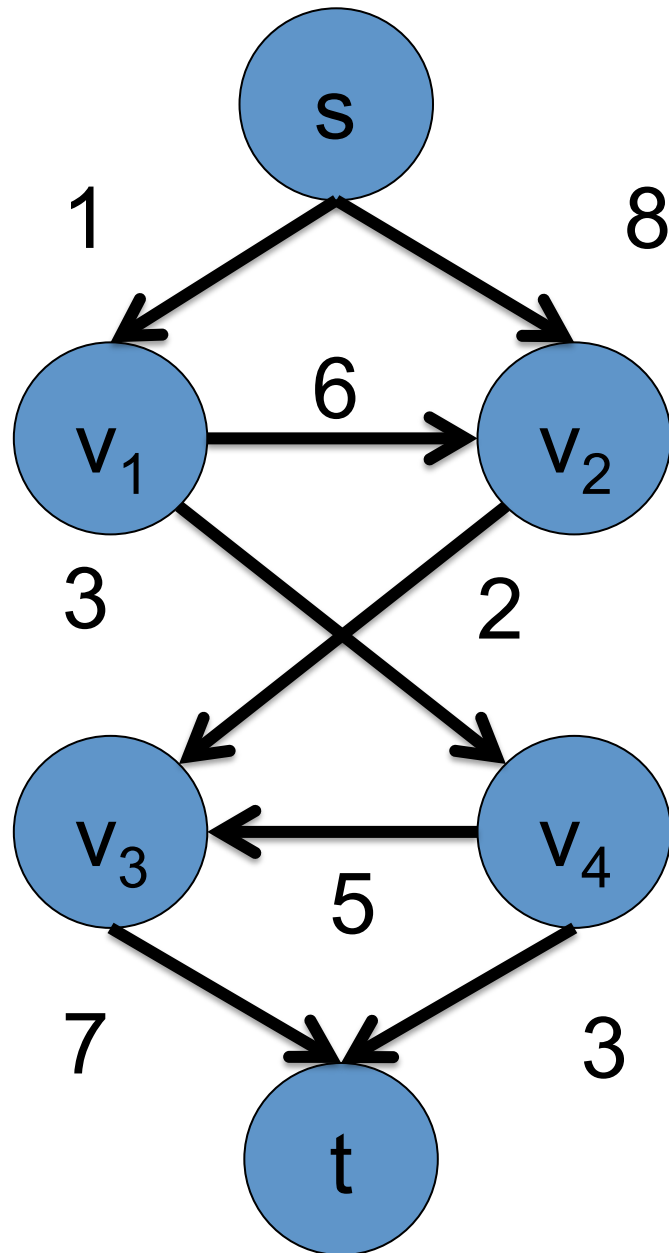
$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$\sum_{(u,v) \in A} \text{flow}((u,v))$$

= Outgoing flow

s-t Flow



Function flow: $A \rightarrow R$

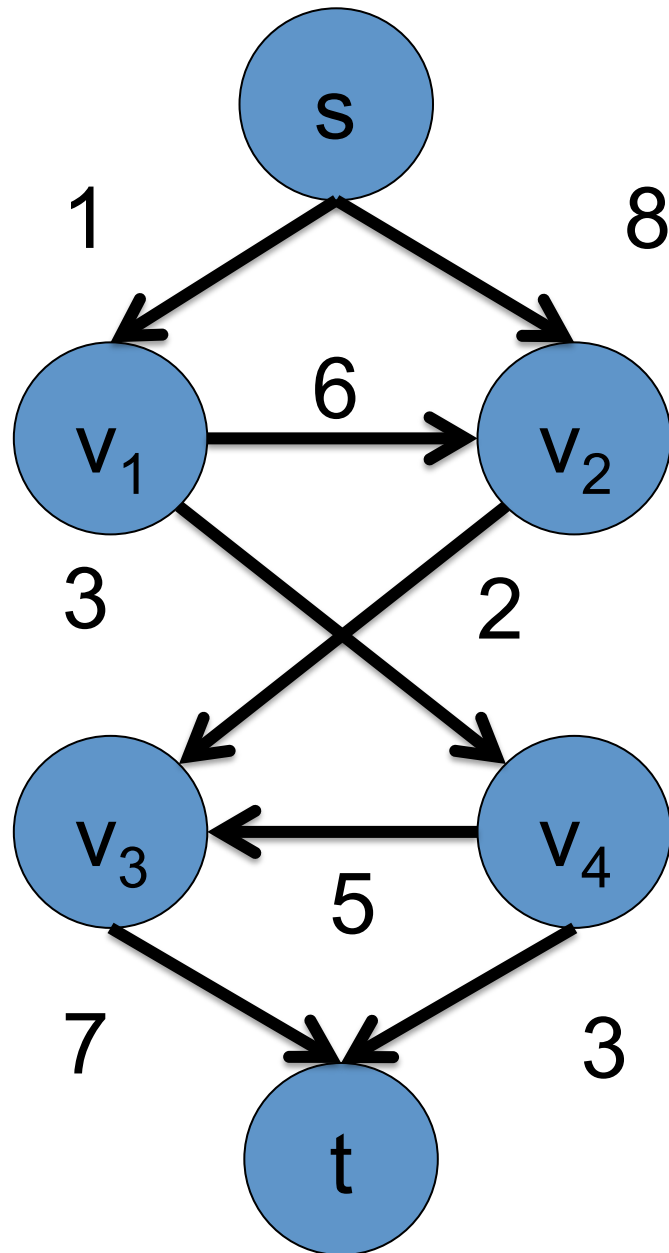
$$\text{flow}(a) \leq c(a)$$

$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$\begin{aligned} & \sum_{(u,v) \in A} \text{flow}((u,v)) \\ &= \sum_{(v,u) \in A} \text{flow}((v,u)) \end{aligned}$$

s-t Flow



Function flow: $A \rightarrow R$

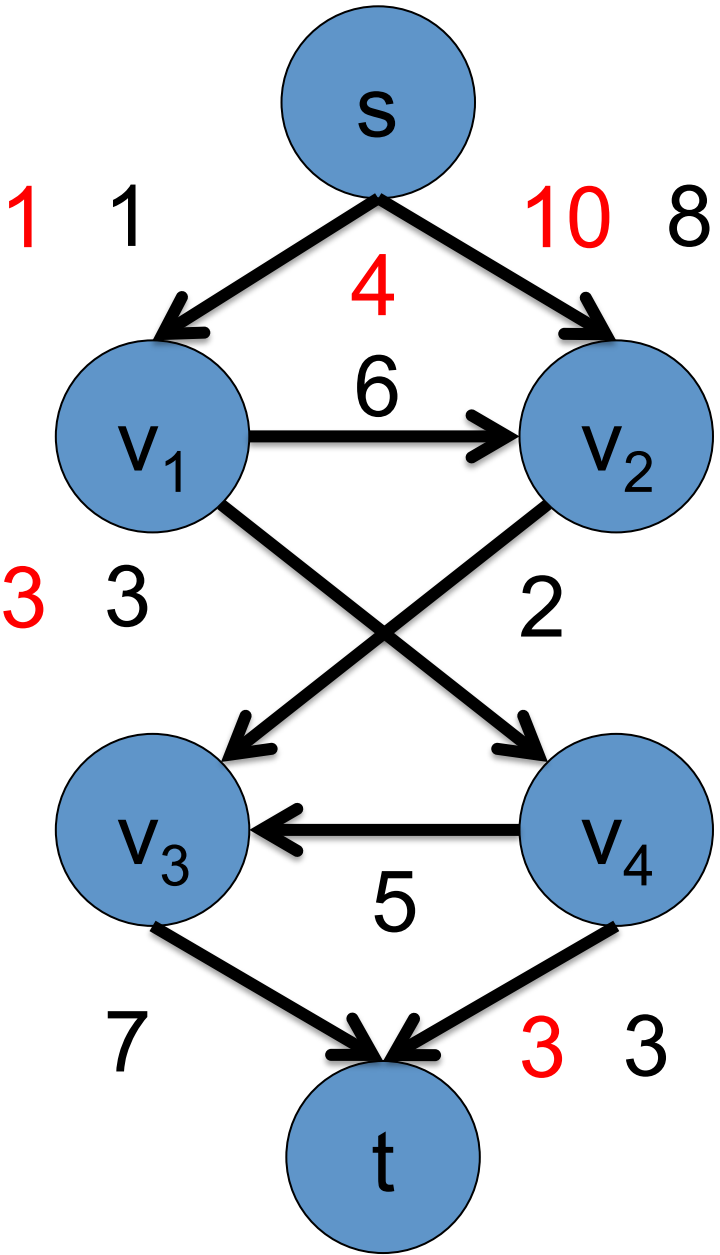
$$\text{flow}(a) \leq c(a)$$

$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$

s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

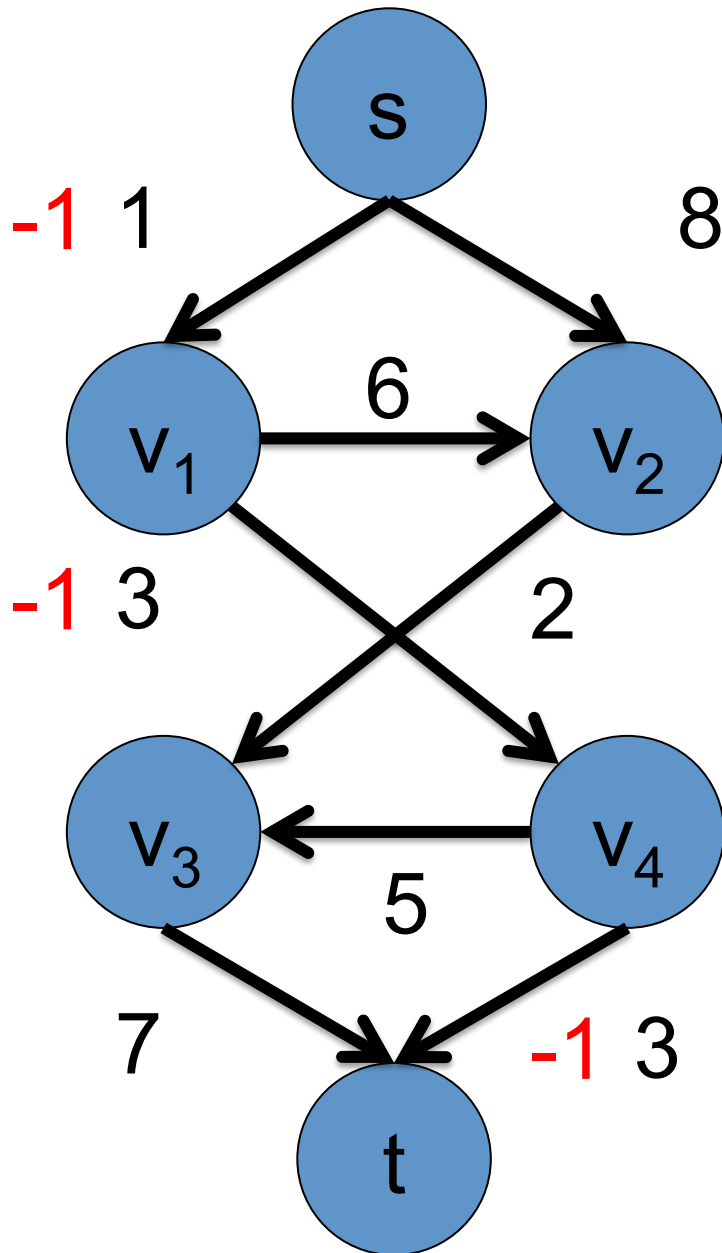
$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$



s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

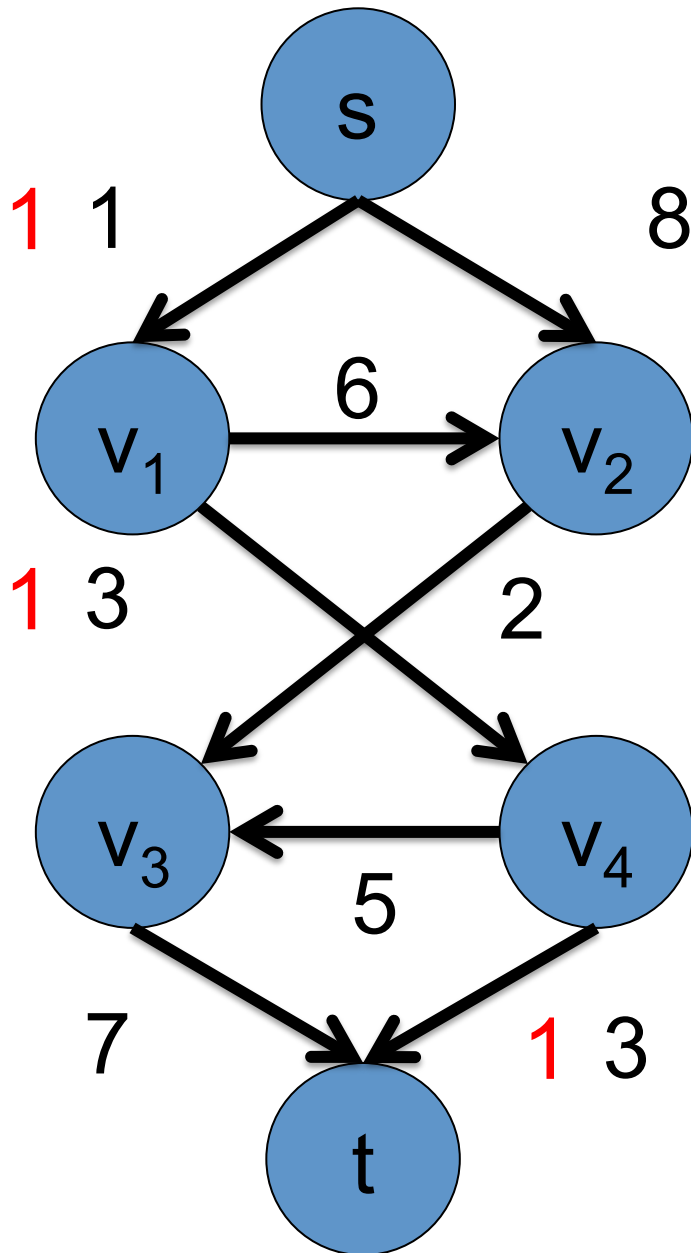
$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$

X

s-t Flow



Function flow: $A \rightarrow R$

$$\text{flow}(a) \leq c(a)$$

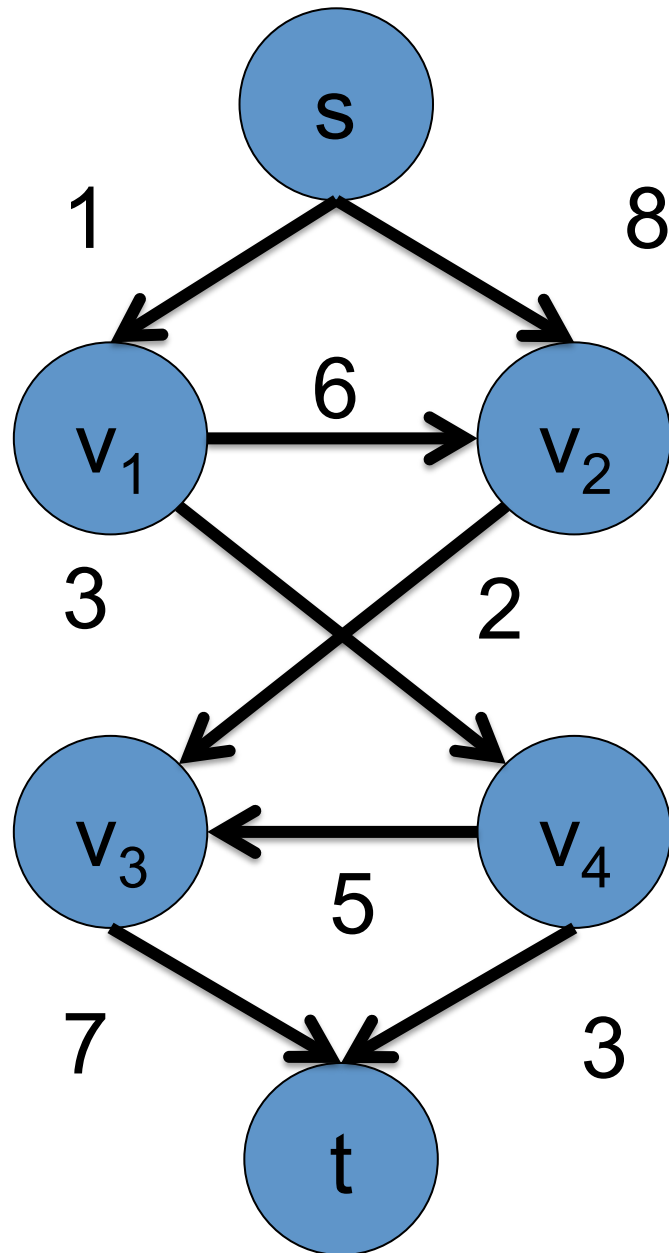
$$\text{flow}(a) \geq 0$$

For all $v \in V \setminus \{s, t\}$

$$E_{\text{flow}}(v) = 0$$

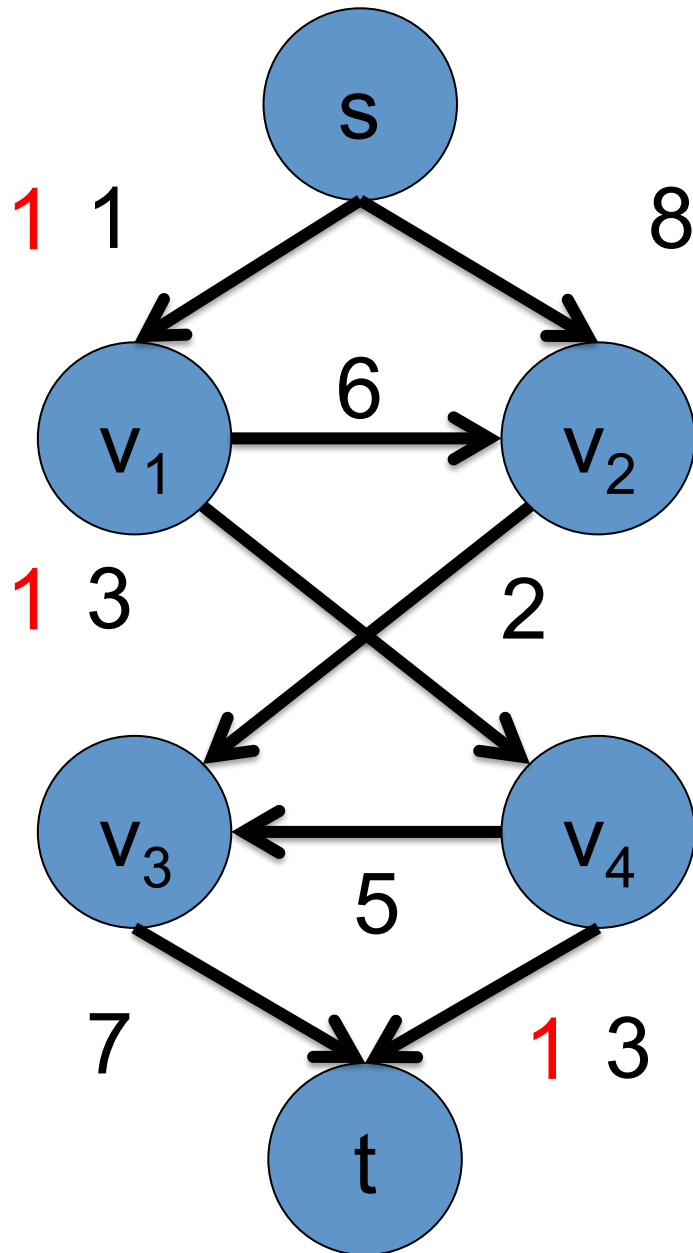


Value of s-t Flow



Outgoing flow of s
- Incoming flow of s

Value of s-t Flow



$$-E_{\text{flow}}(s) \quad E_{\text{flow}}(t)$$

$$\sum_{(s,v) \in A} \text{flow}((s,v))$$

$$- \sum_{(u,s) \in A} \text{flow}((u,s))$$

Value = 1

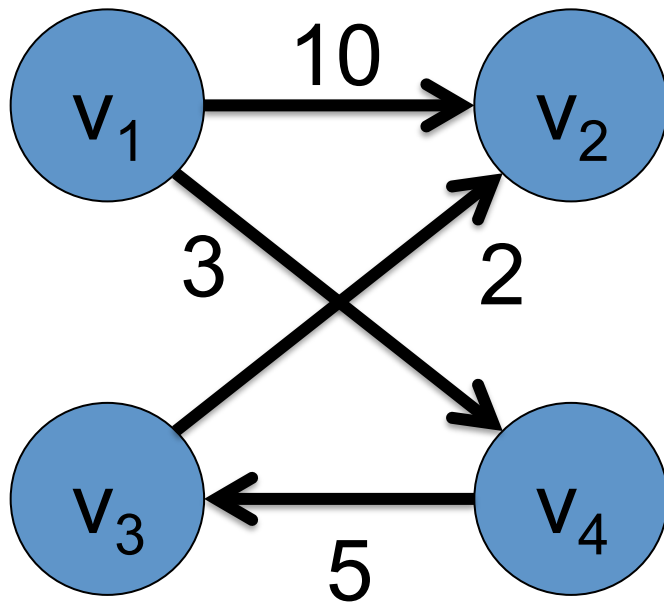
Outline

- Preliminaries
 - Functions and Excess Functions
 - s-t Flow
 - **s-t Cut**
 - Flows vs. Cuts
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- Algorithms
- Energy minimization with max flow/min cut

Cut

$$D = (V, A)$$

Let U be a subset of V



C is a set of arcs such that

- $(u, v) \in A$
- $u \in U$
- $v \in V \setminus U$

C is a cut in the digraph D

Cut

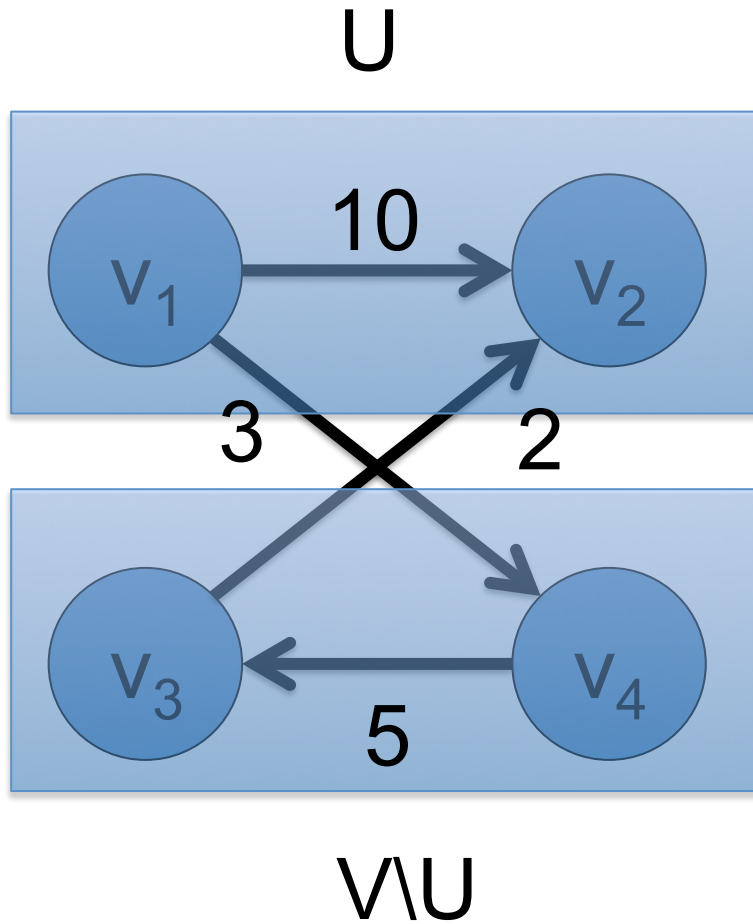
$$D = (V, A)$$

What is C?

$$\{(v_1, v_2), (v_1, v_4)\} ?$$

$$\{(v_1, v_4), (v_3, v_2)\} ?$$

$$\{(v_1, v_4)\} ?$$



Cut

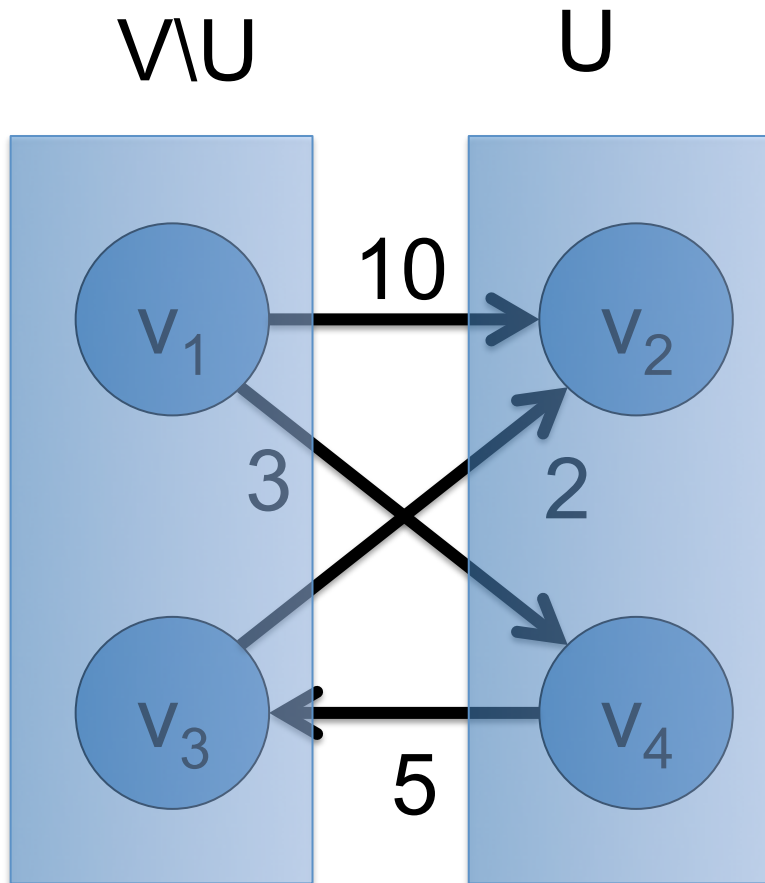
$$D = (V, A)$$

What is C?

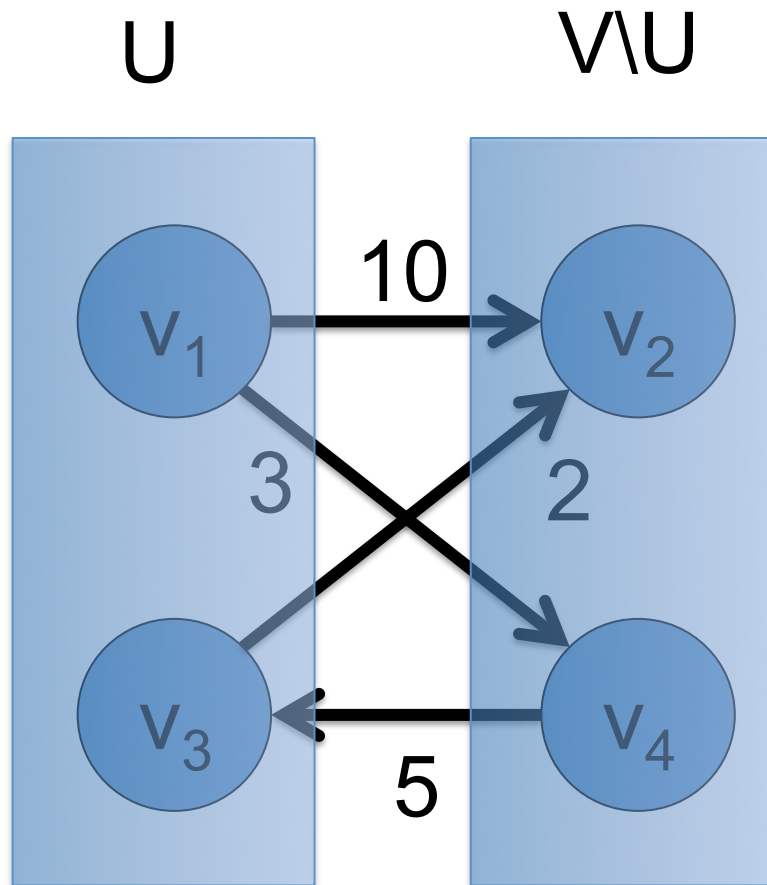
$$\{(v_1, v_2), (v_1, v_4), (v_3, v_2)\} ?$$

✓ $\{(v_4, v_3)\} ?$

$$\{(v_1, v_4), (v_3, v_2)\} ?$$



Cut



$$D = (V, A)$$

What is C ?



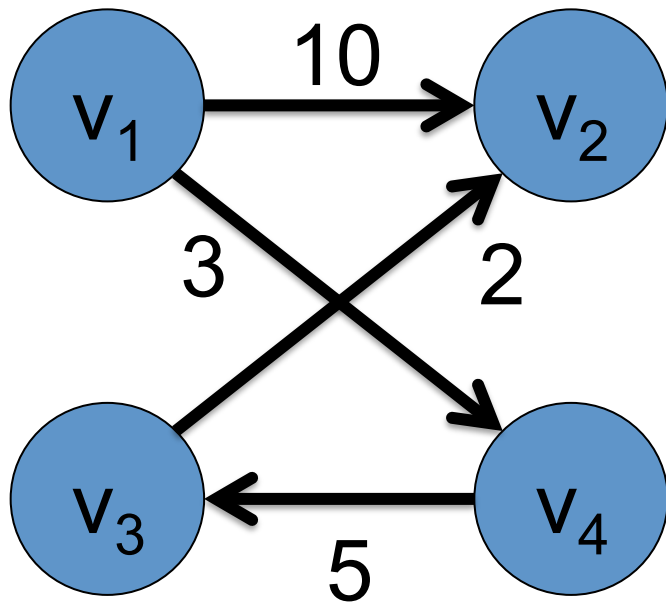
$\{(v_1, v_2), (v_1, v_4), (v_3, v_2)\}$?

$\{(v_3, v_2)\}$?

$\{(v_1, v_4), (v_3, v_2)\}$?

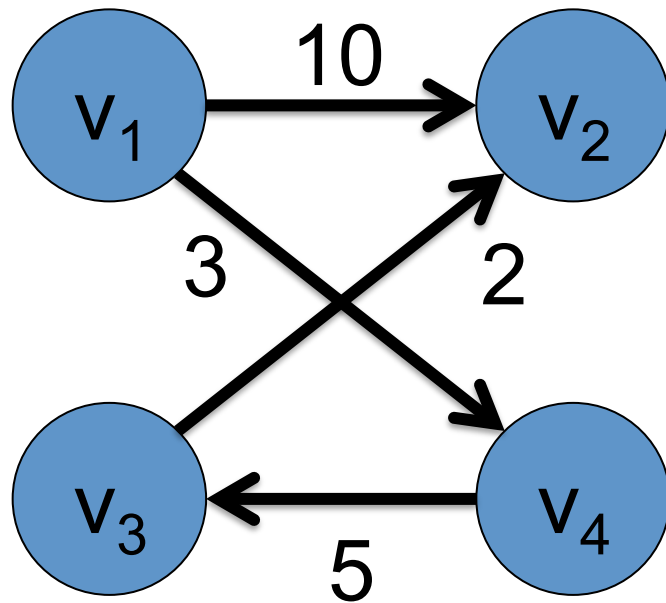
Cut

$$D = (V, A)$$



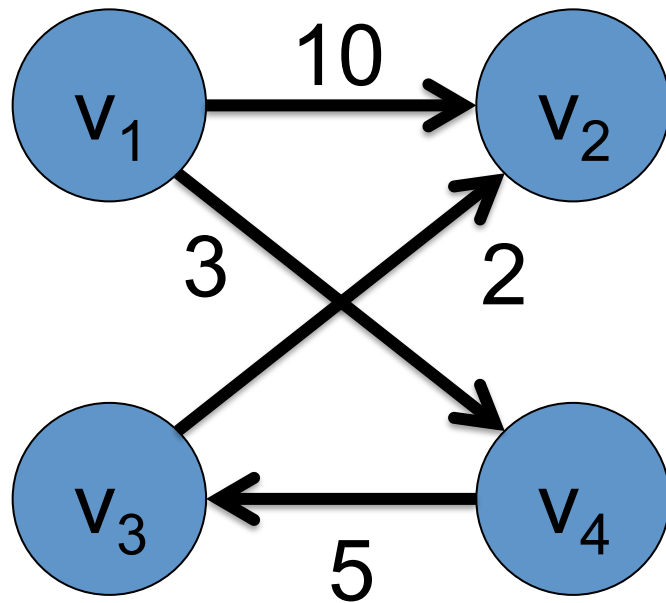
$$C = \text{out-arcs}(U)$$

Capacity of Cut



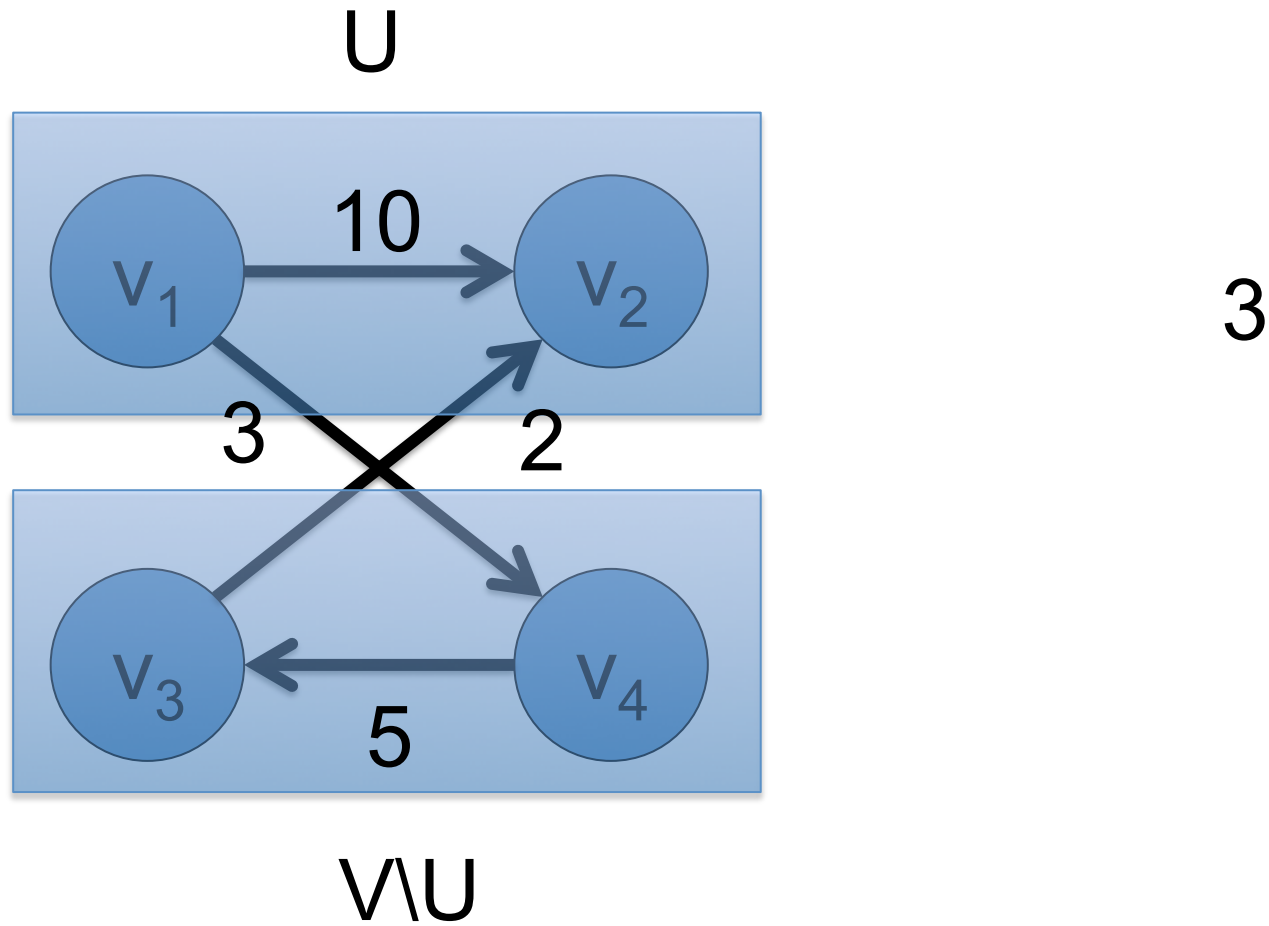
Sum of capacity of all arcs in C

Capacity of Cut

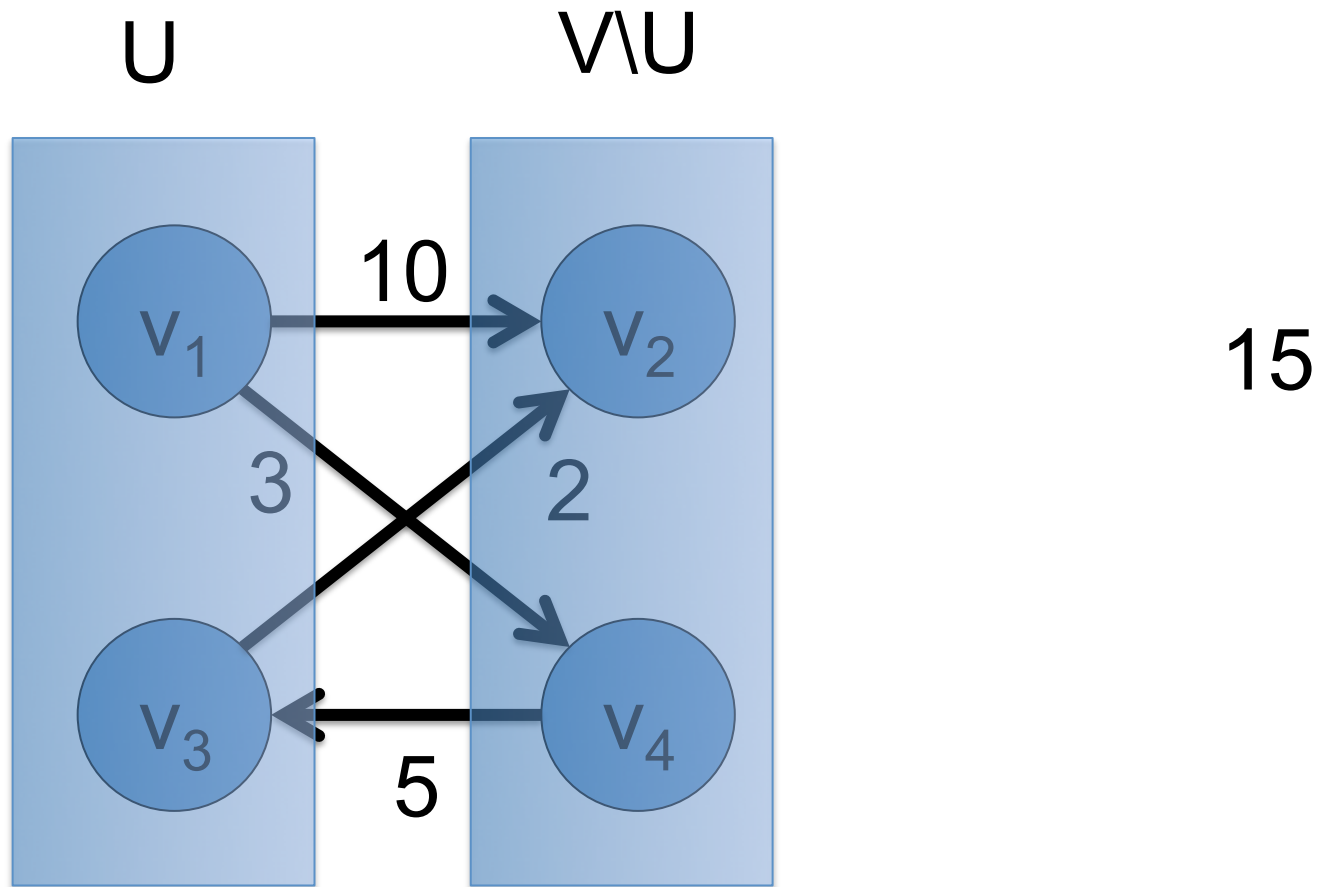


$$\sum_{a \in C} c(a)$$

Capacity of Cut

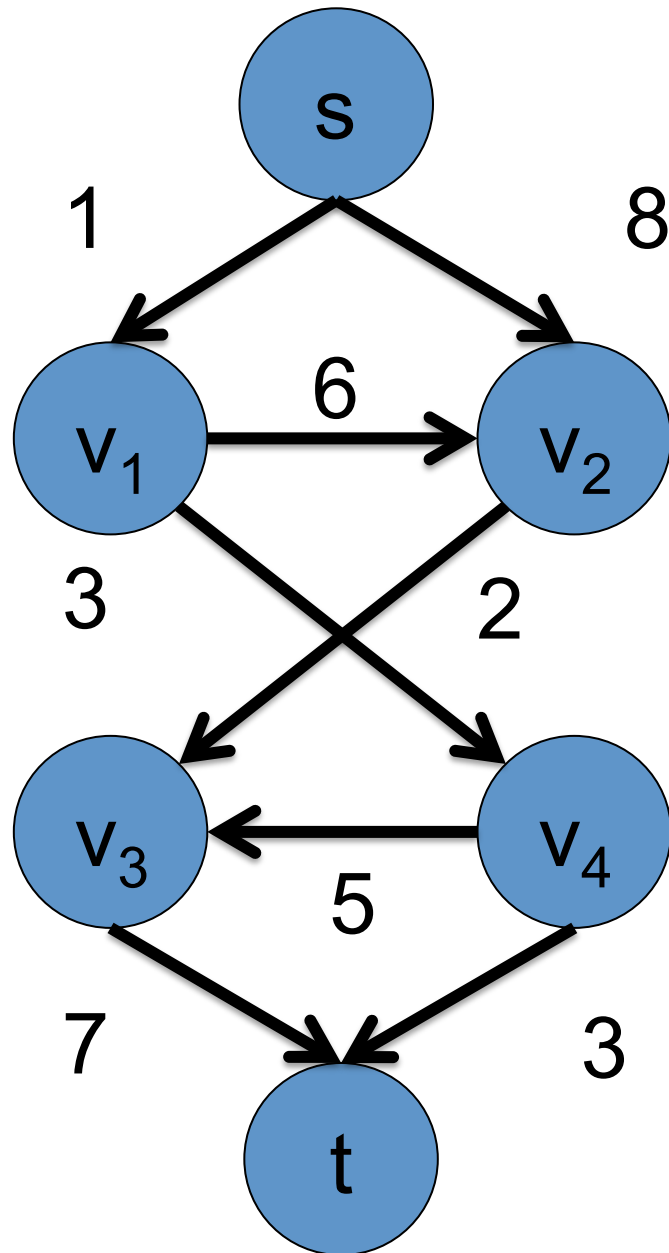


Capacity of Cut



s-t Cut

$$D = (V, A)$$



A source vertex “s”

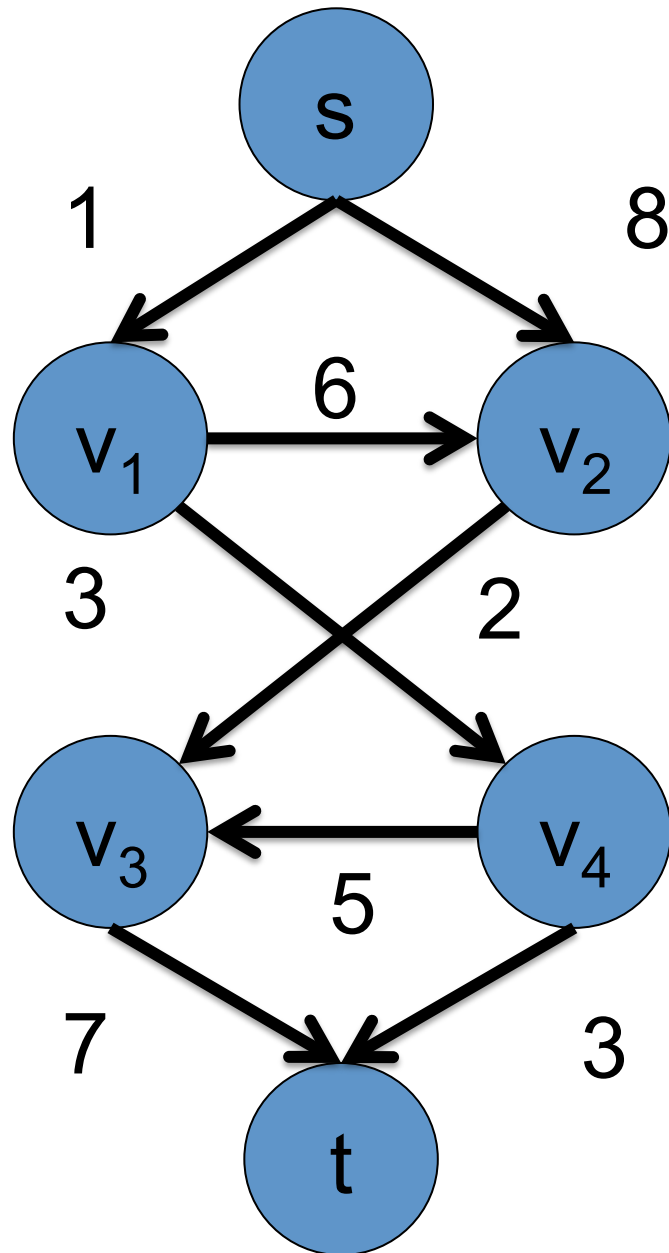
A sink vertex “t”

C is a cut such that

- $s \in U$
- $t \in V \setminus U$

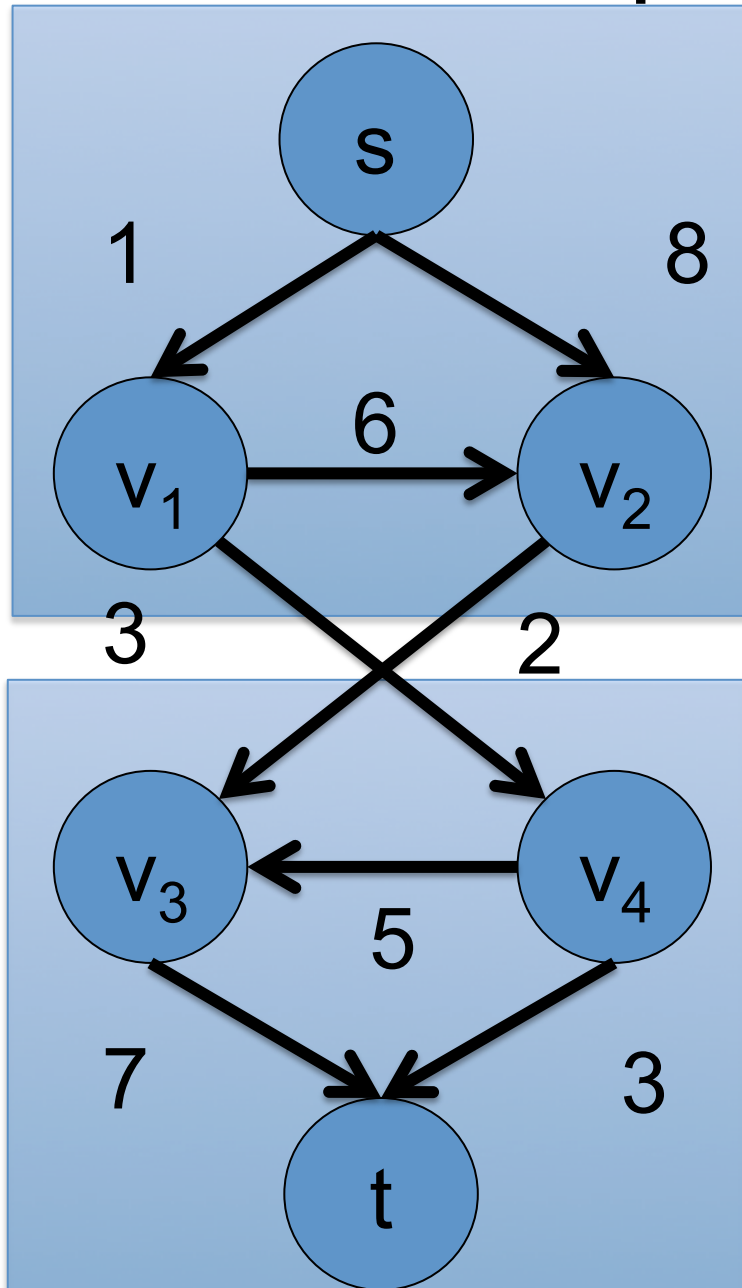
C is an s-t cut

Capacity of s-t Cut



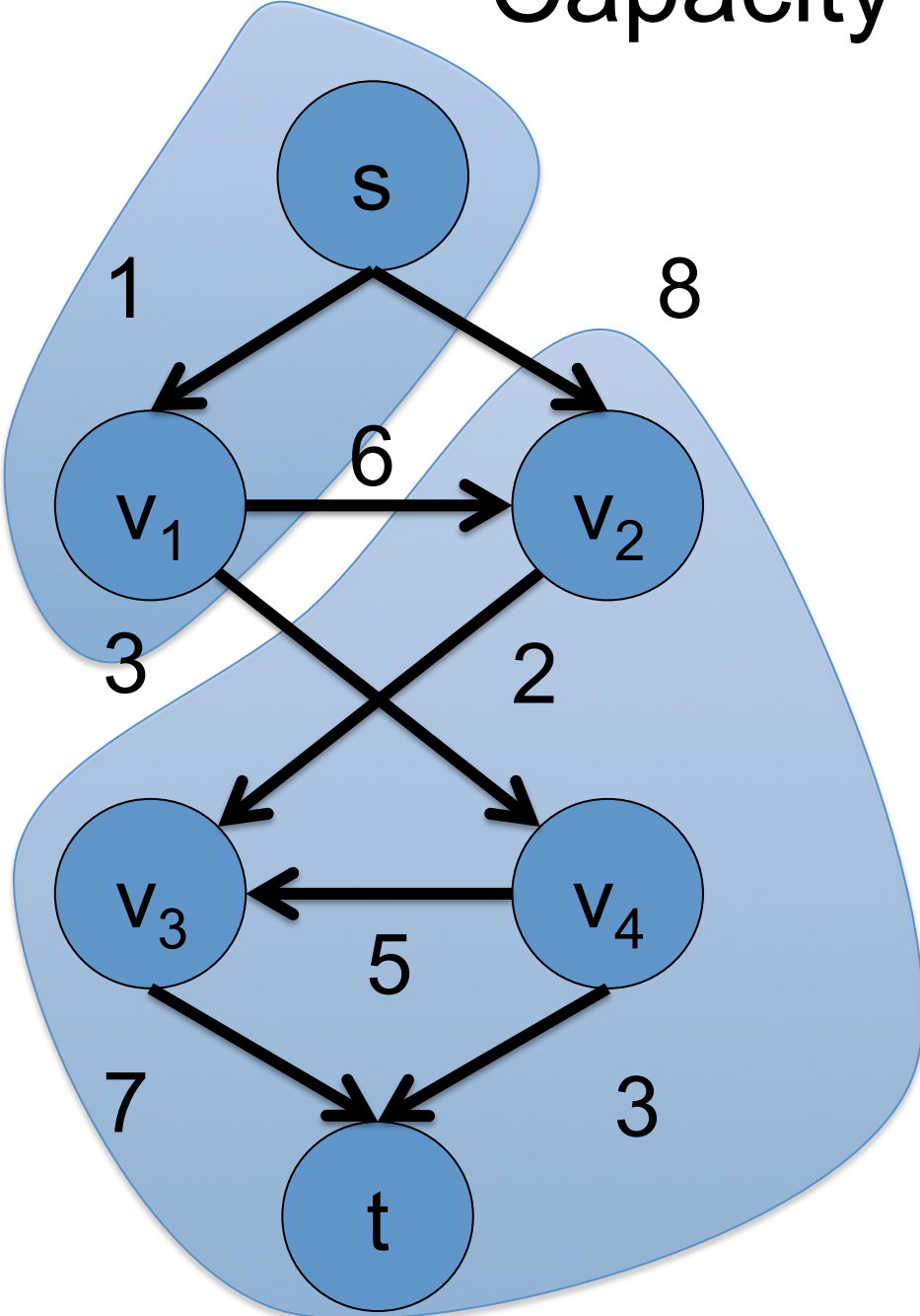
$$\sum_{a \in C} c(a)$$

Capacity of s-t Cut



5

Capacity of s-t Cut



17

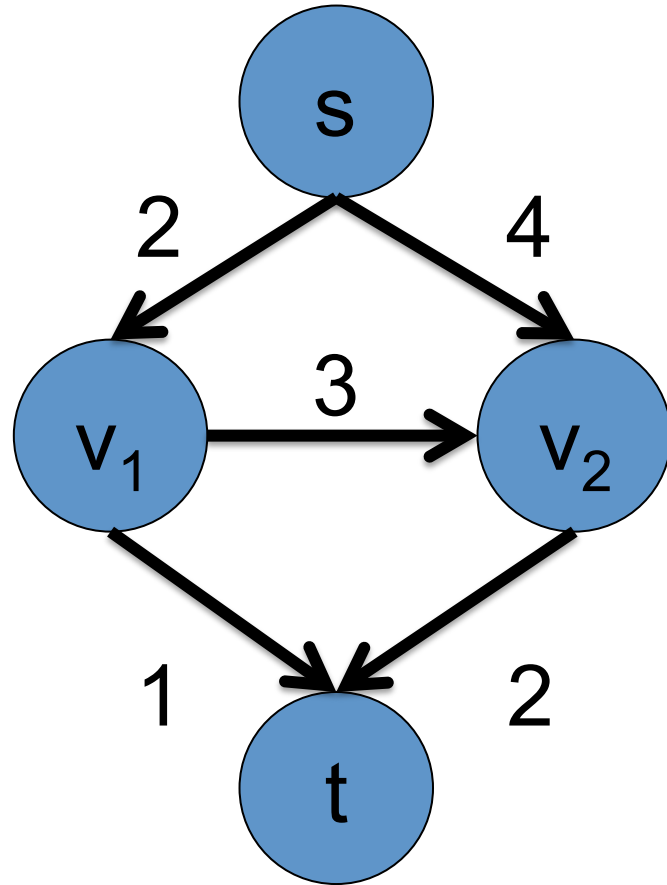
Outline

- Preliminaries
 - s-t Flow
 - s-t Cut
 - **Flows vs. Cuts**
- Maximum Flow
- Algorithms
- Energy minimization with max flow/min cut

Outline

- Preliminaries
- **Maximum Flow**
 - Residual Graph
 - Max-Flow Min-Cut Theorem
- Algorithms
- Energy minimization with max flow/min cut

Maximum Flow Problem



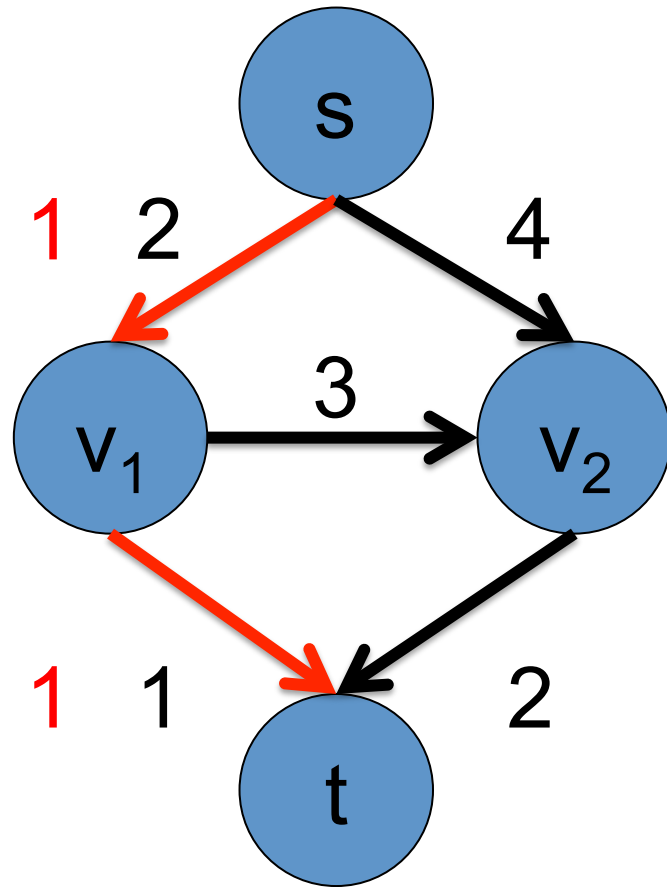
Find the flow with the maximum value !!

$$\sum_{(s,v) \in A} \text{flow}((s,v))$$

$$- \sum_{(u,s) \in A} \text{flow}((u,s))$$

First suggestion to solve this problem !!

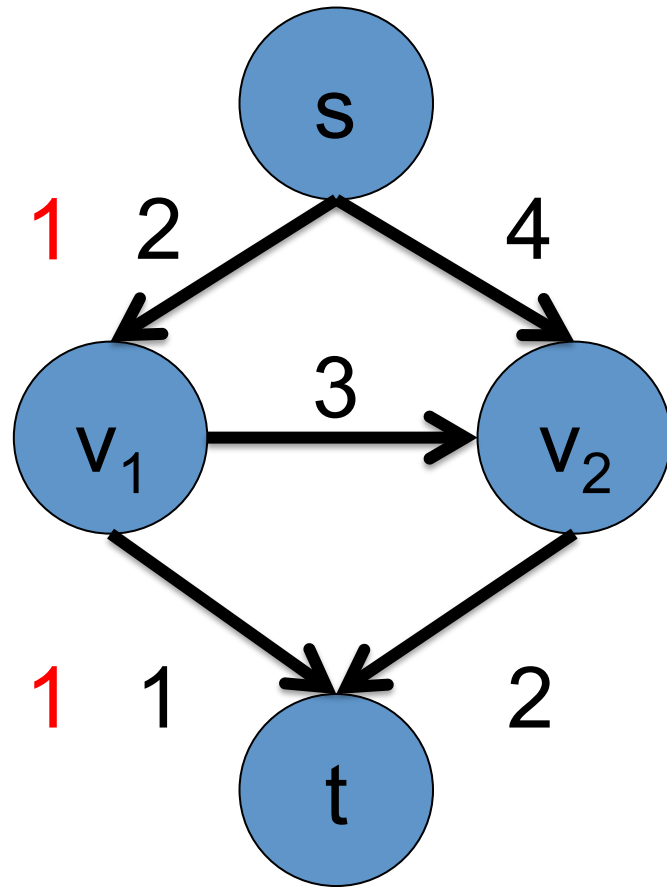
Passing Flow through s-t Paths



Find an s-t path where
 $\text{flow}(a) < c(a)$ for all arcs

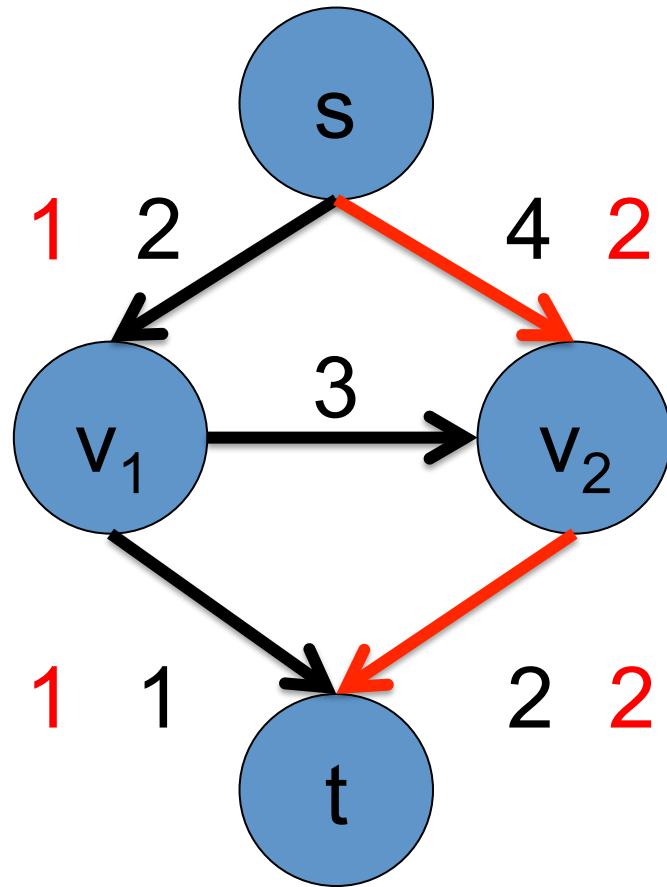
Pass maximum allowable
flow through the arcs

Passing Flow through s-t Paths



Find an s-t path where
 $\text{flow}(a) < c(a)$ for all arcs

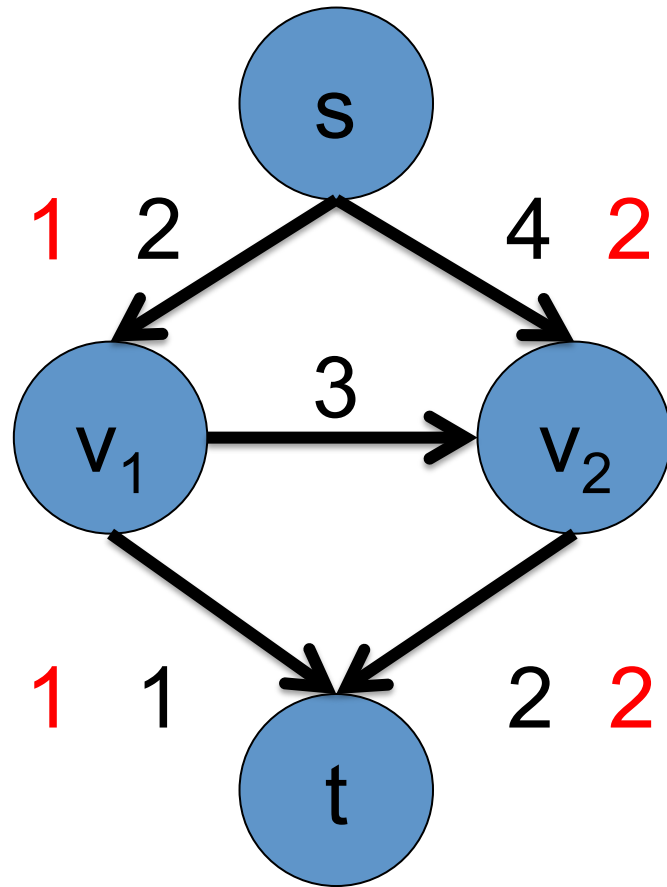
Passing Flow through s-t Paths



Find an s-t path where
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Pass maximum allowable
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Passing Flow through s-t Paths



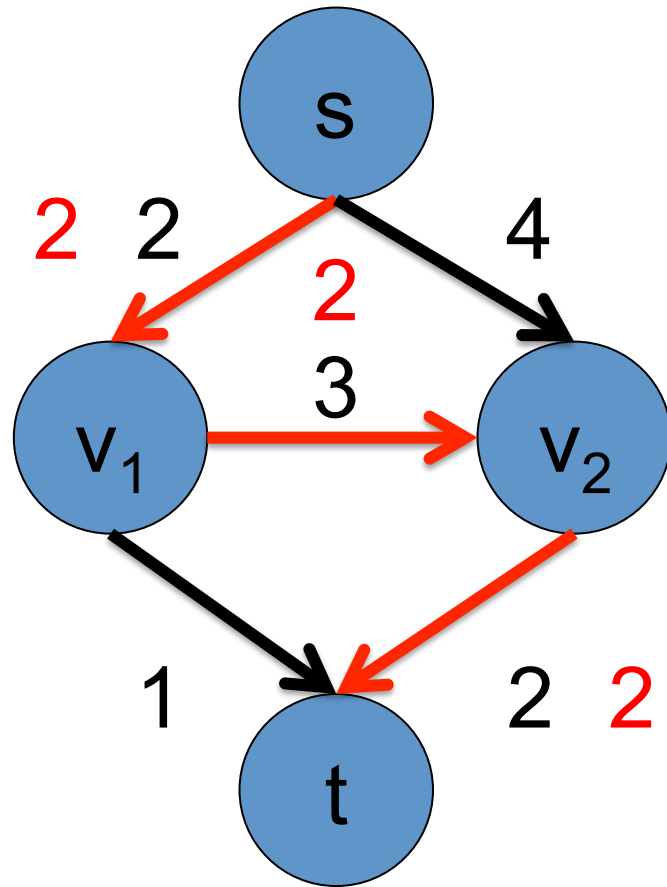
Find an s-t path where
 $\text{flow}(a) < c(a)$ for all arcs

No more paths. Stop.

Will this give us maximum flow?

NO !!!

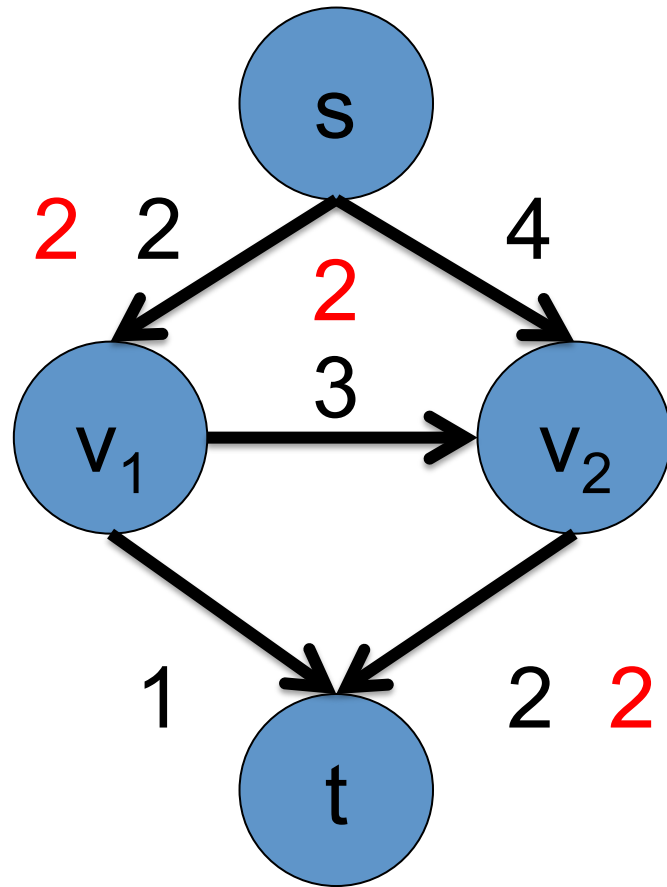
Passing Flow through s-t Paths



Find an s-t path where
 $\text{flow}(a) < c(a)$ for all arcs

Pass maximum allowable
flow through the arcs

Passing Flow through s-t Paths



Find an s-t path where
 $\text{flow}(a) < c(a)$ for all arcs

No more paths. Stop.

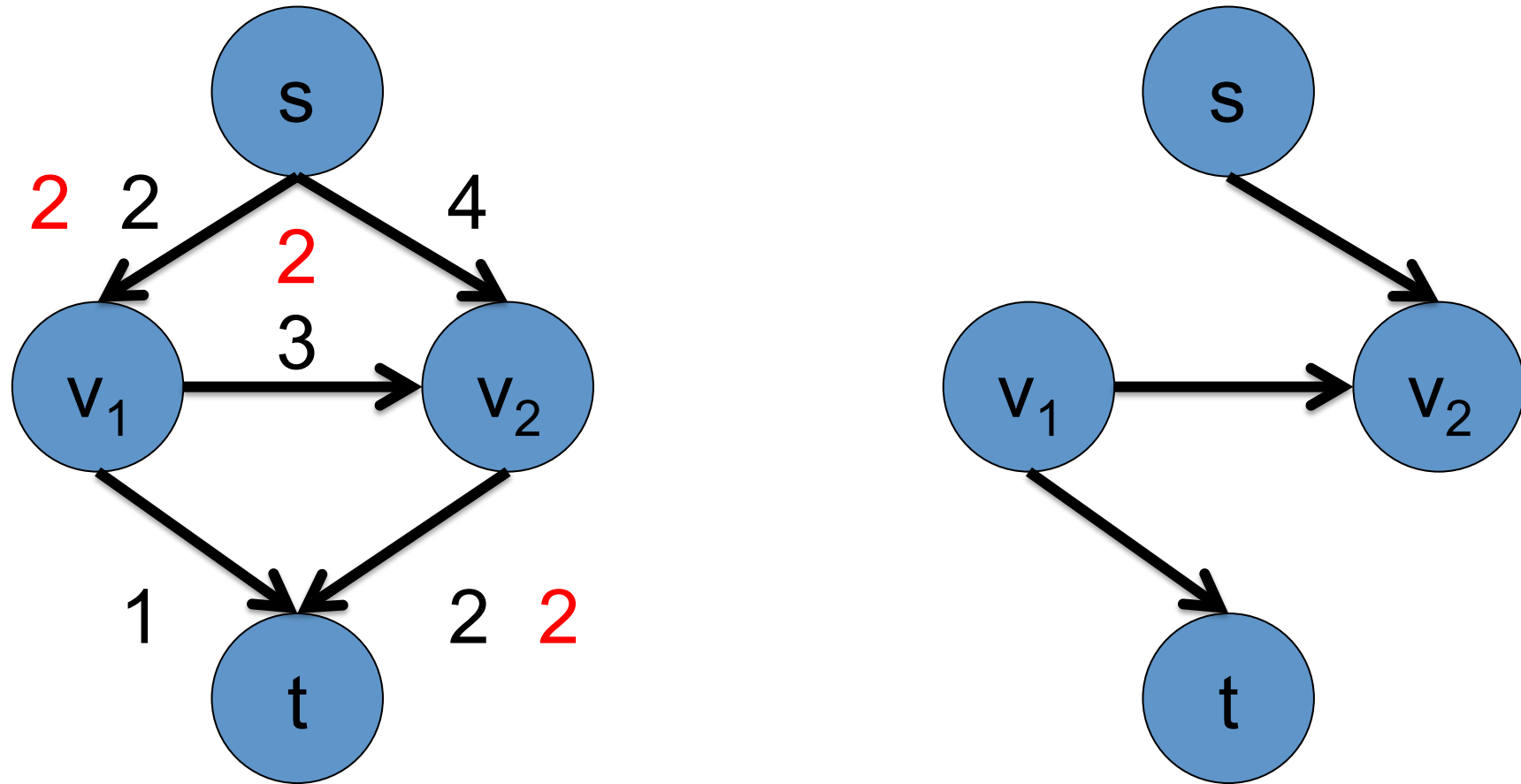
Another method?

Incorrect Answer !!

Outline

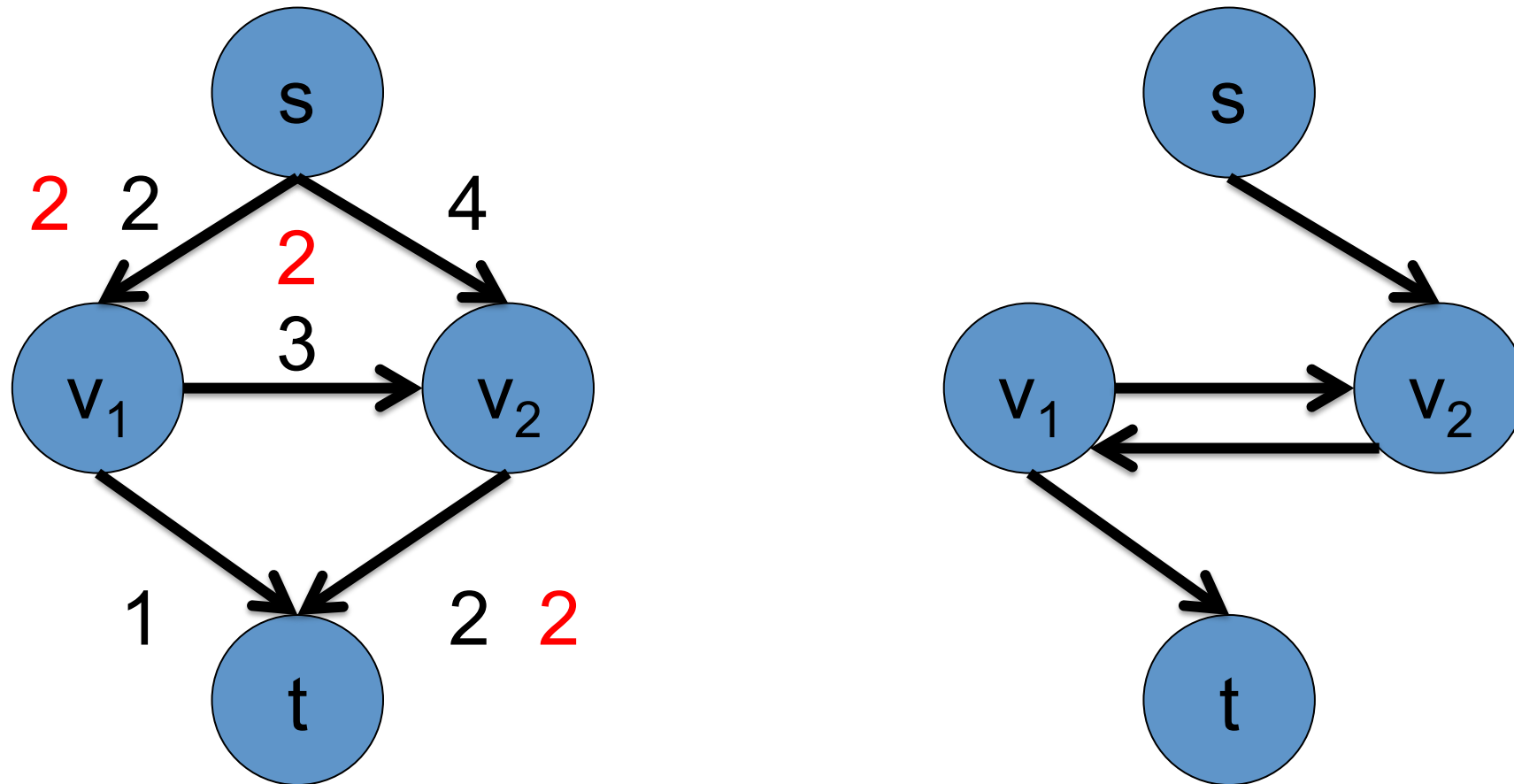
- Preliminaries
- Maximum Flow
 - **Residual Graph**
 - Max-Flow Min-Cut Theorem
- Algorithms
- Energy minimization with max flow/min cut

Residual Graph



Arcs where $\text{flow}(a) < c(a)$

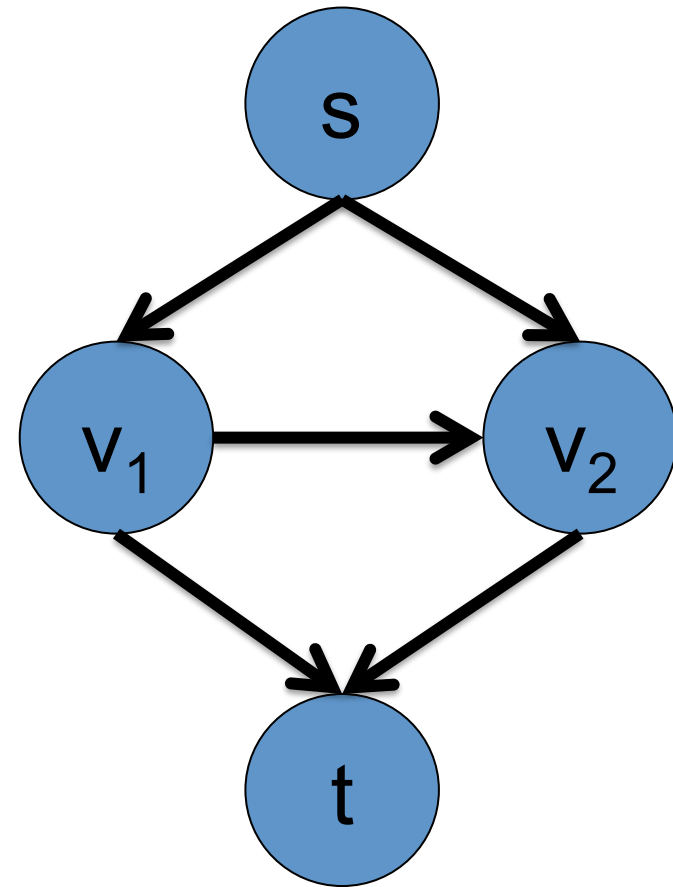
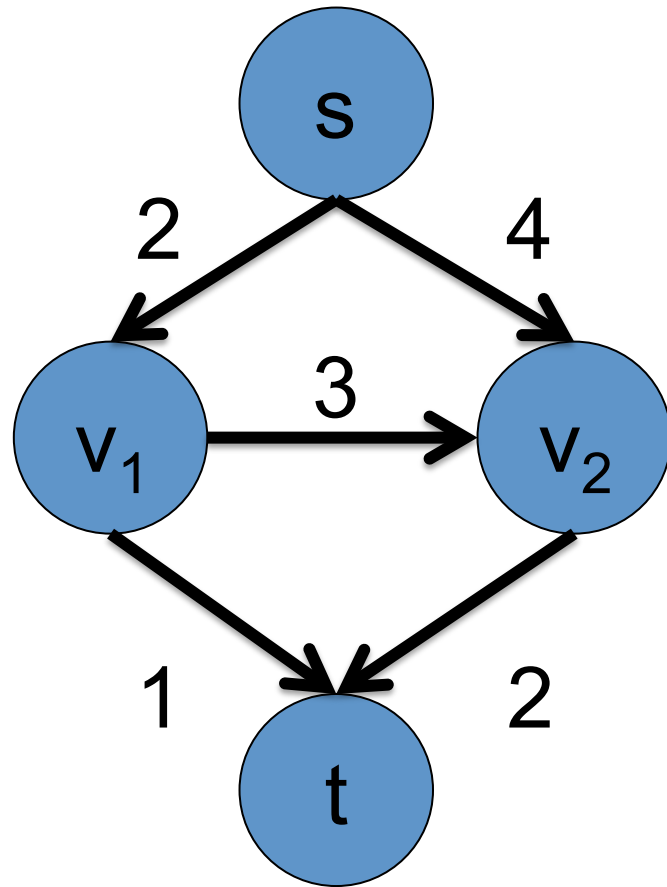
Residual Graph



Including arcs to s and from t is not necessary

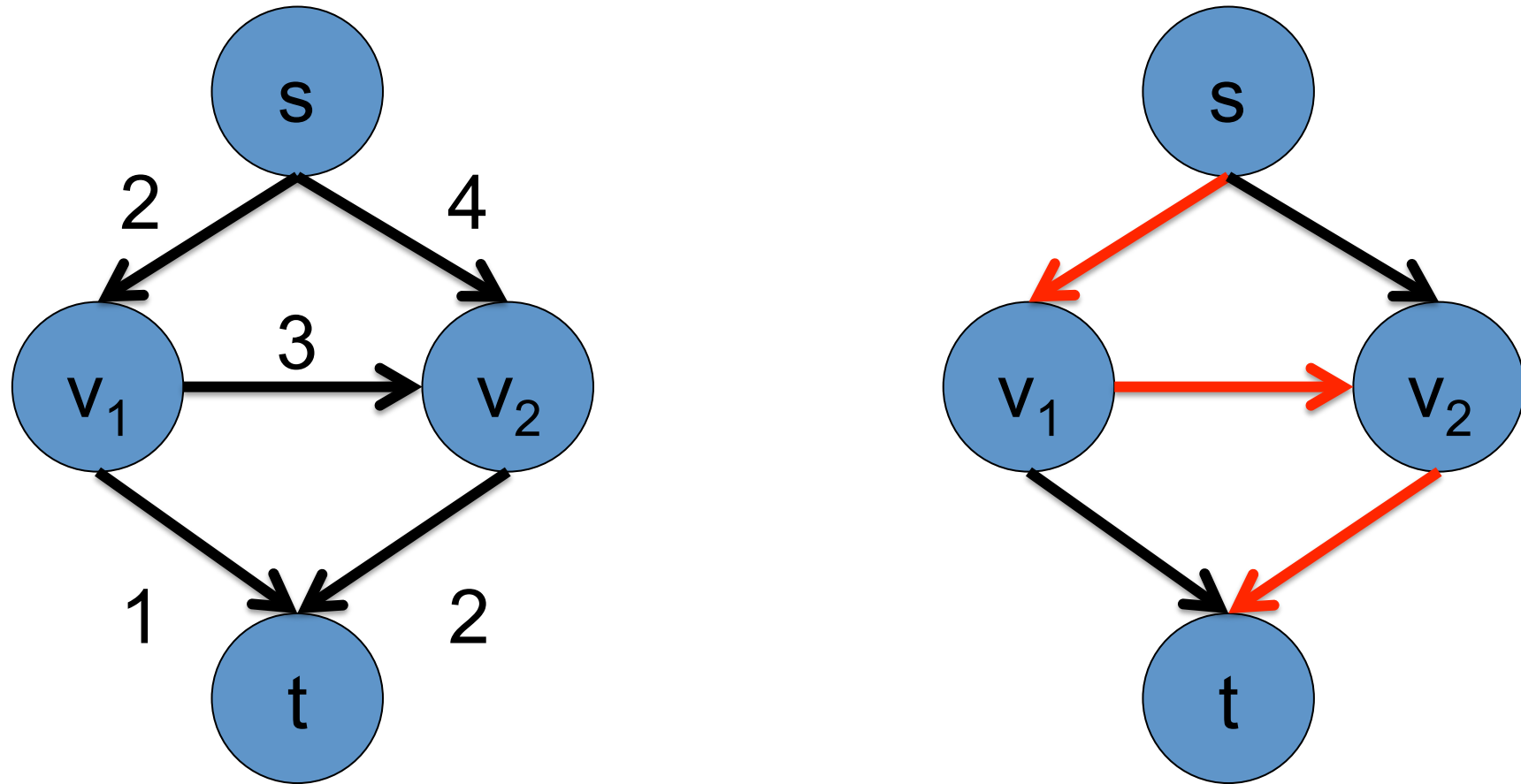
Inverse of arcs where $\text{flow}(a) > 0$

Maximum Flow using Residual Graphs



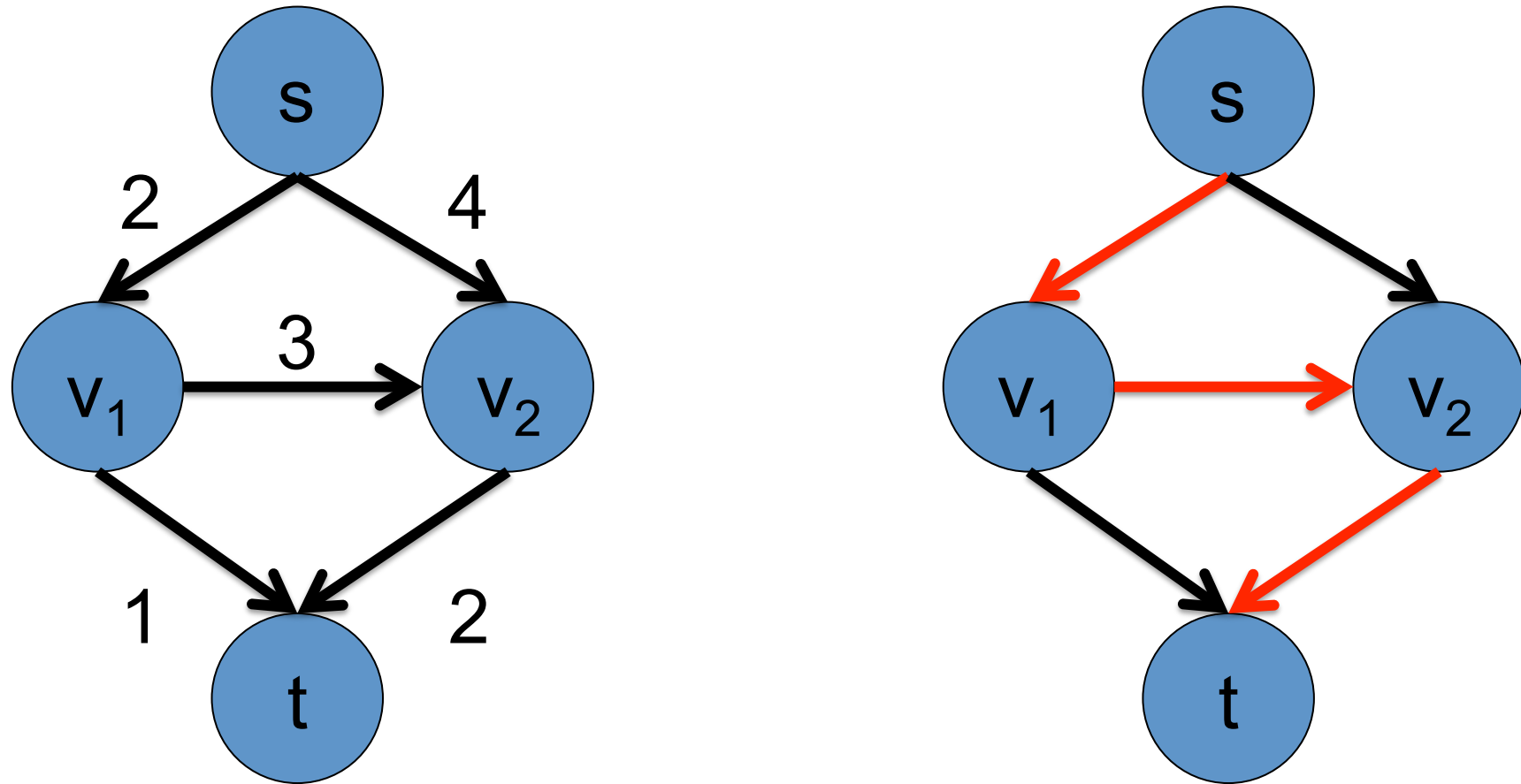
Start with zero flow.

Maximum Flow using Residual Graphs



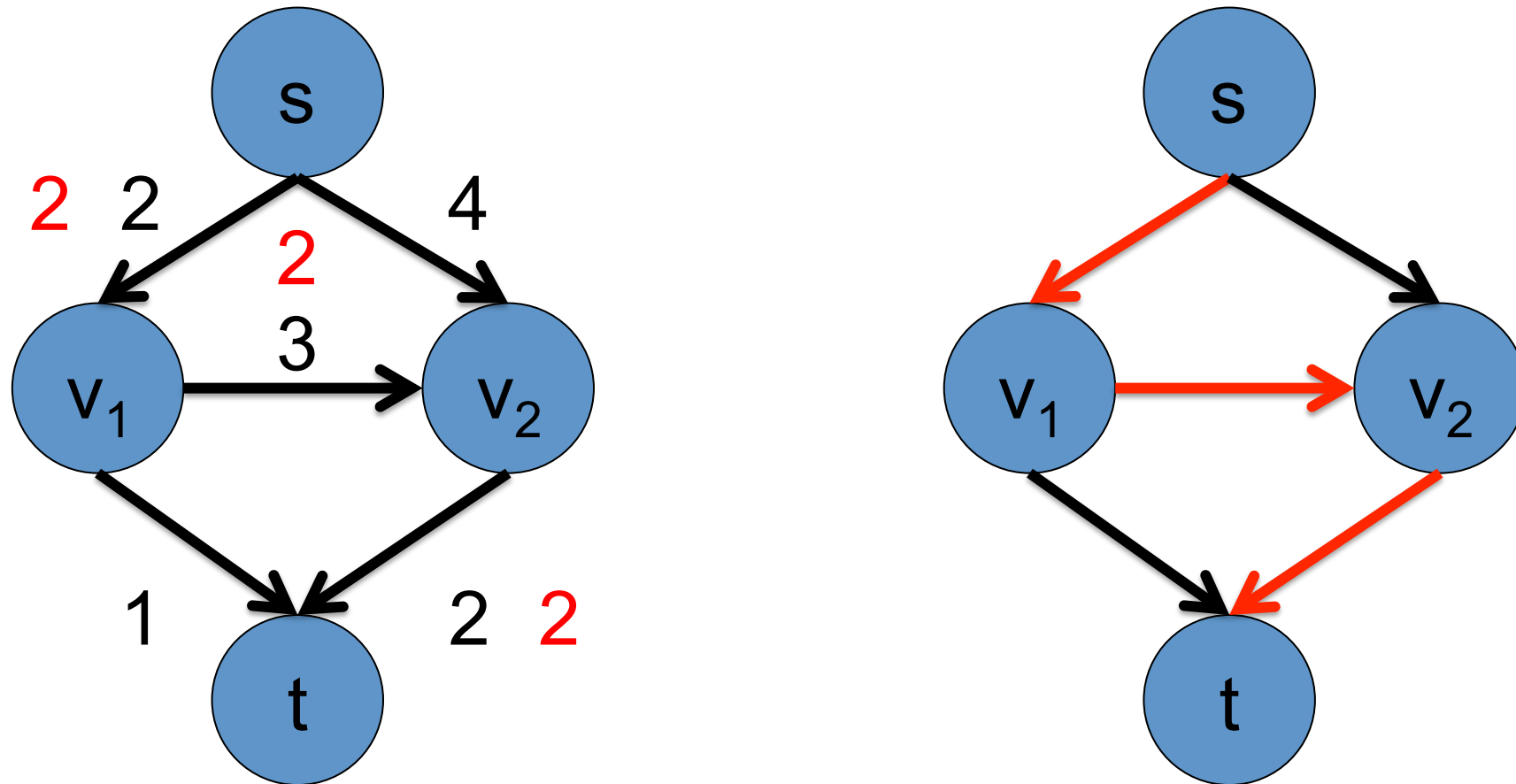
Find an s-t path in the residual graph.

Maximum Flow using Residual Graphs



For inverse arcs in path, subtract flow K .

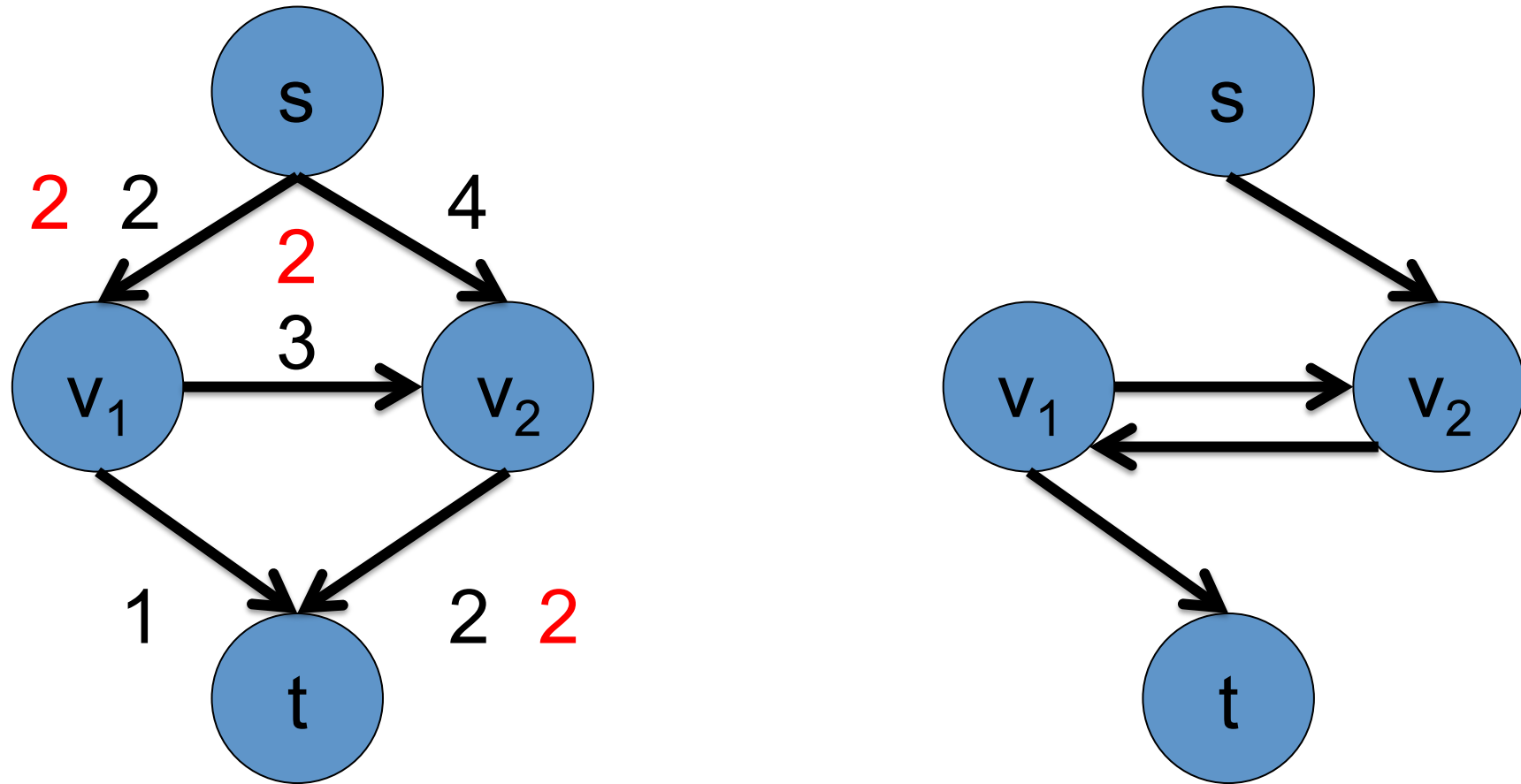
Maximum Flow using Residual Graphs



Choose maximum allowable value of K .

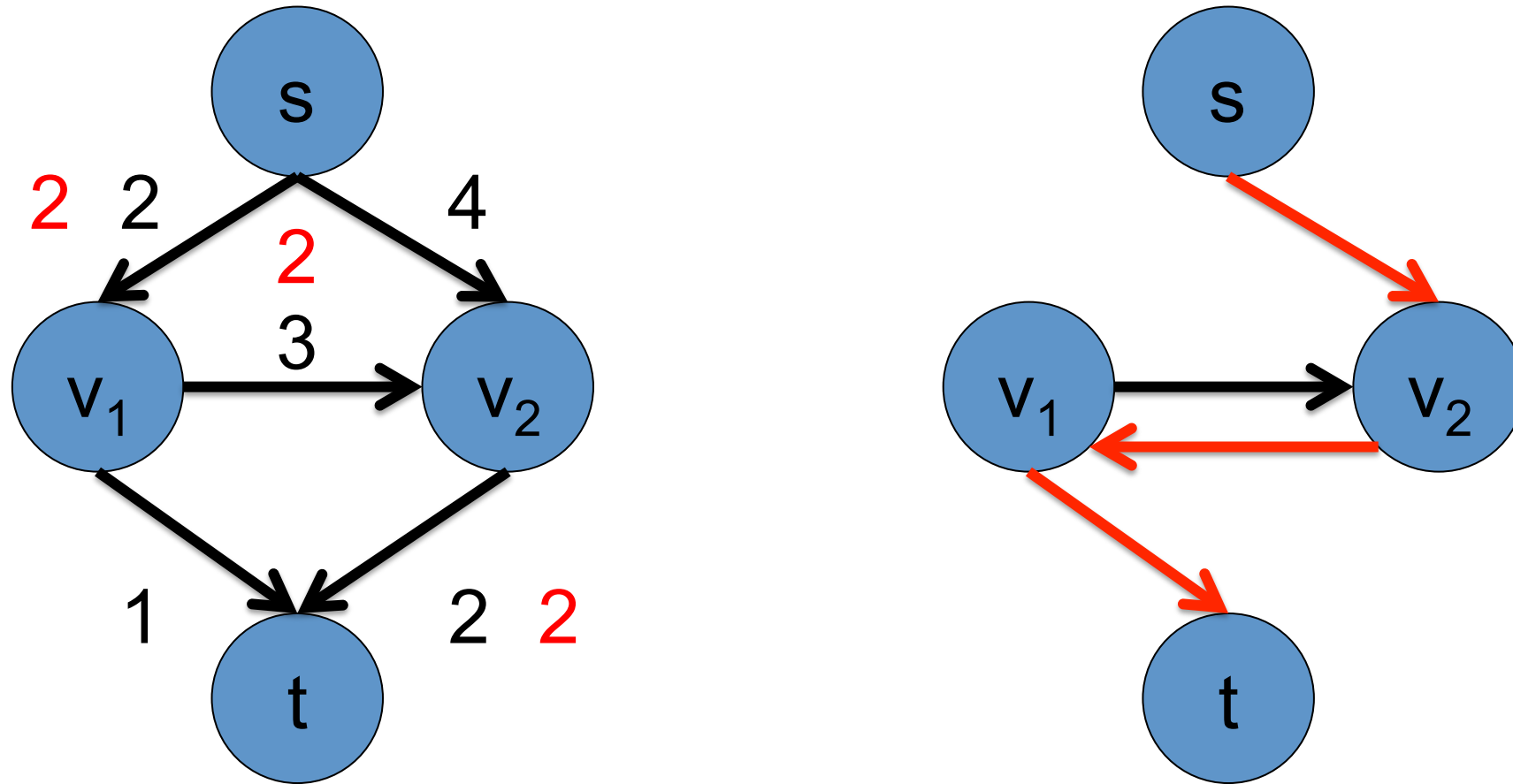
For forward arcs in path, add flow K .

Maximum Flow using Residual Graphs



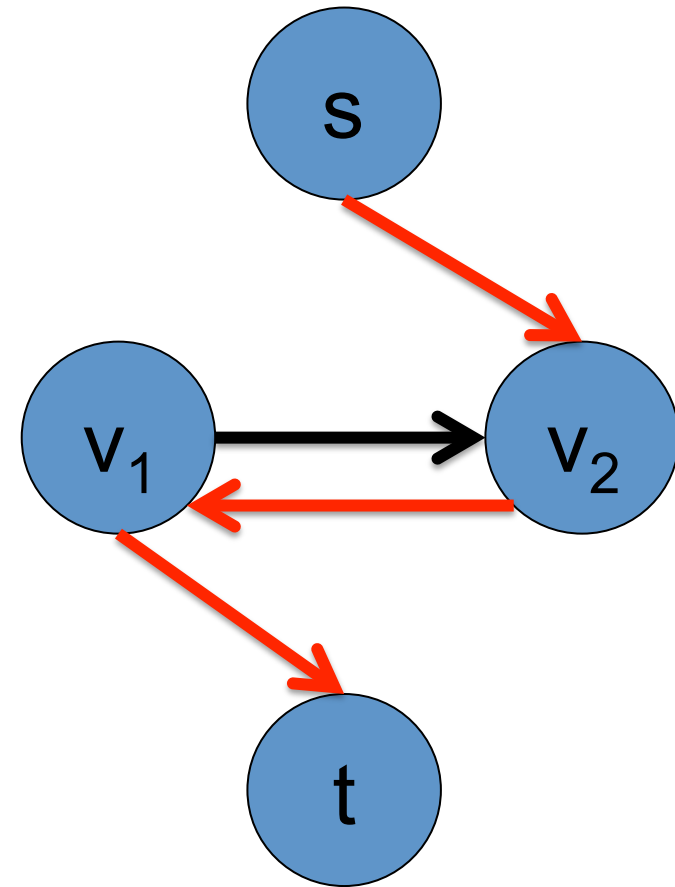
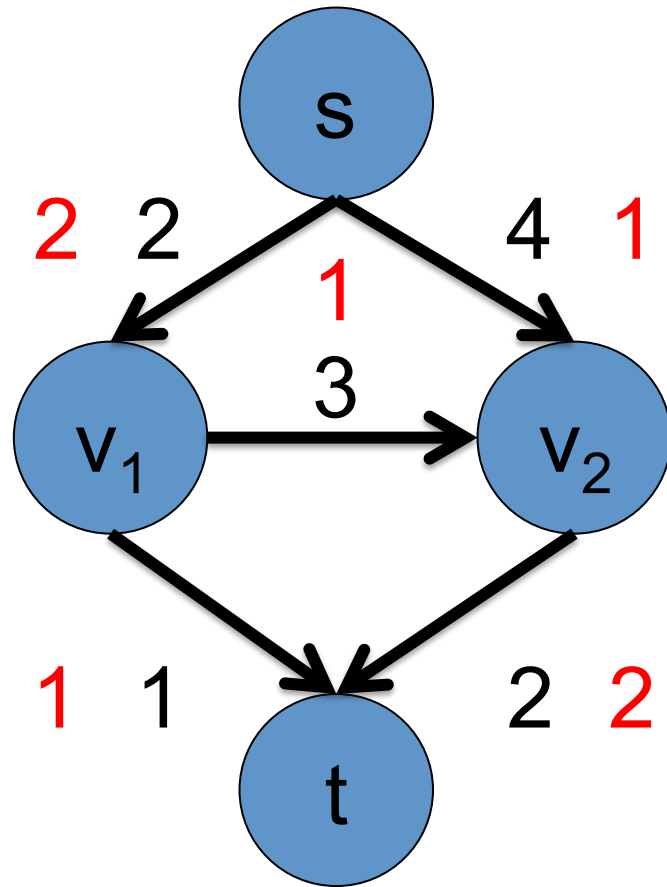
Update the residual graph.

Maximum Flow using Residual Graphs



Find an s-t path in the residual graph.

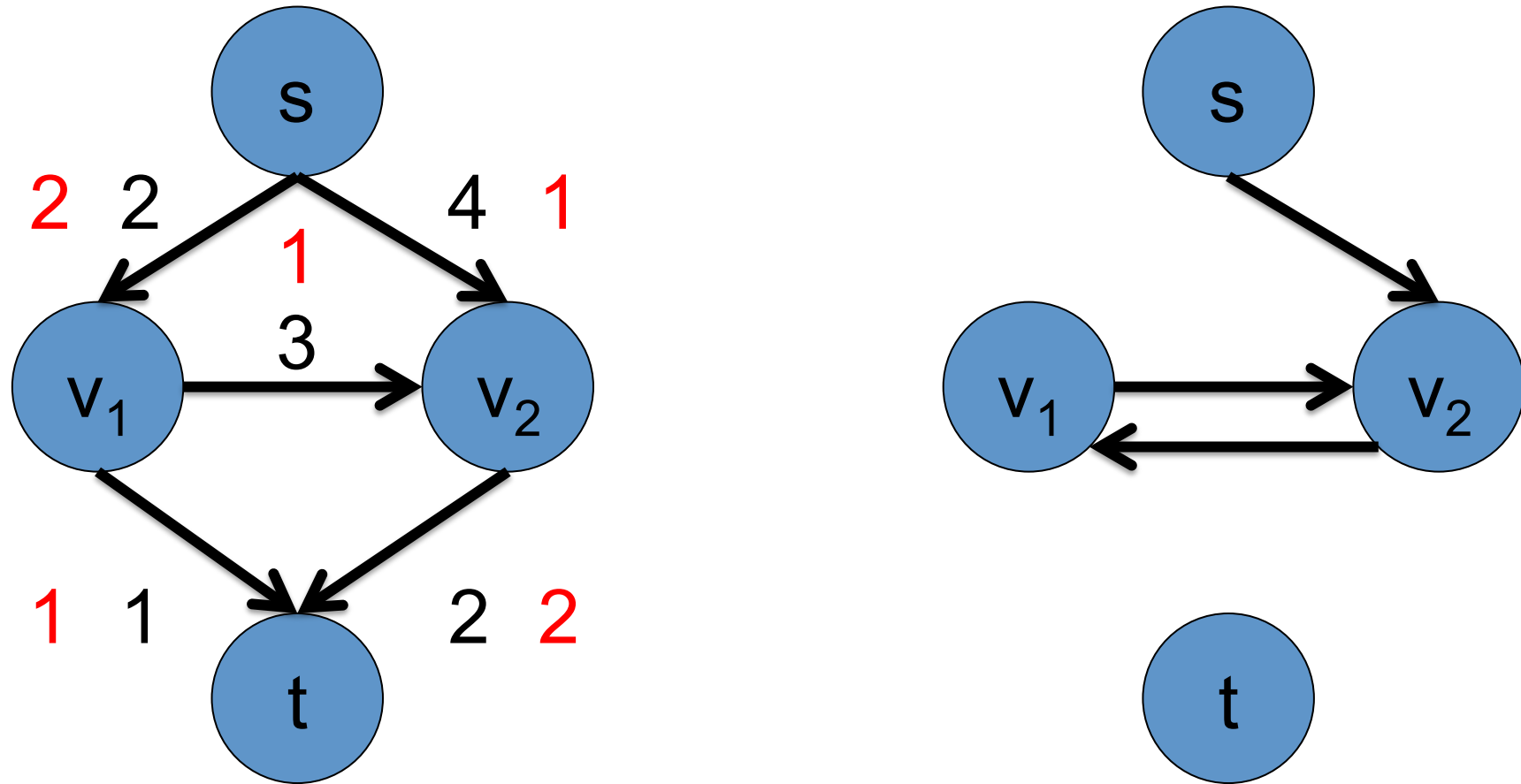
Maximum Flow using Residual Graphs



Choose maximum allowable value of K .

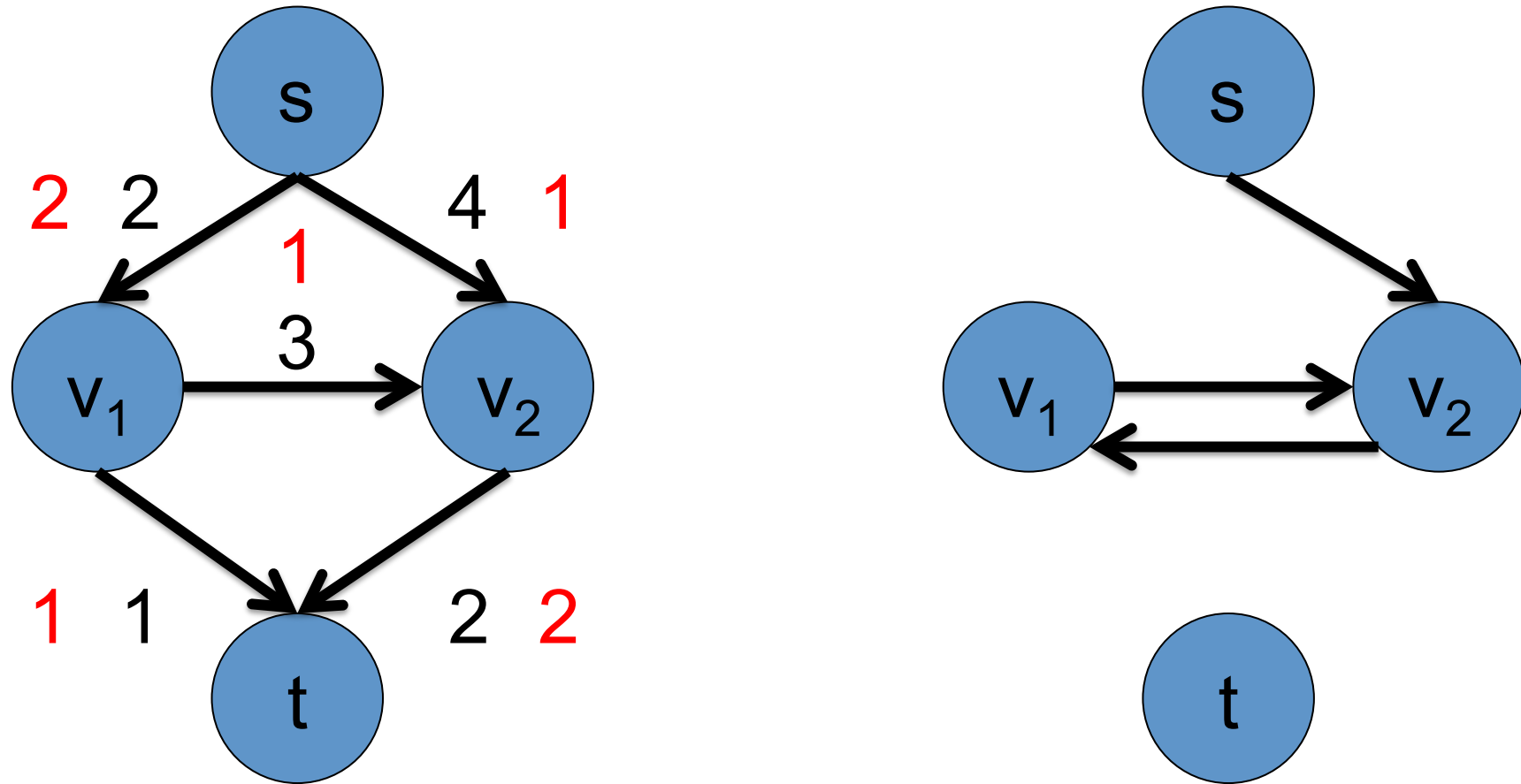
Add K to (s, v_2) and (v_1, t) . Subtract K from (v_1, v_2) .

Maximum Flow using Residual Graphs



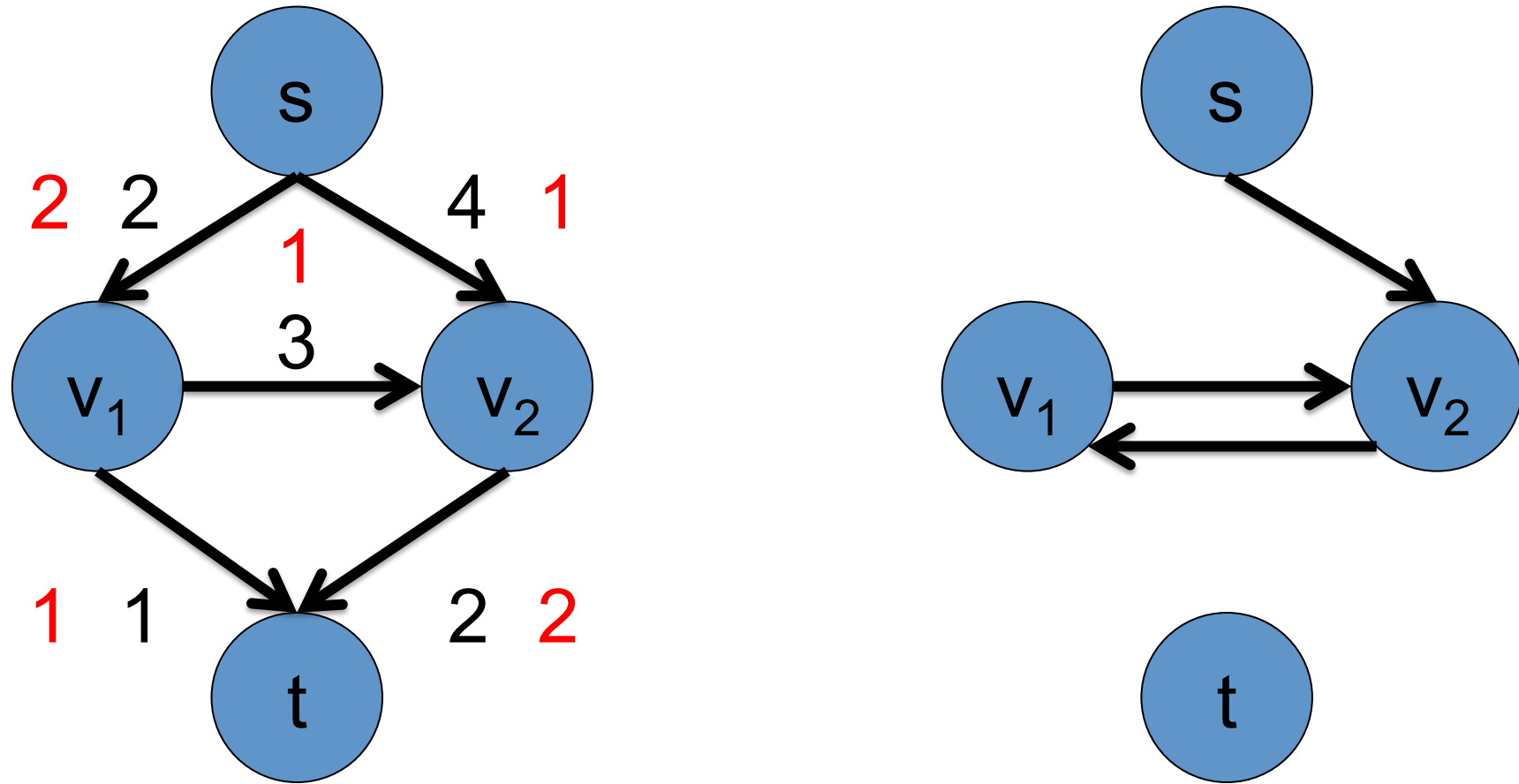
Update the residual graph.

Maximum Flow using Residual Graphs



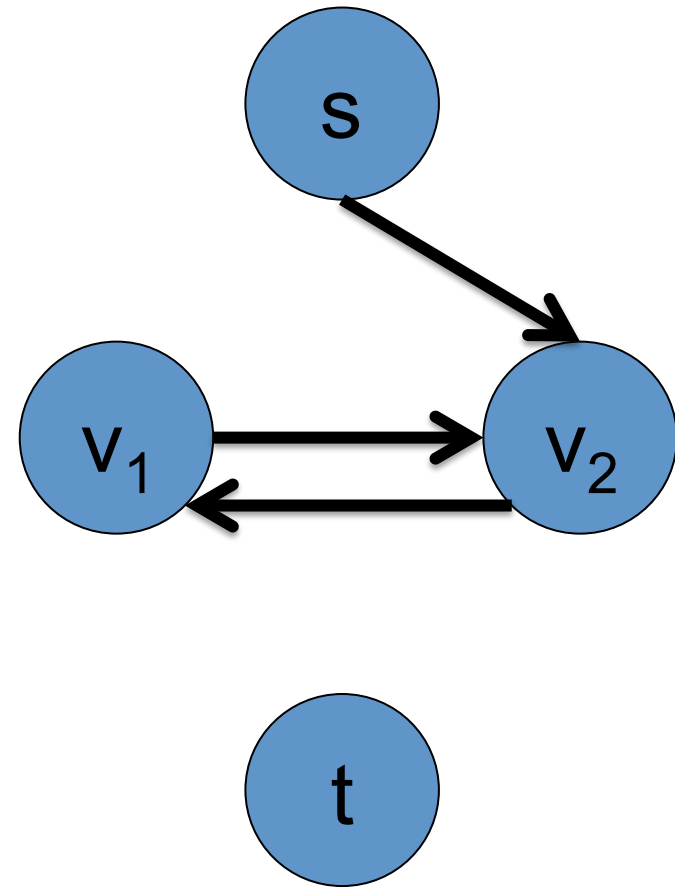
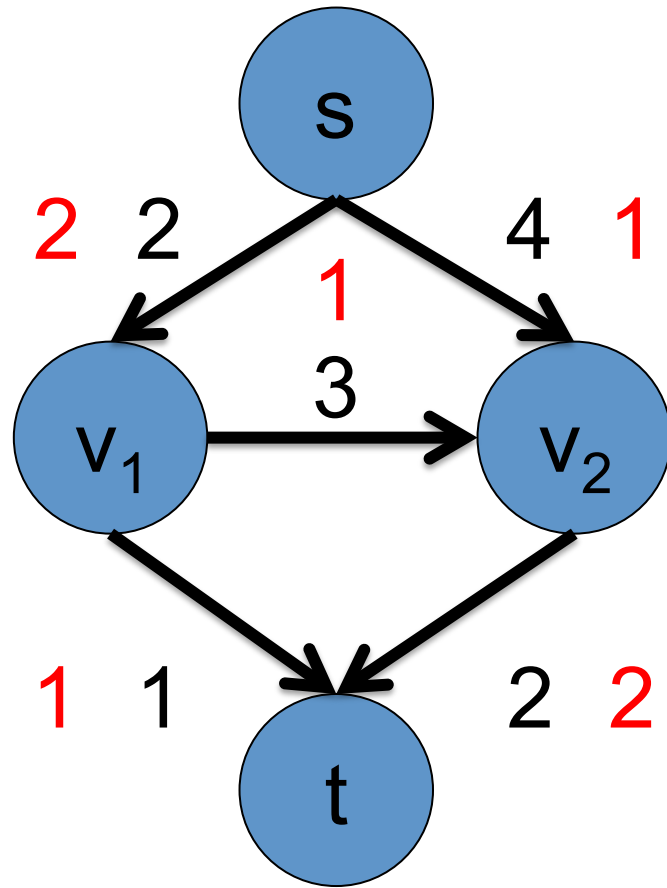
Find an s-t path in the residual graph.

Maximum Flow using Residual Graphs



No more s-t paths. Stop.

Maximum Flow using Residual Graphs



Correct Answer.

Outline

- Preliminaries
- Maximum Flow
 - Residual Graph
 - Max-Flow Min-Cut Theorem
- **Algorithms**
- Energy minimization with max flow/min cut

History of Maxflow Algorithms

Augmenting Path and Push-Relabel

n: #nodes

m: #edges

U: maximum edge weight

year	discoverer(s)	bound
1951	Dantzig	$O(n^2mU)$
1955	Ford & Fulkerson	$O(m^2U)$
1970	Dinitz	$O(n^2m)$
1972	Edmonds & Karp	$O(m^2 \log U)$
1973	Dinitz	$O(nm \log U)$
1974	Karzanov	$O(n^3)$
1977	Cherkassky	$O(n^2m^{1/2})$
1980	Galil & Naamad	$O(nm \log^2 n)$
1983	Sleator & Tarjan	$O(nm \log n)$
1986	Goldberg & Tarjan	$O(nm \log(n^2/m))$
1987	Ahuja & Orlin	$O(nm + n^2 \log U)$
1987	Ahuja et al.	$O(nm \log(n\sqrt{\log U/m}))$
1989	Cheriyān & Hagerup	$E(nm + n^2 \log^2 n)$
1990	Cheriyān et al.	$O(n^3 / \log n)$
1990	Alon	$O(nm + n^{8/3} \log n)$
1992	King et al.	$O(nm + n^{2+\epsilon})$
1993	Phillips & Westbrook	$O(nm(\log_{m/n} n + \log^{2+\epsilon} n))$
1994	King et al.	$O(nm \log_{m/(n \log n)} n)$
1997	Goldberg & Rao	$O(m^{3/2} \log(n^2/m) \log U)$ $O(n^{2/3}m \log(n^2/m) \log U)$

**Algorithms
assume non-
negative edge
weights**

[Slide credit: Andrew Goldberg]

History of Maxflow Algorithms

Augmenting Path and Push-Relabel

n: #nodes

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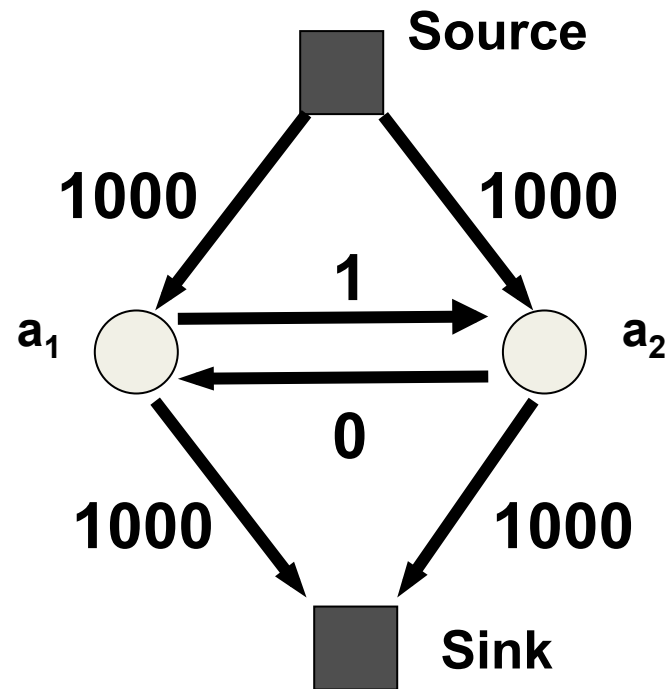
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**Algorithms
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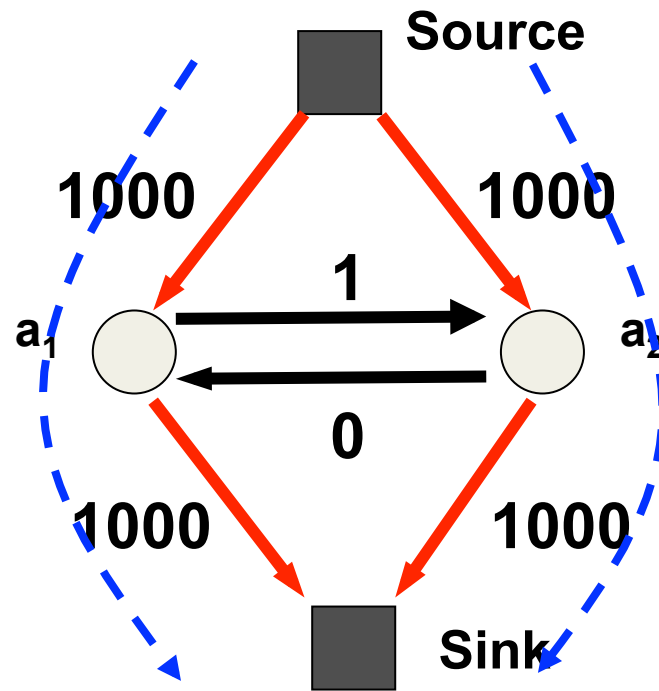
Augmenting Path based Algorithms

Ford Fulkerson: Choose **any** augmenting path



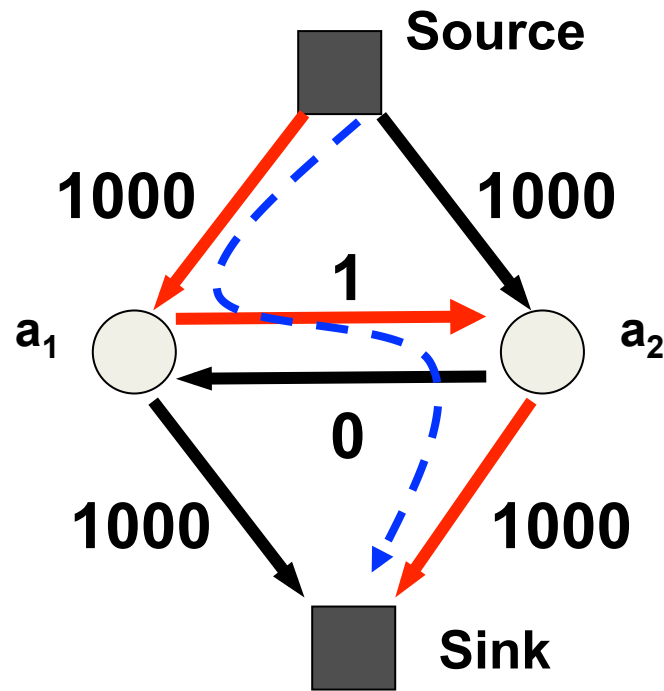
Augmenting Path based Algorithms

Ford Fulkerson: Choose **any** augmenting path



Augmenting Path based Algorithms

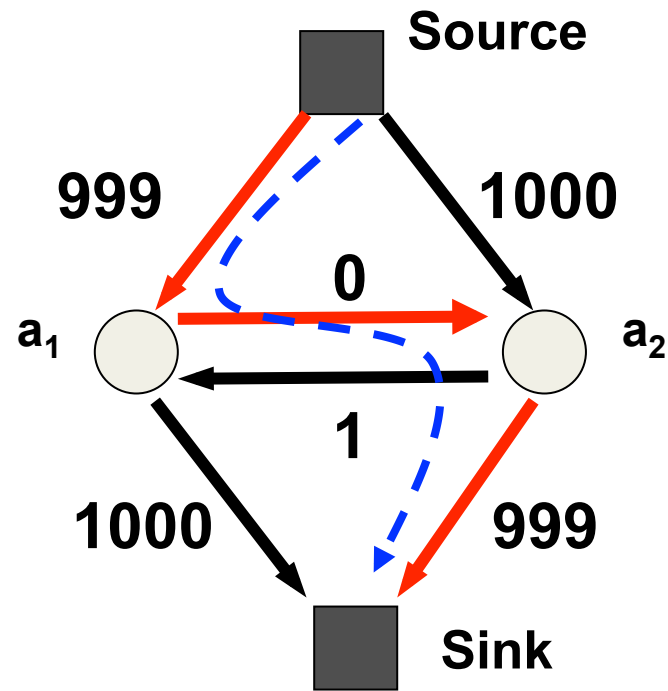
Ford Fulkerson: Choose **any** augmenting path



**Bad
Augmenting
Path**

Augmenting Path based Algorithms

Ford Fulkerson: Choose **any** augmenting path

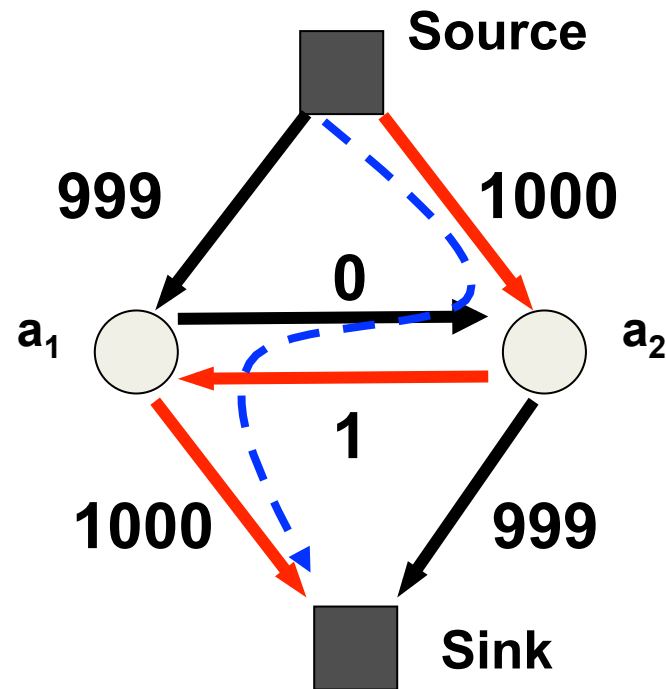


Augmenting Path based Algorithms

n: #nodes

m: #edges

Ford Fulkerson: Choose **any** augmenting path



We will have to perform 2000 augmentations!

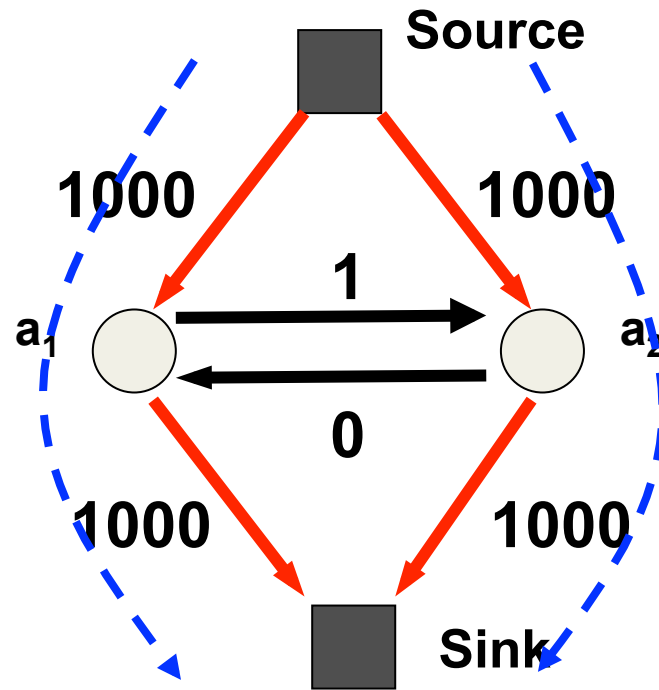
Worst case complexity: $O(m \times \text{Total_Flow})$
(Pseudo-polynomial bound: depends on flow)

Augmenting Path based Algorithms

n: #nodes

m: #edges

Dinitz: Choose **shortest** augmenting path

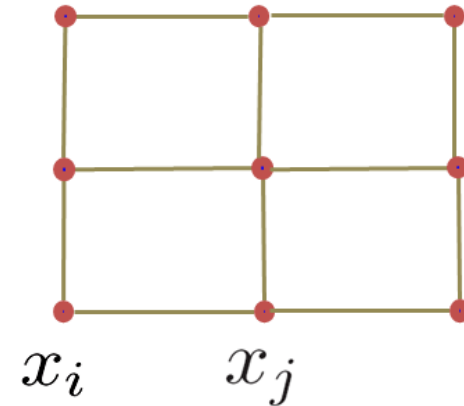


Worst case complexity: $O(m n^2)$

Maxflow in Computer Vision

- **Specialized algorithms for vision problems**

- Grid graphs
- Low connectivity ($m \sim O(n)$)



- **Dual search tree augmenting path algorithm**

[Boykov and Kolmogorov PAMI 2004]

- Finds approximate shortest augmenting paths efficiently
- High worst-case time complexity
- Empirically outperforms other algorithms on vision problems

- **Efficient code available on the web**

e.g., <http://pub.ist.ac.at/~vnk/software.html>

